

MUTUAL ALGEBRAICITY AND CELLULARITY

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ABSTRACT. We prove two results intended to streamline proofs about cellularity that pass through mutual algebraicity. First, we show that a countable structure M is cellular if and only if M is ω -categorical and mutually algebraic. Second, if a countable structure M in a finite relational language is mutually algebraic non-cellular, we show it admits an elementary extension adding infinitely many infinite MA-connected components.

Towards these results, we introduce MA-presentations of a mutually algebraic structure, in which every atomic formula is mutually algebraic. This allows for an improved quantifier elimination and a decomposition of the structure into independent pieces. We also show this decomposition is largely independent of the MA-presentation chosen.

1. INTRODUCTION

A characterization of mutual algebraicity, implying that in many settings non-mutually algebraic structures are complicated, was given in [LT1]. However, in most cases concerning the combinatorics of countable structures or of hereditary classes, the dividing line in “many models” behavior is not given by mutual algebraicity, but by the stronger condition of cellularity. Examples include counting the number of countable structures of a given age [MPW], counting the finite models of a given size in a hereditary class [LT2], and counting the number of substructures of a countable structure up to isomorphism [LM].

In this paper, we provide results describing the gap between mutual algebraicity and cellularity, leading to a general strategy for showing cellularity is a dividing line in a given problem. This strategy is explicitly employed in the companion paper [BLSib] to count the number of structures bi-embeddable with a given countable structure. The first step in this strategy is to produce complicated behavior in the non-mutually algebraic case, likely using the Ryll-Nardzewski-type characterization of mutual algebraicity in a finite relational language from [LT1]. Our first result shows that in the ω -categorical setting, this is already enough.

Theorem 1.1 (Theorem 3.7). *Suppose M is a countable structure in a countable language. Then M is cellular if and only if M is mutually algebraic and ω -categorical.*

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The next step of the strategy is to show that if M is mutually algebraic but not cellular, then models of its theory are sufficiently complex. This is the main part of the paper. At first blush, it is not clear that cellularity and mutual algebraicity are related. By definition, a cellular structure M admits a natural partition into finite pieces. Depending on the language a mutually algebraic structure M is presented in, such a partition might not be readily describable. As examples, if we take T to be the theory of $(\mathbb{Z}, succ)$, the theory of the integers with a binary successor relation, then any model M of T is mutually algebraic, and M is naturally partitioned into its ‘ \mathbb{Z} -classes’. The salient feature is that the binary successor relation is mutually algebraic. By contrast, the arguably simpler theory of $Th(M, E)$, where E is an equivalence relation on M with two classes, both infinite, is mutually algebraic (in fact, cellular), yet the atomic formula $E(x, y)$ is not mutually algebraic. In this case, the cellular partition of M is into singletons, with singletons from the two classes distinguished, which requires naming a constant. In Section 2, we introduce the notion of an *MA-presentation*, in which all atomic relations are mutually algebraic. We show that MA-presented structures M admit a stronger quantifier elimination result and there is a natural equivalence relation \sim that partitions M as in the definition of cellularity. We also define a notion of two structures with the same universe but in different languages to be *associated* and we prove that every mutually algebraic structure is associated to an MA-presented structure. Although there may be several MA-presented structures associated to a given mutually algebraic M , the corresponding partitions into \sim -classes are largely the same.

In Section 3, we use the fact that any mutually algebraic M has many associated MA-presented structures to give several characterizations of cellularity, e.g., a mutually algebraic M is cellular if and only if there is a uniform finite bound on the size of \sim -classes in some/every associated MA-presented M' . As a byproduct of this analysis, we prove the following result, which is used by the authors in [?BLSib].

Theorem 1.2 (Theorem 3.10). *Let \mathcal{L} be finite relational, and suppose M is mutually algebraic but non-cellular countable \mathcal{L} -structure. Then there is some $M^* \succ M$ such that M^* contains infinitely many new infinite MA-connected components, pairwise isomorphic over M . Furthermore, we may take the universe of M^* to be the universe of M together with these new components.*

Also in Section 3, we obtain other characterizations of cellularity, including the following, which does not mention partitions.

Theorem 1.3 (Theorem 3.9). *A countable structure M is monadically ω -categorical if and only if M is cellular.*

2. MUTUAL ALGEBRAICITY AND MA-PRESENTATIONS

Throughout this paper, we only consider languages \mathcal{L} with no function symbols (i.e., \mathcal{L} only has relation and constant symbols). As all the properties

we consider are preserved by passing between a function and its graph, this is not a significant restriction.

Definition 2.1. Given a structure M , an n -ary relation $R(\bar{x})$ is *mutually algebraic* if there is a constant K such that for each $m \in M$, the number of tuples $\bar{m} \in M^n$ such that $R(\bar{m})$ and $m \in \bar{m}$ is at most K .

Note that every unary relation is mutually algebraic.

Definition 2.2. Given an \mathcal{L} -structure M , let \mathcal{L}_M be \mathcal{L} expanded by constant symbols for every element of M . M is *mutually algebraic* if every \mathcal{L}_M -formula is equivalent to a boolean combination of mutually algebraic \mathcal{L}_M formulas.

Theorem 2.3 ([?MA, Theorem 3.3], [?MAO, Proposition 4.1, Theorem 4.2]). *Given an \mathcal{L} -structure M , let \mathcal{L}_M be \mathcal{L} expanded by constant symbols for every element of M . The following are equivalent.*

- (1) M is mutually algebraic.
- (2) Every atomic \mathcal{L} -formula is equivalent to a boolean combination of quantifier-free mutually algebraic \mathcal{L}_M -formulas.
- (3) Every \mathcal{L}_M -formula is equivalent to a boolean combination of \mathcal{L}_M -formulas of the form $\exists \bar{y} \theta(\bar{x}, \bar{y})$, where $\theta(\bar{x}, \bar{y})$ is quantifier-free mutually algebraic.
- (4) M is weakly minimal and trivial.

While Theorem 2.3 is useful, the dependence on \mathcal{L}_M -formulas can be awkward in applications. To alleviate this, we introduce a stronger notion of an MA-presented structure and with Proposition 2.6, we see that such a structure admits a better reduction of quantifier complexity. However, as we will see in Lemma 2.22, every mutually algebraic structure M is *associated* (see Definition 2.18) to an MA-presented structure M' in a new language.

2.1. MA-presented structures.

Definition 2.4. An \mathcal{L} -structure M is *MA-presented* if every atomic \mathcal{L} -formula is mutually algebraic. M is *finitely MA-presented* if, in addition, \mathcal{L} is finite relational.

We begin by bounding the quantifier complexity of formulas in an MA-presented structure.

Definition 2.5. Let $\mathcal{E} = \{\exists \bar{z} \theta(\bar{y}, \bar{z})\}$, where $\theta(\bar{y}, \bar{z})$ is quantifier-free and mutually algebraic \mathcal{L} -formula. We allow $\text{lg}(\bar{z}) = 0$, so as M is MA-presented, every atomic \mathcal{L} -formula is in \mathcal{E} .

Let $\mathcal{E}^* := \{\text{all } \mathcal{L}\text{-formulas equivalent to boolean combinations and adjunctions of free variables of formulas from } \mathcal{E}\}$. As every $\gamma(\bar{x}) \in \mathcal{E}$ is mutually algebraic, the set \mathcal{E} of formulas is **not** closed under adjunction of dummy variables, although \mathcal{E}^* is.

We note a dichotomy among partitioned formulas $\beta(x, \bar{y}) := \exists \bar{z} \theta(x, \bar{y}, \bar{z}) \in \mathcal{E}$ with $\text{lg}(x) = 1$: On one hand, if both $\text{lg}(\bar{y}) = \text{lg}(\bar{z}) = 0$, i.e., if $\beta(x) := \theta(x)$

is quantifier-free and mutually algebraic, then $\neg\beta(x)$ is also quantifier-free, mutually algebraic (recall that every formula in one free variable is mutually algebraic) hence $\neg\beta \in \mathcal{E}$. On the other hand, if at least one of $\text{lg}(\bar{y}), \text{lg}(\bar{z}) > 0$, then for some $K \in \omega$, $M \models \forall \bar{y} \exists^{\leq K} x \beta(x, \bar{y})$, i.e., $\neg\beta(M, \bar{c})$ is cofinite for all $\bar{c} \in M^{\text{lg}(\bar{y})}$.

Proposition 2.6. *Suppose M is an infinite, MA-presented \mathcal{L} -structure. Then every \mathcal{L} -formula is equivalent to a formula in \mathcal{E}^* .*

Proof. This proof is similar to the proof of Theorem 4.2 of [?MAO]. Via Disjunctive Normal Form and the dichotomy above, every formula in \mathcal{E}^* in which the variable x appears in each element of \mathcal{E} can be written as

$$\bigvee_k \left[\bigwedge_i \alpha_{i,k}(x, \bar{y}_{i,k}) \wedge \bigwedge_j \neg\beta_{j,k}(x, \bar{y}_{j,k}) \right]$$

where $\neg\beta_{j,k}(M, \bar{c})$ is cofinite for all j, k , and \bar{c} .

Note that \mathcal{E}^* contains all atomic \mathcal{L} -formulas and is closed under boolean combinations. Thus, to prove the Proposition, it suffices to show that \mathcal{E}^* is closed under projections. By the standard form given above and the fact that \exists commutes with \bigvee , it suffices to show that

$$\delta(\bar{y}) := \exists x \left(\bigwedge_i \alpha_i(x, \bar{y}_i) \wedge \bigwedge_{j=1}^J \neg\beta_j(x, \bar{y}_j) \right) \in \mathcal{E}^*$$

where each $\alpha_i, \beta_j \in \mathcal{E}$ and where $\neg\beta_j(M, \bar{c}_j)$ is cofinite for all j and \bar{c}_j from M . We split into cases.

Case 1. There is no α_i .

In this case, as M is infinite and each $\neg\beta_j(M, \bar{c}_j)$ is cofinite, $\delta(\bar{y}) := \exists x \bigwedge_j \neg\beta_j(x, \bar{y}_j)$ holds for all \bar{y} .

Case 2. There is at least one α_i .

As the conjunction of elements of \mathcal{E} with a common variable symbol is also in \mathcal{E} , we may assume there is exactly one $\alpha^*(x, \bar{y}^*)$. There are now two subcases.

Subcase 2a). $\bar{y}^* = \emptyset$ and $\alpha^*(M)$ is infinite.

Then as in Case 1, $\delta(\bar{y}) := \exists x [\alpha^*(x) \wedge \bigwedge_j \neg\beta_j(x, \bar{y}_j)]$ holds for all \bar{y} .

Subcase 2b). Not Subcase 2a).

In this case, for some $r^* \in \omega$, $M \models \forall \bar{y}^* \exists^{\leq r^*} x \alpha^*(x, \bar{y}^*)$. For each j , put $\gamma_j(x, \bar{y}_j) := \alpha^*(x, \bar{y}^*) \wedge \beta_j(x, \bar{y}_j)$. Thus, $\gamma_j(M, \bar{c}_j) \subseteq \alpha^*(M, \bar{c}^*)$ for all j and all $\bar{c} \in M^{\text{lg}(\bar{y})}$. Also, as α^* and β_j contain a common free variable, namely x ,

each $\gamma_j(x, \bar{y}'_j)$ is mutually algebraic. It follows that $\delta(\bar{y})$ is equivalent to

$$\bigvee_{1 \leq r \leq r^*} \left(\exists^{=r} x \alpha^*(x, \bar{y}^*) \wedge \exists^{<r} x \bigvee_{j=1}^J \gamma_j(x, \bar{y}'_j) \right)$$

so it suffices to show that this formula is in \mathcal{E}^* .

As in the proof of Lemma 2.4 of [8], for each r and for each $S \subseteq \{1, \dots, J\}$, $\exists^{\leq r} x \bigwedge_{j \in S} \gamma_j(x, \bar{y}'_j) \in \mathcal{E}$. From this, it follows easily that $\exists^{=r} x \bigwedge_{j \in S} \gamma_j(x, \bar{y}'_j)$ and $\exists^{<r} x \bigwedge_{j \in S} \gamma_j(x, \bar{y}'_j)$ are in \mathcal{E}^* . Finally, by the inclusion-exclusion principle, it follows that $\bigvee_{j=1}^J \gamma_j(x, \bar{y}'_j)$ is also in \mathcal{E}^* and we finish. \square

This restriction on quantifier complexity will be used in the analysis of the following natural decomposition of an MA-presented structure into independent components, analogous to the decomposition of a graph into connected components.

Definition 2.7. Suppose M is an MA-presented \mathcal{L} -structure. For $a, b \in M$, call b a *mate* of a if there is some \bar{c} from M and the type of (some permutation of) $a\bar{c}$ contains an atomic \mathcal{L} -formula. Being a mate is reflexive and symmetric. Let \sim be the transitive closure of the mate relation. As \sim is an equivalence relation, write $[a]_{\sim}$ for the \sim -class of $a \in M$, and call \sim -classes *MA-connected components*.

Observe that if M is an MA-presented \mathcal{L} -structure, then $M \models \neg R(\bar{d})$ for all \mathcal{L} -atomic $R(\bar{x})$ and all $\bar{d} \in M^{\text{lg}(\bar{x})}$ whenever \bar{d} is not contained in a single \sim -class.

Definition 2.8. Suppose M is an MA-presented \mathcal{L} -structure. A *component map* is a bijection $f : M \rightarrow M$ such that $f([a]_{\sim}) = [f(a)]_{\sim}$ (setwise) for all $a \in M$. An *\mathcal{L} -component map* is a component map such that for each $a \in M$, $f \upharpoonright_{[a]_{\sim}} : [a]_{\sim} \rightarrow [f(a)]_{\sim}$ is an \mathcal{L} -isomorphism.

The following Lemma, similar to [?MA, Proposition 4.4] although working with a single structure, follows immediately from the observation above.

Lemma 2.9. *If M is an MA-presented \mathcal{L} -structure, then $f : M \rightarrow M$ is an automorphism if and only if it is an \mathcal{L} -component map.*

The following definition gives a syntactic characterization of the \sim -relation.

Definition 2.10. A *linked \mathcal{L} -atomic conjunction* is an \mathcal{L} -formula of the form $\theta(\bar{w}) := \bigwedge_{i < n} \alpha_i(\bar{w}_i)$, where each $\alpha_i(\bar{w}_i)$ is \mathcal{L} -atomic and $\bar{w}_i \cap \bar{w}_{i+1} \neq \emptyset$ for all $i < n - 1$. (If $n = 1$, this last condition is vacuous, hence every \mathcal{L} -atomic formula is a linked \mathcal{L} -atomic conjunction.)

Note that if every atomic \mathcal{L} -formula is mutually algebraic (e.g., when M is MA-presented), then as $\phi(\bar{u}) \wedge \psi(\bar{v})$ is mutually algebraic whenever $\phi(\bar{u}), \psi(\bar{v})$ are mutually algebraic and $\bar{u} \cap \bar{v} \neq \emptyset$, it follows every linked \mathcal{L} -atomic conjunction $\theta(\bar{w})$ is mutually algebraic as well.

Lemma 2.11. *Suppose M is an infinite, MA-presented \mathcal{L} -structure. For $a, b \in M$, $a \sim b$ if and only if $\text{tp}(a, b)$ contains $\exists \bar{u} \theta(x, y, \bar{u})$ for some linked, \mathcal{L} -atomic conjunction $\theta(x, y, \bar{u})$.*

Proof. First, assume $a \sim b$. Choose $\{c_0, \dots, c_n\} \subseteq M$ such that $c_0 = a$, $c_n = b$, and c_{i+1} is a mate of c_i for each $i < n$. For each $i < n$, choose an \mathcal{L} -atomic $\alpha_i(x_i, x_{i+1}, \bar{z}_i)$ witnessing c_{i+1} is a mate of c_i . Without loss, assume \bar{z}_i and \bar{z}_j are disjoint for distinct i, j . Then, writing \bar{w}_i for $x_i x_{i+1} \bar{z}_i$, $\theta(\bar{w}) := \bigwedge_{i < n} \alpha_i(\bar{w}_i)$ is a linked \mathcal{L} -atomic conjunction and $M \models \exists \bar{u} \theta(a, b, \bar{u})$, where $\bar{u} := \bar{w} \setminus \{x_0, x_n\}$.

Conversely, assume $M \models \exists \bar{u} \theta(a, b, \bar{u})$, where $\theta(\bar{w}) := \bigwedge_{i < n} \alpha_i(\bar{w}_i)$, with each α_i \mathcal{L} -atomic and $\bar{w}_i \cap \bar{w}_{i+1} \neq \emptyset$. For each $i < n$, choose a variable symbol $x_i \in \bar{w}_i \cap \bar{w}_{i+1}$. Choose $\bar{c} \in M^{\text{lg}(\bar{u})}$ such that $M \models \theta(a, b, \bar{c})$, let \bar{c}_i denote the restriction of \bar{c} corresponding to \bar{w}_i , and let d_i be the singleton in \bar{c}_i corresponding to x_i . Without loss, assume $a \in \bar{c}_0$ and $b \in \bar{c}_{n-1}$. Then a, d_0 and b, d_{n-1} are mates, as are d_i, d_{i+1} for each $i < n-1$. Thus, $a \sim b$. \square

For elements $a, b \in \text{acl}(\emptyset)$, $a \sim b$ can be rather strange (e.g., if c, d are constant symbols then the formula $\phi(x, y) := (x = c \wedge y = d)$ is mutually algebraic). However, outside of $\text{acl}(\emptyset)$, the \sim relation characterizes dependence. The following definition appears in [?LT1].

Definition 2.12. Let M be any \mathcal{L} -structure. A mutually algebraic \mathcal{L} -formula $\theta(\bar{z})$ supports an infinite array if there is an infinite set $\{\bar{d}_i : i \in \omega\}$ such that $M \models \theta(\bar{d}_i)$ for each i and $\bar{d}_i \cap \bar{d}_j = \emptyset$ for all $i \neq j$.

Lemma 2.13. *Suppose M is any structure and $\theta(\bar{z})$ is a mutually algebraic formula. Then $\theta(\bar{z})$ supports an infinite array if and only if $\theta(\bar{z})$ is not algebraic, i.e., $M \models \exists^\infty \bar{z} \theta(\bar{z})$.*

Proof. Left to right is obvious, so assume $\theta(\bar{z})$ is not algebraic. Choose any infinite set $\{\bar{d}_i : i \in \omega\}$ of distinct realizations of $\theta(\bar{z})$. By the finite Δ -system lemma, there is a finite set R and an infinite $I \subseteq \omega$ such that $\bar{d}_i \cap \bar{d}_j = R$ for all distinct $i, j \in I$. As I is infinite and $\theta(\bar{z})$ is mutually algebraic, we must have $R = \emptyset$, hence $\theta(\bar{z})$ supports an infinite array. \square

Lemma 2.14. *Suppose M is MA-presented, $\bar{c} \in M^k$, and $\bar{c} \cap \text{acl}(\emptyset) = \emptyset$.*

- (1) *If $\text{tp}(\bar{c})$ contains a quantifier-free, mutually algebraic \mathcal{L} -formula $\theta(\bar{z})$ with $\text{lg}(\bar{z}) = \text{lg}(\bar{c})$, then $\text{tp}(\bar{c})$ contains a linked \mathcal{L} -atomic conjunction $\rho(\bar{z})$.*
- (2) *If $\text{tp}(\bar{c})$ contains a mutually algebraic \mathcal{L} -formula $\delta(\bar{x})$ with $\text{lg}(\bar{x}) = \text{lg}(\bar{c})$, then there are $\bar{z} \supseteq \bar{x}$, $\bar{d} \supseteq \bar{c}$ with $\text{lg}(\bar{d}) = \text{lg}(\bar{z})$, and a linked \mathcal{L} -atomic conjunction $\rho(\bar{z}) \in \text{tp}(\bar{d})$.*
- (3) *If \bar{c} satisfies a quantifier-free mutually algebraic \mathcal{L}_M -formula $\theta(\bar{x}, \bar{m})$ with $\bar{c} \cap \text{acl}(\bar{m}) = \emptyset$. Then there are $\bar{z} \supseteq \bar{x}$, $\bar{d} \supseteq \bar{c}$ with $\text{lg}(\bar{d}) = \text{lg}(\bar{z})$, and a linked \mathcal{L} -atomic conjunction $\delta(\bar{z}) \in \text{tp}(\bar{d})$.*

Proof. By passing to a large enough elementary extension, we may assume M is $|\mathcal{L}|^+$ -saturated. For (1), by writing $\theta(\bar{z})$ in Disjunctive Normal Form

and choosing a disjunct in $\text{tp}(\bar{c})$, we may assume $\theta(\bar{z})$ has the form

$$\bigwedge_{i \in S} \alpha_i(\bar{z}_i) \wedge \bigwedge_{j \in U} \neg \beta_j(\bar{z}_j)$$

where each α_i, β_j are \mathcal{L} -atomic. Now, since M is MA-presented, each $\alpha_i(\bar{z}_i), \beta_j(\bar{z}_j)$ is mutually algebraic, and since $\bar{c} \cap \text{acl}(\emptyset) = \emptyset$, the saturation of M implies that $\theta(\bar{z})$ supports an infinite array. Thus, by Lemma A.1 of [?LT1], there is an $S_0 \subseteq S$ such that $\rho(\bar{z}) := \bigwedge_{i \in S_0} \alpha_i(\bar{z}_i)$ is a linked, \mathcal{L} -atomic conjunction with $\bigcup\{\bar{z}_i : i \in S_0\} = \bar{z}$. As $\theta(\bar{z}) \vdash \rho(\bar{z})$, $\rho(\bar{z}) \in \text{tp}(\bar{c})$.

For (2), by Proposition 2.6, we may assume $\delta(\bar{x})$ has the form there. By choosing a disjunct satisfied by \bar{c} , we may assume $\delta(\bar{x})$ has the form

$$\bigwedge_{i \in S} \phi_i(\bar{x}_i) \wedge \bigwedge_{j \in U} \neg \psi_j(\bar{x}_j)$$

where each $\phi_i(\bar{x}_i), \psi_j(\bar{x}_j)$ is from \mathcal{E} , and so is mutually algebraic. As $\delta(\bar{x})$ supports an infinite array, again we have a subset $S_0 \subseteq S$ with $\gamma(\bar{x}) := \bigwedge_{i \in S_0} \phi_i(\bar{x}_i)$, and $\{\bar{x}_i : i \in S_0\}$ linked with $\bigcup\{\bar{x}_i : i \in S_0\} = \bar{x}$. Now, each $\phi_i(\bar{x}_i)$ has the form $\exists \bar{w}_i \eta_i(\bar{x}_i, \bar{w}_i)$ where $\eta_i(\bar{x}_i, \bar{w}_i)$ is quantifier-free and mutually algebraic. Thus, $\gamma(\bar{x})$ is equivalent to a formula $\exists \bar{w}^* \theta(\bar{x}, \bar{w}^*)$ with $\theta(\bar{x}, \bar{w}^*)$ quantifier free and mutually algebraic since $\{\eta_i(\bar{x}_i, \bar{w}_i) : i \in S_0\}$ are linked. Additionally, $M \models \gamma(\bar{c})$. Write θ as $\theta(\bar{z})$ with $\bar{z} \supseteq \bar{x}$. Choose $\bar{d} \supseteq \bar{c}$ such that $M \models \theta(\bar{d})$. As $\theta(\bar{z})$ is mutually algebraic, $\bar{d} \cap \text{acl}(\emptyset) = \emptyset$. Thus, we can apply (1) to $\theta(\bar{z})$ and \bar{d} , getting $\rho(\bar{z})$ as required.

(3) Let $M_{\bar{m}}$ be the expansion of M naming the constants of \bar{m} . Applying (1) to $\theta(\bar{x}, \bar{m})$ in $M_{\bar{m}}$, $\text{tp}(\bar{c})$ contains a linked, $\mathcal{L}_{\bar{m}}$ -atomic conjunction $\rho(\bar{x}, \bar{m}) := \bigwedge_{i \in S_0} \alpha_i(\bar{x}_i, \bar{m}_i)$ with $\bigcup\{\bar{x}_i : i \in S_0\} = \bar{x}$. As M is MA-presented, every atomic formula $\alpha_i(\bar{x}, \bar{w}_i)$ is mutually algebraic, hence by choosing disjoint variables for each \bar{w}_i and letting $\bar{w}^\# = \langle \bar{w}_i : i \in S_0 \rangle$, we have $M \models \exists \bar{w}^\# \delta(\bar{c}, \bar{w}^\#)$ where $\delta(\bar{c}, \bar{w}^\#) := \bigwedge_{i \in S_0} \alpha_i(\bar{x}_i, \bar{w}_i)$. Choose any $\bar{d} \supseteq \bar{c}$ with $M \models \delta(\bar{d})$ and we finish, taking $\bar{z} := \bar{x} \hat{\ } \bar{w}^\#$. \square

Proposition 2.15. *Suppose M is an MA-presented \mathcal{L} -structure and let $a \in M \setminus \text{acl}(\emptyset)$. Then $[a]_{\sim} = \text{acl}(a) \setminus \text{acl}(\emptyset)$.*

Proof. Assume M is MA-presented and fix $a \in M \setminus \text{acl}(\emptyset)$. First, assume $b \sim a$. By Lemma 2.11 choose a linked \mathcal{L} -atomic conjunction $\theta(x, y, \bar{u})$ and \bar{c} such that $M \models \theta(a, b, \bar{c})$. As θ is mutually algebraic, it follows that $b \in \text{acl}(a)$ and $a \in \text{acl}(b)$. Thus, if $b \in \text{acl}(\emptyset)$, we would have $a \in \text{acl}(\emptyset)$ as well, which it is not. So $b \in \text{acl}(a) \setminus \text{acl}(\emptyset)$.

Conversely, assume $b \in \text{acl}(a) \setminus \text{acl}(\emptyset)$. As M mutually algebraic implies $\text{Th}(M)$ is weakly minimal and so acl satisfies exchange, $a \in \text{acl}(b)$. Thus, there is an integer K and \mathcal{L} -formulas $\phi(x, y), \psi(x, y) \in \text{tp}(a, b)$ such that $M \models \exists^{\leq K} x \phi(x, b)$ and $M \models \exists^{\leq K} y \psi(a, y)$. Let

$$\delta(x, y) := \phi(x, y) \wedge \psi(x, y) \wedge \exists^{\leq K} u \phi(u, y) \wedge \exists^{\leq K} v \psi(x, v)$$

Then $\delta(x, y) \in \text{tp}(a, b)$ and is mutually algebraic. Additionally, $ab \cap \text{acl}(\emptyset) = \emptyset$, so $b \sim a$ follows from Lemma 2.14(2) and Lemma 2.11. \square

2.2. Associating a mutually algebraic structure with an MA-presented structure.

In this section, we define what we mean by an *association* of structures in different languages and prove that every mutually algebraic structure is associated to an MA-presented structure. This procedure is far from unique, but we explore how different the \sim -relations can be.

Definition 2.16. Suppose $\mathcal{L} \subseteq \mathcal{L}^*$ are languages and M is an \mathcal{L} -structure. An \mathcal{L}^* -structure M^* is a q.f. expansion by definitions of M if M^* is an expansion of M (i.e., M^* and M have the same universes and the interpretations of every \mathcal{L}_M -formula are the same in both structures) such that every atomic \mathcal{L}^* -formula is equivalent to a quantifier-free \mathcal{L}_M -formula.

The following facts are easy.

Fact 2.17. *Suppose $\mathcal{L} \subseteq \mathcal{L}^*$ and M^* is a q.f. expansion of M by definitions. Then:*

- (1) *Every q.f. \mathcal{L}_M^* -formula is equivalent to a q.f. \mathcal{L}_M -formula.*
- (2) *Every \mathcal{L}_M^* -formula $\theta^*(\bar{x})$ is equivalent to an \mathcal{L}_M -formula $\gamma(\bar{x})$.*
- (3) *For every $N \succeq M$ there is a unique $N^* \succeq M^*$ such that N^* is a q.f. expansion of N by definitions; and*
- (4) *For every $N^* \succeq M^*$, the \mathcal{L} -reduct $N \succeq M$ and N^* is a q.f. expansion by definitions of N .*
- (5) *If \mathcal{L}^* (and hence \mathcal{L}) are finite, relational, then there is some finite \bar{m} from M so that every \mathcal{L}^* -formula $\phi(\bar{x})$ is equivalent in M^* to some $\mathcal{L}_{\bar{m}}$ -formula $\psi(\bar{x}, \bar{m})$.*

Proof. (1) By taking boolean combinations, one immediately gets that every q.f. \mathcal{L}^* -formula is equivalent to a q.f. \mathcal{L}_M -formula. Then (1) follows from this by taking specializations.

(2) As M and M^* have the same universes, one proves by induction on the complexity of \mathcal{L}^* -formulas that every \mathcal{L}^* -formula is equivalent to an \mathcal{L}_M -formula. By taking specializations, (2) follows.

(3) Given any atomic \mathcal{L}^* -formula $\alpha(\bar{x})$, let $\theta(\bar{x})$ be a q.f. \mathcal{L} -formula equivalent to it. Now interpret $\alpha(N^*)$ as $\theta(N)$. This works, and is the unique way of obtaining such an N^* .

(4) Easy.

(5) As \mathcal{L}^* is finite relational, let $\{\alpha_i(\bar{x}_i) : i < K\}$ enumerate the (finitely many) \mathcal{L}^* -atomic formulas. By (1), for each i , choose an \mathcal{L} -formula $\psi_i(\bar{x}_i, \bar{y}_i)$ and $\bar{m}_i \in M^{\text{lg}(\bar{y}_i)}$ such that $\psi_i(\bar{x}_i, \bar{m}_i)$ equivalent on M^* to $\alpha_i(\bar{x}_i)$. Let \bar{m} be any finite tuple containing every \bar{m}_i . It follows by induction on the complexity of \mathcal{L}^* -formulas that every \mathcal{L}^* -formula is equivalent on M^* to an $\mathcal{L}_{\bar{m}}$ -formula. \square

Definition 2.18. Suppose \mathcal{L}_1 and \mathcal{L}_2 are disjoint languages. An \mathcal{L}_1 -structure M_1 and an \mathcal{L}_2 -structure M_2 are *associated* if there is an $\mathcal{L}_1 \cup \mathcal{L}_2$ -structure M^* that is a q.f. expansion by definitions of both M_1 and M_2 .

The most obvious example of associated structures is that any expansion M' of a structure M by naming constants is associated to M . The following results are immediate from Fact 2.17.

Fact 2.19. *Suppose M_1, M_2 are associated structures in languages $\mathcal{L}_1, \mathcal{L}_2$, respectively. Then:*

- (1) M_1 and M_2 have the same universe, which we denote by M ;
- (2) Every (q.f.) $\mathcal{L}_1(M)$ -formula is equivalent to a (q.f.) $\mathcal{L}_2(M)$ -formula; and
- (3) For $\ell = 1, 2$, if $N_\ell \succeq M_\ell$ then there is a unique $N_{3-\ell} \succeq M_{3-\ell}$ with universe N_ℓ such that N_1 is associated to N_2 .

Definition 2.20. Suppose M is a mutually algebraic \mathcal{L} -structure. A set \mathcal{A} of quantifier-free, mutually algebraic \mathcal{L}_M -formulas is *acceptable* if every $R \in \mathcal{L}$ is equivalent to a boolean combination of formulas from \mathcal{A} . If \mathcal{A} is an acceptable set of formulas, let $D(\mathcal{A}) := \bigcup \{\bar{a} : \phi(\bar{x}, \bar{a}) \in \mathcal{A}\}$ denote the parameters from M used in \mathcal{A} .

Note that by Theorem 2.3(2), every M admits an acceptable set \mathcal{A} . Moreover, if \mathcal{L} is finite, then \mathcal{A} and hence $D(\mathcal{A})$ can be chosen to be finite as well.

Definition 2.21. Suppose M is a mutually algebraic \mathcal{L} -structure and \mathcal{A} is an acceptable set of \mathcal{L}_M -formulas. Let

$$\mathcal{L}_{\mathcal{A}} := \{\text{constant symbols of } \mathcal{L}\} \cup \{c_d : d \in D(\mathcal{A})\} \cup \{R_\phi(\bar{x}) : \phi(\bar{x}, \bar{a}) \in \mathcal{A}\}$$

and let $M_{\mathcal{A}}$ be the $\mathcal{L}_{\mathcal{A}}$ -structure with universe M , $c^{M_{\mathcal{A}}} = c^M$ for every $c \in \mathcal{L}$, $c_d^{M_{\mathcal{A}}} = d$ for every $d \in D(\mathcal{A})$, and $R_\phi^{M_{\mathcal{A}}} := \{\bar{b} \in M^{\text{lg}(\bar{x})} : M \models \phi(\bar{b}, \bar{a})\}$.

Lemma 2.22. *Let M be any mutually algebraic \mathcal{L} -structure, \mathcal{A} an acceptable set of \mathcal{L}_M -formulas, and let $M_{\mathcal{A}}$ be the $\mathcal{L}_{\mathcal{A}}$ -structure constructed as above. Then:*

- (1) $M_{\mathcal{A}}$ is MA -presented.
- (2) M and $M_{\mathcal{A}}$ are associated.
- (3) For every $Y \subseteq M$ and every subset $X \subseteq M^k$,
 - (a) X is Y -definable in $M_{\mathcal{A}}$ if and only if X is $(D(\mathcal{A}) \cup Y)$ -definable in M .
 - (b) $\text{acl}_{M_{\mathcal{A}}}(Y) = \text{acl}_M(D(\mathcal{A}) \cup Y)$.

Proof. The verifications of (1) and (2) are just unpacking definitions. It is also clear that (3b) follows immediately from (3a), which itself follows from induction on the complexity of formulas, using that the constants of \mathcal{L} are present in $\mathcal{L}_{\mathcal{A}}$ and interpreted correctly to handle the case of atomic formulas. \square

Now, $\text{acl}_{M_{\mathcal{A}}}(\emptyset)$ can vary widely among associated MA-presentations of M (e.g., the set of ‘extra parameters’ $D(\mathcal{A})$ can be increased arbitrarily), but we will see that \sim -classes sufficiently far away from $\text{acl}_{M_{\mathcal{A}}}(\emptyset)$ are invariant under our choice of MA-presentation. We begin with the following lemma.

Lemma 2.23. *Suppose M_1 and M_2 are both associated MA-presented \mathcal{L}_1 , \mathcal{L}_2 -structures respectively with common mutual expansion M^* , and suppose $a \notin \text{acl}_{M^*}(\emptyset)$. Then $[a]_{\sim_{M_1}} = [a]_{\sim_{M_2}}$.*

Proof. By passing to sufficiently saturated elementary extensions, we may assume that M_1 , M_2 and M^* are all $(|\mathcal{L}_1| + |\mathcal{L}_2|)^+$ -saturated. By symmetry, it suffices to prove that if $a \sim_{M_1} b$, then $a \sim_{M_2} b$. For this, choose a finite tuple \bar{c} with $a, b \in \bar{c}$ and a quantifier-free, mutually algebraic \mathcal{L}_1 -formula $\phi(\bar{x}) \in \text{tp}(\bar{c})$. Since $a \notin \text{acl}_{M^*}(\emptyset)$, the mutual algebraicity of $\phi(\bar{x})$ implies that $\bar{c} \cap \text{acl}_{M^*}(\emptyset) = \emptyset$.

As M_2 is associated to M_1 , choose a quantifier-free $\mathcal{L}_2(M)$ -formula $\theta(\bar{x}, \bar{m})$ equivalent to $\phi(\bar{x})$ on M . Thus,

$$M^* \models \forall \bar{x}(\phi(\bar{x}) \leftrightarrow \theta(\bar{x}, \bar{m}))$$

As M^* is sufficiently saturated, choose \bar{m}' from M^* so that $\text{stp}_{M^*}(\bar{m}') = \text{stp}_{M^*}(\bar{m})$ and \bar{m}' independent from \bar{c} over \emptyset (in M^*). As $\text{tp}_{M^*}(\bar{c}, \bar{m}') = \text{tp}_{M^*}(\bar{c}, \bar{m})$, we have $M^* \models \theta(\bar{c}, \bar{m}')$, and by the independence we have $\bar{c} \cap \text{acl}_{M^*}(\bar{m}') = \emptyset$. But now, since $\theta(\bar{x}, \bar{m}')$ is a quantifier-free $\mathcal{L}_2(M)$ -formula, $M_2 \models \theta(\bar{c}, \bar{m}')$ and since M_2 is a reduct of M^* , $\bar{c} \cap \text{acl}_{M_2}(\bar{m}') = \emptyset$. Thus, by Lemma 2.14(3) there are $\bar{z} \supseteq \bar{x}$, $\bar{d} \supseteq \bar{c}$ with $\text{lg}(\bar{d}) = \text{lg}(\bar{z})$, and a linked \mathcal{L}_2 -atomic conjunction $\delta(\bar{z}) \in \text{tp}(\bar{d})$. As $a, b \in \bar{d}$, \bar{d} witnesses that $a \sim_{M_2} b$. \square

It is tempting to weaken the hypothesis of Lemma 2.23 to simply ‘ $a \notin \text{acl}_{M_1}(\emptyset) \cup \text{acl}_{M_2}(\emptyset)$ ’ but the following example shows this is not sufficient.

Example 1. Let \mathcal{L}_1 contain a single binary relation R_1 and let M_1 be an \mathcal{L}_1 -structure containing infinitely many directed \mathbb{Z} -chains, \mathbb{Z}^+ -chains, and \mathbb{Z}^- -chains. In some \mathbb{Z} -chain of M_1 , pick a point a and its successor b , although we do not name them by constants. Let \mathcal{L}_2 contain a single binary relation R_2 , and let M_2 be the \mathcal{L}_2 -structure with the same universe as M_1 and $R_2^{M_2}$ interpreted as $R_1^{M_1} \setminus \{(a, b)\}$. It is easily checked that M_1 and M_2 are associated, that $\text{acl}_{M_1}(\emptyset) = \text{acl}_{M_2}(\emptyset) = \emptyset$, and that the \sim -components are the chains.

Now M_2 looks the same as M_1 , except that the \mathbb{Z} -chain containing m and m' has been broken to produce a \mathbb{Z}^+ -chain with ‘minimum’ element b and a \mathbb{Z}^- -chain with ‘maximum’ element a . Hence $a \sim_{M_1} b$, but $a \not\sim_{M_2} b$.

Despite this example, if $M_{\mathcal{A}_1}$ and $M_{\mathcal{A}_2}$ are each associated to the same mutually algebraic M as in Definition 2.21, we can do better.

Proposition 2.24. *Suppose M is a mutually algebraic \mathcal{L} -structure with two acceptable sets \mathcal{A}_1 and \mathcal{A}_2 of \mathcal{L}_M -formulas. For $\ell = 1, 2$, let acl_ℓ denote the*

algebraic closure relation on $M_{\mathcal{A}_\ell}$ and let $[a]_\ell$ denote the \sim -component of a in $M_{\mathcal{A}_\ell}$. For any $a \in M \setminus (\text{acl}_1(\emptyset) \cup \text{acl}_2(\emptyset))$ we have $[a]_1 = [a]_2$.

Proof. To ease notation, for $\ell = 1, 2$, write M_ℓ in place of $M_{\mathcal{A}_\ell}$ and D_ℓ in place of $D(\mathcal{A}_\ell)$. As M_1 and M_2 are associated, let M^* be their common q.f. expansion. In light of Lemma 2.23, it suffices to show that $a \notin \text{acl}_{M^*}(\emptyset)$.

As M, M_1, M_2, M^* all have the same universe and the same interpretation of every constant $c \in \mathcal{L}$, it follows by induction on formulas as in Lemma 2.22 that every $X \subseteq M^k$ that is 0-definable in M^* is $(D_1 \cup D_2)$ -definable in M . It then follows that $\text{acl}_{M^*}(\emptyset) \subseteq \text{acl}_M(D_1 \cup D_2)$. However, as $a \notin \text{acl}_\ell(\emptyset)$ for both $\ell = 1, 2$, it follows from Lemma 2.22 that $a \notin (\text{acl}_M(D_1) \cup \text{acl}_M(D_2))$. But, as M is mutually algebraic, algebraic closure is trivial, hence $a \notin \text{acl}_M(D_1 \cup D_2)$, hence not in $\text{acl}_{M^*}(\emptyset)$. Thus, $[a]_1 = [a]_2$ follows from Lemma 2.23. \square

We close with one additional instance where the hypotheses of Lemma 2.23 can be weakened.

Corollary 2.25. *Suppose M is mutually algebraic and $N \succeq M$. Then the partition of $(N \setminus M)$ given by \sim does not depend on the choice of an associated MA-presented structure, i.e., if N_1, N_2 are both MA-presented and associated to N , then $[a]_{\sim_{N_1}} = [a]_{\sim_{N_2}}$ for every $a \in (N \setminus M)$.*

Proof. Choose any MA-presented N_1, N_2 associated to N . As N_1 and N_2 are associated to each other, let N^* be their common mutual expansion. As $a \in N \setminus M$, we have $a \notin \text{acl}_{N^*}(\emptyset)$. Thus, $[a]_{\sim_{N_1}} = [a]_{\sim_{N_2}}$ by Lemma 2.23. \square

3. CELLULARITY

We begin by stating our definition of cellularity.

Definition 3.1. A structure M is *cellular* if, for some integers n and $\langle k_i : i \in [n] \rangle$, it admits a partition $\{K, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$ satisfying the following.

- (1) K is finite and $\text{lg}(\bar{c}_{i,j}) = k_i$.
- (2) For a fixed i , $\{\bar{c}_{i,j} \mid j \in \omega\}$ are pairwise isomorphic over K . We may thus enumerate each as $\bar{c}_{i,j} = (c_{i,j}^1, \dots, c_{i,j}^{k_i})$.
- (3) For every $i \in [n]$ and $\sigma \in S_\infty$, there is a $\sigma_i^* \in \text{Aut}(M)$ mapping each $\bar{c}_{i,j}$ onto $\bar{c}_{i,\sigma(j)}$ by sending $c_{i,j}^\ell$ to $c_{i,\sigma(j)}^\ell$ for $1 \leq \ell \leq k_i$, and fixing $M \setminus \bigcup_{j \in \omega} \bar{c}_{i,j}$ pointwise.

The original definition of cellularity from [Schmerl] and its rephrasing in [MPW] require $n = 1$. But given a partition as in Definition 3.1, we may produce one with $n = 1$ by taking \bar{c}_j to be the concatenation of $\bar{c}_{i,j}$ for each $i \in [n]$, so our definition is equivalent.

Note that like ω -categoricity, ω -stability, and mutual algebraicity, the cellularity of a structure M is preserved under reducts. Thus, we begin with the following Lemma that does not depend on our choice of language.

Lemma 3.2. *If M is cellular, then M is ω -categorical, ω -stable, and mutually algebraic.*

Proof. Call a cellular partition $\{K, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$ of M *indecomposable* if, for each $i \in [n]$ if we partition each $\bar{c}_{i,j}$ as $\bar{d}_{i,j} \hat{\ } \bar{e}_{i,j}$ with both sides non-empty, then

$$\{K, \{\bar{c}_{i',j} \mid i' \in [n], i' \neq i, j \in \omega\} \cup \{\bar{d}_{i,j} : j \in \omega\} \cup \{\bar{e}_{i,j} : j \in \omega\}\}$$

is *not* a cellular decomposition of M .

As $\text{lg}(\bar{c}_{i,j})$ drops each time a cellular partition is decomposed, it follows that every cellular M has an indecomposable cellular partition. Fix an indecomposable cellular partition $\{K, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$ witnessing the cellularity of M . Form a new finite relational language

$$\mathcal{L}^* := \{c_k : k \in K\} \cup \{U_{i,\ell} : i \in [n], 1 \leq \ell \leq k_i\} \cup \{R_i(\bar{x}) : i \in [n]\}$$

where each $U_{i,\ell}$ is unary and each R_i is a k_i -ary relation symbol. Let M^* be the \mathcal{L}^* -structure with universe M , and interpretations $c_k^{M^*} = k$ for each $k \in K$, $U_{i,\ell}^{M^*} := \{c_{i,j}^\ell : j \in \omega\}$, and $R_i^{M^*} := \{\bar{c}_{i,j} : j \in \omega\}$. It is easily verified that the ‘grid-like structure’ M^* is ω -categorical, ω -stable, and mutually algebraic. Moreover, by the automorphism condition in the definition of cellularity, it follows that every \mathcal{L} -definable set in M is \mathcal{L}^* -definable in M^* , hence M is a reduct M^* . \square

The following Lemma crucially requires \mathcal{L}^* to be finite relational. As an example, naming infinitely many constants destroys ω -categoricity.

Lemma 3.3. *Suppose $\mathcal{L}, \mathcal{L}'$ are both finite relational languages and the \mathcal{L} -structure M is associated to the \mathcal{L}' -structure M' . Then*

- (1) M is ω -categorical if and only if M' is.
- (2) M is cellular if and only if M' is.

Proof. For both parts, as both properties are preserved under reducts, it suffices to prove this in the special case where $\mathcal{L} \subset \mathcal{L}'$ and M' is a q.f. expansion by definitions of M . As \mathcal{L}' is finite relational, by Fact 2.17(5), choose a finite \bar{m} from M' so that every \mathcal{L}' -formula is equivalent in M' to an $\mathcal{L}_{\bar{m}}$ -formula.

For (1), assume M is ω -categorical. Let $k := \text{lg}(\bar{m})$. Then, up to equivalence in M' , for every $n \geq 1$, the number of \mathcal{L}' -formulas $\phi(\bar{x})$ with at most n free variables is at most the number of \mathcal{L} -formulas with at most $(n+k)$ free variables. As the latter number is finite, M' is ω -categorical by Ryll-Nardzewski’s theorem.

For (2), assume M is cellular. Fix a partition $\{K, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$ witnessing this. By Lemma 3.2, M is ω -categorical, hence M' is ω -categorical by (1). Let

$$K^* := K \cup \bigcup \{\bar{c}_{i,j} : \bar{c}_{i,j} \cap \bar{m} \neq \emptyset\}$$

Note that K^* is finite, and after reindexing (removing finitely many tuples) M' can be partitioned as $\{K^*, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$. It is easily checked that this partition witnesses that M' is cellular. \square

As a warm-up to proving Theorem 3.7, we first characterize cellularity among infinite, finitely MA-presented structures.

Proposition 3.4. *The following are equivalent for an infinite, finitely MA-presented \mathcal{L} -structure M :*

- (1) M is cellular;
- (2) M is ω -categorical;
- (3) There is a uniform finite bound on $|\text{acl}(a)|$ for $a \in M$;
- (4) There is a uniform finite bound on $|[a]_{\sim}|$ for $a \in M$.

We remark that this list could be extended. For example such an M is cellular if and only if the binary relation \sim is \mathcal{L} -definable.

Proof. (1) \Rightarrow (2) is by Lemma 3.2.

(2) \Rightarrow (3): This direction holds for any ω -categorical M . As there are only finitely many distinct 1-types, and only finitely many inequivalent 2-formulas $\delta(x, y)$, there is an integer K so that $M \models \exists^{<K} x \delta(x, a)$ among all $a \in M$ and all algebraic formulas $\delta(x, a)$. From this, a uniform bound on $|\text{acl}(a)|$ can be found.

(3) \Rightarrow (4) is immediate, since by Proposition 2.15, $[a]_{\sim} \subseteq \text{acl}(a)$ for every $a \in M$.

(4) \Rightarrow (1): Choose enumerations \bar{c} for each $[a]_{\sim}$ -class in M . As L is finite and there is a uniform bound on $\text{lg}(\bar{c})$, the set of quantifier-free types $qftp(\bar{c})$ occurring is finite. Let $\{p_i(\bar{x}_i) : i \in [n]\}$ enumerate the quantifier-free types that occur infinitely often, and let $K := \bigcup \{\bar{c} : qftp(\bar{c}) \text{ occurs only finitely many times}\}$. Then $K \subseteq M$ is finite. For each $i \in [n]$, let $\{\bar{c}_{i,j} : j \in \omega\}$ enumerate the tuples \bar{c} realizing $p_i(\bar{x}_i)$. Recalling that each $\bar{c}_{i,j}$ is an enumeration of a \sim -class, Lemma 2.9 implies that $\{K, \{\bar{c}_{i,j} \mid i \in [n], j \in \omega\}\}$ is a cellular decomposition of M . \square

Next, we relax the assumption that a mutually algebraic M is finitely MA-presented, but keep the assumption of a finite relational language. For this, we begin with the following lemma.

Lemma 3.5. *Suppose M and M' are associated mutually algebraic structures in finite relational languages \mathcal{L} and \mathcal{L}' , respectively. If there is a uniform finite bound on $|\text{acl}_M(a)|$ for each $a \in M$, then there is a uniform finite bound on $|\text{acl}_{M'}(a)|$ for each $a \in M$.*

Proof. Suppose $|\text{acl}_M(a)| \leq N$ for every $a \in M$. Then, for any finite set $A \subseteq M$, since $\text{acl}_M(A) = \bigcup \{\text{acl}_M(a) : a \in A\}$ for every $a \in A$, $|\text{acl}_M(A)| \leq N \cdot |A|$. Now, it follows from Fact 2.17(5) that there is a finite \bar{m} from M so that every \mathcal{L}' -formula is equivalent on M to an $\mathcal{L}_{\bar{m}}$ -formula. Fix any $a \in M$. From the above, if $b \in \text{acl}_{M'}(a)$, then $b \in \text{acl}_M(a\bar{m})$. Thus, if $k = |\bar{m}|$, then $|\text{acl}_{M'}(a\bar{m})| \leq N(k+1)$. \square

Theorem 3.6. *The following are equivalent for an infinite, mutually algebraic \mathcal{L} -structure M in a finite relational language \mathcal{L} .*

- (1) M is cellular;
- (2) M is ω -categorical;
- (3) There is a uniform finite bound on $|\text{acl}(a)|$ for all $a \in M$;
- (4) For every finitely MA-presented M' associated to M , there is a uniform, finite bound on $|[a]_{\sim}|$ for $a \in M'$;
- (5) For some finitely MA-presented M' associated to M , there is a uniform, finite bound on $|[a]_{\sim}|$ for $a \in M'$.

Proof. (1) \Rightarrow (2) is given by Lemma 3.2.

(2) \Rightarrow (3): As in the proof of Proposition 3.4, (3) holds for any ω -categorical M .

(3) \Rightarrow (4): Choose any finitely MA-presented M' associated to M . By (3) and Lemma 3.5, there is a uniform bound on $|\text{acl}(a)|$ among $a \in M$, hence there is a uniform bound on $|[a]_{\sim}|$ by Proposition 3.4 applied to M' .

(4) \Rightarrow (5): This follows immediately from Lemma 2.22, noting that a finite acceptable set \mathcal{A} can be found.

(5) \Rightarrow (1): Given M' witnessing (5), M' is cellular by Proposition 3.4, hence M is cellular by Lemma 3.3. \square

Finally, we want to identify cellular structures ‘in the rough’ where the original language might not be finite. We do know that on one hand, that if M cellular, it is describable in a finite relational language (see e.g., the proof of Lemma 3.2).

On the other hand, extending to an infinite language (even with only relation symbols) can be troublesome.

Example 2. Let $\mathcal{L} = \{U_n : n \in \omega\}$ and let M be a countable model of the theory of ‘independent unary predicates.’ Then M is mutually algebraic, and in fact, MA-presented in this language. Additionally, $[a]_{\sim} = \{a\} = \text{acl}(a)$ for all $a \in M$, hence there are uniform, finite bounds (namely one) for the sizes of these sets. Despite that, M is not cellular.

Theorem 3.7. *Suppose M is a countable structure in a countable language. Then M is cellular if and only if M is mutually algebraic and ω -categorical.*

Proof. That cellularity implies mutual algebraicity and ω -categoricity is Lemma 3.2. Conversely, assume M is mutually algebraic and ω -categorical. We reduce to the case of a finite relational language, and so finish by Theorem 3.6. Recall Lachlan’s theorem, [?Lach, Proposition 1.6] that a superstable, ω -categorical theory in countable language is ω -stable.

Thus, as M is mutually algebraic implies $\text{Th}(M)$ is weakly minimal and superstable, our ω -categorical M is ω -stable as well. Thus, by e.g., [?Pill, Ch. 3, Lemmas 1.7, 3.9], $\text{Th}(M)$ is interdefinable with some reduct to a finite relational language. It follows that M is cellular by Theorem 3.6. \square

From these results, we obtain a characterization of cellularity that does not mention mutual algebraicity.

Definition 3.8. A countable structure M is *monadically ω -categorical* if every expansion of M by (finitely many) unary predicates is ω -categorical.

Theorem 3.9. *A countable structure M is monadically ω -categorical if and only if M is cellular.*

Proof. It is easy to see that if M is cellular, then any expansion by finitely many unary predicates still admits a partition as in Definition 3.1, although the value of n for the partition might increase.

Now suppose M is monadically ω -categorical. If M is unstable, then by [7BS, Lemma 8.1.3], there is some unary expansion M' of M and formula $\phi(x, y)$ (on singletons) with the strict order property in M' , so some $\psi(x, y)$ defines a partial order with infinite chains. Adding another unary predicate defining an infinite discrete subset of one such chain yields infinitely many 2-types.

Suppose M is also stable. If M is mutually algebraic, we are finished by Theorem 3.7, so suppose not. Then by [7Tri], there is some unary expansion M' of M admitting a definable equivalence relation with infinitely many infinite classes. Enumerating ω of these classes and adding a unary predicate picking n elements from the n^{th} class yields infinitely many 1-types. \square

We remark that the restriction of adding only finitely many unary predicates is necessary, see e.g., Example 2 – If \mathcal{L} is just equality and M is an infinite \mathcal{L} -structure, then M is cellular, but there is an expansion by ω unary predicates that is not cellular.

We close with a technical result that is used by the authors in [7BLsib]. This allows us to produce an extension of a mutually algebraic non-cellular model analogous to an extension of a non-mutually algebraic model by infinitely many infinite arrays all realizing the same quantifier-free type. Thus we may reproduce the “many models” behavior of the non-mutually algebraic case in the mutually algebraic non-cellular case.

Theorem 3.10. *Let \mathcal{L} be finite relational, and suppose M is mutually algebraic but non-cellular countable \mathcal{L} -structure. Then there is some $M^* \succ M$ such that M^* contains infinitely many new infinite MA-connected components, pairwise isomorphic over M . Furthermore, we may take the universe of M^* to be the universe of M together with these new components.*

Proof. By naming finitely many constants, we may assume M is MA-presented. By Proposition 3.4, there is no finite bound on the size of MA-connected components in M .

Claim. We may find an elementary extension of M adding a single new infinite MA-component.

Proof of Claim. We proceed by compactness. Let r be the maximum arity of the language \mathcal{L} and expand \mathcal{L}_M by new constants $\{c_{i,j} \mid i \in \omega, j \in [r]\}$. As notation, let $\bar{c}_i = (c_{i,1}, \dots, c_{i,r})$. Consider the following theory:

- (1) The elementary diagram of M .

- (2) For every $i \in \omega$, $c_{i,j}$ is not equal to any $m \in M$.
- (3) For every $i \in \omega$, some $R \in \mathcal{L}$ holds on some initial subtuple of \bar{c}_i . We now let \bar{c}'_i denote the maximal initial subtuple such that some $R \in \mathcal{L}$ holds on \bar{c}'_i .
- (4) For every $i \in \omega$, $\bar{c}'_i \cap \bar{c}'_{i+1} \neq \emptyset$.
- (5) For every $i \neq j \in \omega$, $\bar{c}'_i \neq \bar{c}'_j$.

If this theory were satisfiable, we would get an elementary extension of M , and $[c_{0,0}]_{\sim}$ would be the desired infinite MA-connected component. So consider a finite subset S of the sentences, and let $F \subset M$ be the finite set of realizations of \mathcal{L}_M -constants mentioned in S . Let $n \in \omega$ be such that $n - 1$ is the largest first index of any $c_{i,j}$ occurring in S . We must interpret $(\bar{c}'_0, \dots, \bar{c}'_{n-1})$ so that it satisfies a linked \mathcal{L} -atomic conjunction, avoids F , and the subtuples are distinct. Given $x, y \in M$, let $d(x, y)$ be the minimum number of conjuncts of any \mathcal{L} -atomic conjunction satisfied by any tuple containing x and y (and $d(x, y) = \infty$ if they are in different MA-connected components). Then it suffices to find some $x \in M$ such that $|[x]_{\sim}| > nr$ and $d(x, f) > n$ for every $f \in F$, since we may then find a tuple containing x and avoiding F that satisfies a linked \mathcal{L} -atomic conjunction with n conjuncts (and with the subtuples for each conjunct distinct), and may take \bar{c}'_i to be the subtuple satisfying the i^{th} conjunct.

Since M is MA-presented, there is some $K \in \omega$ such that for every $m \in M$, $|\{x \in M \mid d(m, x) \leq k\}| \leq K^k$ for every $k \in \omega$. As there is no finite bound on the size of MA-connected components in M , there is some MA-connected component of size greater than $|F| \cdot K^n$, and so we may take the desired x in that component. We may also need to interpret constants appearing in $\bigcup_i \{\bar{c}_i \setminus \bar{c}'_i\}$, but we may choose all such points to lie in a different MA-connected component. \diamond

Working inside a sufficiently saturated elementary extension of M , we may take an infinite set $\{c_i \mid i \in \omega\}$ of realizations of $\text{tp}(c_{0,0}/M)$, pairwise independent over M . Then $\{[c_i]_{\sim} \mid i \in \omega\}$ are the desired infinite MA-connected components, pairwise isomorphic over M . Furthermore, by [[?MA](#), Proposition 4.2], $M \prec M \cup \{[c_i]_{\sim} \mid i \in \omega\}$. \square

REFERENCES