

Borel Complexity and Potential Canonical Scott Sentences

Douglas Ulrich, Richard Rast, and Michael C. Laskowski*

Department of Mathematics

University of Maryland, College Park

November 15, 2016

Abstract

We define and investigate *HC*-forcing invariant formulas of set theory, whose interpretations in the hereditarily countable sets are well behaved under forcing extensions. This leads naturally to a notion of cardinality $||\Phi||$ for sentences Φ of $L_{\omega_1, \omega}$, which counts the number of sentences of $L_{\infty, \omega}$ that, in some forcing extension, become a canonical Scott sentence of a model of Φ . We show this cardinal bounds the complexity of $(\text{Mod}(\Phi), \cong)$, the class of models of Φ with universe ω , by proving that $(\text{Mod}(\Phi), \cong)$ is not Borel reducible to $(\text{Mod}(\Psi), \cong)$ whenever $||\Psi|| < ||\Phi||$. Using these tools, we analyze the complexity of the class of countable models of four complete, first-order theories T for which $(\text{Mod}(T), \cong)$ is properly analytic, yet admit very different behavior. We prove that both ‘Binary splitting, refining equivalence relations’ and Koerwien’s example [14] of an eni-depth 2, ω -stable theory have $(\text{Mod}(T), \cong)$ non-Borel, yet neither is Borel complete. We give a slight modification of Koerwien’s example that also is ω -stable, eni-depth 2, but is Borel complete. Additionally, we prove that $I_{\infty, \omega}(\Phi) < \beth_{\omega_1}$ whenever $(\text{Mod}(\Phi), \cong)$ is Borel.

1 Introduction

In their seminal paper [3], Friedman and Stanley define and develop a notion of *Borel reducibility* among classes C of structures with universe ω in a fixed, countable language L that are Borel and invariant under permutations of ω . It is well known (see e.g., [11] or [4]) that such classes are of the form $\text{Mod}(\Phi)$, the set of models of Φ whose universe is precisely ω for some sentence $\Phi \in L_{\omega_1, \omega}$. A *Borel reduction* is a Borel function $f : \text{Mod}(\Phi) \rightarrow \text{Mod}(\Psi)$ that satisfies $M \cong N$ if and only if $f(M) \cong f(N)$. One says that $\text{Mod}(\Phi)$ is *Borel reducible* to $\text{Mod}(\Psi)$, written $(\text{Mod}(\Phi), \cong) \leq_B (\text{Mod}(\Psi), \cong)$ or more typically $\Phi \leq_B \Psi$, if there is a Borel reduction $f : \text{Mod}(\Phi) \rightarrow \text{Mod}(\Psi)$; and the two classes are *Borel equivalent* if there are Borel reductions in both directions. As Borel reducibility is transitive, \leq_B induces a pre-order on $\{\text{Mod}(\Phi) : \Phi \in L_{\omega_1, \omega}\}$. In [3], Friedman and Stanley show that among Borel invariant classes, there is a maximal class with respect to \leq_B . We say Φ is *Borel complete* if it is in this maximal class. Examples include the theories of graphs, linear orders, groups, and fields.

It is easily seen that for any $\Phi \in L_{\omega_1, \omega}$, the isomorphism relation on $\text{Mod}(\Phi)$ is analytic, but in many cases it is actually Borel. The isomorphism relation on any Borel complete class is properly analytic (i.e., not Borel) but prior to this paper there were few examples known of classes $\text{Mod}(\Phi)$ where the isomorphism

*All three authors were partially supported by NSF Research Grant DMS-1308546.

relation is properly analytic, but not Borel complete. Indeed, the authors are only aware of the example of abelian p -groups that first appeared in [3]. The class of abelian p -groups is expressible by a sentence of $L_{\omega_1, \omega}$, but no first-order example of this phenomenon was known.

To date, the study of classes $(\text{Mod}(\Phi), \cong)$ is much more developed when the isomorphism relation is Borel than when it is not. Here, however, we define the set of *short* sentences Φ (see Definition 3.6), that properly contain the set of sentences for which isomorphism is Borel but exclude the Borel complete sentences, and develop criteria for concluding $\text{Mod}(\Phi) \not\leq_B \text{Mod}(\Psi)$ among pairs Φ, Ψ of short sentences. Furthermore, Theorem 3.10 asserts that there can never be a Borel reduction from a non-short Φ into a short Ψ . From this technology, we are able to discern more about $\text{Mod}(\Phi)$ when \cong is Borel. For example, Corollary 3.12 asserts that

If \cong is Borel on $\text{Mod}(\Phi)$, then $I_{\infty, \omega}(\Phi) < \beth_{\omega_1}$.

That is, among models of Φ of any cardinality, there are fewer than \beth_{ω_1} pairwise back-and-forth inequivalent models.

In this paper we work in ZFC, and so formalize all of model theory within ZFC. We do this in any reasonable way. For instance, countable languages L are construed as being elements of HC (the set of hereditarily countable sets) and the set of sentences of $L_{\omega_1, \omega}$ form a subclass of HC. Moreover, basic operations, such as describing the set of subformulas of a given formula, are absolute class functions of HC. As well, if $L \in \text{HC}$ is a language, then for every L -structure M , the canonical Scott sentence $\text{css}(M)$ is in particular a set; and if M is countable then $\text{css}(M) \in L_{\omega_1, \omega}$ is an element of HC. Moreover the class function css is highly absolute.

One of our central ideas is the following. Given $\Phi \in L_{\omega_1, \omega}$, let $\text{CSS}(\Phi) \subseteq \text{HC}$ denote the set $\{\text{css}(M) : M \in \text{Mod}(\Phi)\}$. Then any Borel map $f : \text{Mod}(\Phi) \rightarrow \text{Mod}(\Psi)$ induces an HC-definable function $f^* : \text{CSS}(\Phi) \rightarrow \text{CSS}(\Psi)$. This leads us to the investigation of definable subclasses of HC and definable maps between them. We begin by restricting our notion of classes to those definable by *HC-forcing invariant* formulas (see Definition 2.1). Three straightforward consequences of the Product Forcing Lemma allow us to prove that these classes are well-behaved. It is noteworthy that we do not define HC-forcing invariant formulas syntactically. Whereas it is true that every Σ_1 -formula is HC-forcing invariant, determining precisely which classes are HC-forcing invariant depends on our choice of \mathbb{V} .

The second ingredient of our development is that every set A in \mathbb{V} is *potentially in HC*, i.e., there is a forcing extension $\mathbb{V}[G]$ of \mathbb{V} (indeed a Levy collapse suffices) such that $A \in \text{HC}^{\mathbb{V}[G]}$. Given an HC-forcing invariant $\varphi(x)$, we define φ_{ptl} – that is, the *potential solutions to φ* – to be those $A \in \mathbb{V}$ for which $(\text{HC}^{\mathbb{V}[G]}, \epsilon) \models \varphi(A)$ whenever $A \in \text{HC}^{\mathbb{V}[G]}$. The definition of HC-forcing invariance makes this notion well-defined. Thus, given a sentence Φ , one can define $\text{CSS}(\Phi)_{\text{ptl}}$, which should be read as the class of ‘potential canonical Scott sentences’ i.e., the class of all $\varphi \in L_{\infty, \omega}$ such that in some forcing extension $\mathbb{V}[G]$, φ is the canonical Scott sentence of some countable model of Φ . We define Φ to be *short* if $\text{CSS}(\Phi)_{\text{ptl}}$ is a **set** as opposed to a proper class and define the *potential cardinality* of Φ , denoted $||\Phi||$, to be the (usual) cardinality of $\text{CSS}(\Phi)_{\text{ptl}}$ if Φ is short, or ∞ otherwise. By tracing all of this through, with Theorem 3.10(2), we see that

If $\|\Psi\| < \|\Phi\|$, then there cannot be a Borel reduction $f : \text{Mod}(\Phi) \rightarrow \text{Mod}(\Psi)$.

Another issue that is raised by our investigation is a comparison of the class the potential canonical Scott sentences $\text{CSS}(\Phi)_{\text{ptl}}$ with the class $\text{CSS}(\Phi)_{\text{sat}}$, consisting of all sentences of $L_{\infty, \omega}$ that are canonical Scott sentences of some model $M \models \Phi$ with $M \in \mathbb{V}$. Clearly, the latter class is contained in the former, and we call Φ *grounded* (see Definition 3.8) if equality holds. We show that the incomplete theory REF of refining equivalence relations is grounded. By contrast, the theory TK, defined in Section 6, is a complete, ω -stable theory for which $|\text{CSS}(\text{TK})_{\text{sat}}| = \beth_2$, while $\text{CSS}(\text{TK})_{\text{ptl}}$ is a proper class.

Sections 4-6 apply this technology. Section 4 discusses continuous actions by compact groups on Polish spaces. In addition to being of interest in its own right, the results there are also used in Section 6. Sections 5 and 6 discuss four complete first-order theories that are not very complicated stability-theoretically, yet the isomorphism relation is properly analytic in each case. We summarize our findings by:

REF(inf) is the theory of ‘infinitely splitting, refining equivalence relations’. Its language is $L = \{E_n : n \in \omega\}$. It asserts that each E_n is an equivalence relation, E_0 consists of a single class, each E_{n+1} refines E_n , and each E_n -class is partitioned into infinitely many E_{n+1} -classes. REF(inf) is one of the standard examples of a stable, unsuperstable theory. Then:

- REF(inf) is Borel complete, in fact, it is λ -Borel complete for all infinite λ (see Definition 3.15);
- Therefore, REF(inf) is not short;
- Therefore, \cong is not Borel.
- REF(inf) is grounded, i.e., $\text{CSS}(\text{REF}(\text{inf}))_{\text{sat}} = \text{CSS}(\text{REF}(\text{inf}))_{\text{ptl}}$.

REF(bin) is the theory of ‘binary splitting, refining equivalence relations’. The language is also $L = \{E_n : n \in \omega\}$. The axioms of REF(bin) assert that each E_n is an equivalence relation, E_0 is trivial, each E_{n+1} refines E_n , and each E_n -class is partitioned into exactly two E_{n+1} -classes. REF(bin) is superstable (in fact, weakly minimal) but is not ω -stable. Then:

- T_2 (i.e., ‘countable sets of reals’) is Borel reducible into $\text{Mod}(\text{REF}(\text{bin}))$, but
- REF(bin) is short, with $\|\text{REF}(\text{bin})\| = \beth_2$.
- Therefore, T_3 (i.e., ‘countable sets of countable sets of reals’) is not Borel reducible into $\text{Mod}(\text{REF}(\text{bin}))$;
- Therefore, REF(bin) is not Borel complete.
- \cong is not Borel.
- REF(bin) is grounded.

\mathbf{K} is the Koerwien theory, originating in [14] and defined in Section 6. Koerwien proved that \mathbf{K} is complete, ω -stable, eni-NDOP, and of eni-depth 2. Then:

- T_2 (i.e., ‘countable sets of reals’) is Borel reducible into $\text{Mod}(\mathbf{K})$, and
- \mathbf{K} is short, with $|\text{Mod}(\mathbf{K})| = \beth_2$.
- Therefore, T_3 (i.e., ‘countable sets of countable sets of reals’) is not Borel reducible into $\text{Mod}(\mathbf{K})$;
- Therefore, $\text{Mod}(\mathbf{K})$ is not Borel complete;
- Nonetheless, \cong is not Borel (this was proved by Koerwien in [14]).
- Whether \mathbf{K} is grounded or not remains open.

\mathbf{TK} is a ‘tweaked version of \mathbf{K} ’ and is also defined in Section 6. \mathbf{TK} is also complete, ω -stable, eni-NDOP, of eni-depth 2, and is very much like the theory \mathbf{K} , however the automorphism groups of models of \mathbf{TK} induce a more complicated group of elementary permutations of $\text{acl}(\emptyset)$ than do the automorphism groups of models of \mathbf{K} . Then:

- \mathbf{TK} is Borel complete, hence not short, hence \cong is not Borel.
- \mathbf{TK} is not grounded; in fact $\text{CSS}(\mathbf{TK})_{\text{ptl}}$ is a proper class, while $|\text{CSS}(\mathbf{TK})_{\text{sat}}| = \beth_2$.
- Since $|\text{CSS}(\mathbf{TK})_{\text{sat}}| = \beth_2$, \mathbf{TK} is not λ -Borel complete for sufficiently large λ .

Many of the ideas for this arose from the first author, who read and abstracted ideas about absolutely $\Delta_{\frac{1}{2}}$ -formulas in [3] and ideas that were discussed in Chapter 9 of [5]. Only recently, the authors became aware of work on ‘pinned equivalence relations’ as surveyed in e.g., [17]. The whole of this paper was written independently of the development there. In terms of notation, HC always denotes $H(\aleph_1)$, the set of all sets whose transitive closure is countable. For a forcing extension $\mathbb{V}[G]$ of \mathbb{V} , $\text{HC}^{\mathbb{V}[G]}$ denotes those sets that are hereditarily countable in $\mathbb{V}[G]$. Throughout, we only consider set forcings, so when we write ‘for $\mathbb{V}[G]$ a forcing extension of $\mathbb{V} \dots$ ’ we are quantifying over all (set) forcing notions $\mathbb{P} \in \mathbb{V}$ and all \mathbb{P} -generic filters G .

Finally, we are grateful to the referee for their very careful reading and their finding simplifications to some of our arguments.

2 A notion of cardinality for some classes of HC

We develop a notion of cardinality on certain well-behaved subclasses of HC in terms of the existence or non-existence of certain injective maps. Behind the scenes, we rely heavily on the fact that all sets A in \mathbb{V} are ‘potentially’ elements of HC , the set of hereditarily countable sets. For example, if κ is the cardinality of the transitive closure of A and we take \mathbb{P} be the Levy collapsing poset $\text{Coll}(\kappa^+, \omega_1)$ that collapses κ^+ to ω_1 , then for any choice G of a generic filter, $A \in \text{HC}^{\mathbb{V}[G]}$.

2.1 HC-forcing invariant formulas

We begin with our principal definitions.

Definition 2.1. Suppose $\varphi(x)$ is any formula of set theory, possibly with a hidden parameter from HC.

- $\varphi(\text{HC}) = \{a \in \text{HC} : (\text{HC}, \in) \models \varphi(a)\}$.
- If $\mathbb{V}[G]$ is a forcing extension of \mathbb{V} , then $\varphi(\text{HC})^{\mathbb{V}[G]} = \{a \in \text{HC}^{\mathbb{V}[G]} : \mathbb{V}[G] \models 'a \in \varphi(\text{HC})'\}$.
- $\varphi(x)$ is *HC-forcing invariant* if, for every twice-iterated forcing extension $\mathbb{V}[G][G']$,

$$\varphi(\text{HC})^{\mathbb{V}[G][G']} \cap \text{HC}^{\mathbb{V}[G]} = \varphi(\text{HC})^{\mathbb{V}[G]}$$

The reader is cautioned that when computing $\varphi(\text{HC})^{\mathbb{V}[G]}$, the quantifiers of φ range over $\text{HC}^{\mathbb{V}[G]}$ as opposed to the whole of $\mathbb{V}[G]$. Visibly, the class of HC-forcing invariant formulas is closed under boolean combinations. Note that by Shoenfield's Absoluteness Theorem, e.g., Theorem 25.20 of [8], any Σ_1^1 subset of \mathbb{R} is HC-forcing invariant. There is also the closely related Lévy Absoluteness Principle, which has various forms (e.g., Theorem 9.1 of [2] or Section 4 of [9]); we give a version more convenient to us.

Lemma 2.2. *If $\mathbb{V}[G]$ is any forcing extension, and if $\varphi(x)$ is a Σ_1 formula of set theory, then for every $a \in \text{HC}$, $\text{HC} \models \varphi(a)$ if and only if $\text{HC}^{\mathbb{V}[G]} \models \varphi(a)$. In particular, Σ_1 -formulas are HC-forcing invariant.*

Proof. This can be proved using Shoenfield's Absoluteness Theorem, as in Exercise 25.4 of [8] (which readily relativizes to allow a parameter from HC). More directly, let $\varphi(x)$ be $\exists y \psi(y, x)$, where $\psi(y, x)$ is Δ_0 . Let $\mathbb{V}[G]$ be any forcing extension of \mathbb{V} and assume $\text{HC}^{\mathbb{V}[G]} \models \varphi(a)$ for some $a \in \text{HC}^{\mathbb{V}}$.

In $\mathbb{V}[G]$, there is a countable, transitive set M containing a and a witness b to $\varphi(a)$. Choose a bijection $j : M \rightarrow \omega$ with $j(b) = 1$ and $j(a) = 0$. The image of \in restricted to M is a well-founded, extensional relation E such that $(\omega, E) \models \psi(1, 0)$.

Let $\kappa = (\omega_1)^{\mathbb{V}[G]}$. Working in \mathbb{V} , there is an ω -tree T whose ill-foundedness witnesses the existence of E as above. Without going into all the details, T consists of certain pairs $(\sigma(x_0, \dots, x_{n-1}), r)$, where $\sigma(x_0, \dots, x_{n-1})$ is a formula of set theory in the listed variables, $r : n \rightarrow \kappa$ is a function, and we arrange that when $\langle (\sigma_n, r_n) : n \in \omega \rangle$ is a branch through T , then $p := \bigcup_n \sigma_n$ is a complete type in the variables $(x_n : n \in \omega)$, such that the induced model (ω, E') is extensional, and $(\omega, E') \models \psi(1, 0)$, and 0 codes a , and $r := \bigcup_n r_n : \omega \rightarrow \kappa$ is a rank function witnessing that (ω, E') is well-founded.

As (ω, E) is well-founded in $\mathbb{V}[G]$, T is ill-founded in $\mathbb{V}[G]$. As well-foundedness is Δ_1 , this implies T is ill-founded in \mathbb{V} as well. Thus, $\psi(x, a)$ has a witness in $\text{HC}^{\mathbb{V}}$. \square

For more complicated formulas, whether or not $\varphi(x)$ is HC-forcing invariant or not may well depend on the choice of set-theoretic universe. For example, consider the formula $\varphi(x) := (x = \emptyset) \vee (\mathbb{V} \neq \mathbb{L})$. Then $\varphi(\text{HC})$ is equal to $\{\emptyset\}$ if $\text{HC} \subseteq \mathbb{L}$, and $\varphi(\text{HC}) = \text{HC}$ otherwise. Because ' $\text{HC} \not\subseteq \mathbb{L}$ ' holding in \mathbb{V} implies that

it holds in any forcing extension $\mathbb{V}[H]$, it follows that the formula $\varphi(x)$ is HC-forcing invariant if and only if $\text{HC} \not\subseteq \mathbb{L}$.

Before continuing, we state three set-theoretic lemmas that form the lynchpin of our development. Lemma 2.3 is a simple consequence of our definitions. Lemma 2.4 is well-known. It is mentioned in the proof of Theorem 9.4 in [5]; a full proof is given in the more recent [10]. Lemma 2.5 is just a rephrasing of Lemma 2.4. The key tool for all of these lemmas is the Product Forcing Lemma, see e.g., Lemma 15.9 of [8], which states that given any $\mathbb{P}_1 \times \mathbb{P}_2$ -generic filter G , if G_ℓ is the projection of G onto \mathbb{P}_ℓ , then $G = G_1 \times G_2$, each G_ℓ is \mathbb{P}_ℓ -generic, and $\mathbb{V}[G] = \mathbb{V}[G_1][G_2] = \mathbb{V}[G_2][G_1]$ (i.e., G_ℓ meets every dense subset of \mathbb{P}_ℓ in $\mathbb{V}[G_{3-\ell}]$).

Lemma 2.3. *Suppose $\mathbb{P}_1, \mathbb{P}_2 \in \mathbb{V}$ are notions of forcing and $\varphi(x)$ is HC-forcing invariant (possibly with a hidden parameter from HC). If $A \in V$ and, for $\ell = 1, 2$, H_ℓ is \mathbb{P}_ℓ -generic and $\mathbb{V}[H_\ell] \models A \in \text{HC}$, then $\mathbb{V}[H_1] \models A \in \varphi(\text{HC})$ if and only if $\mathbb{V}[H_2] \models A \in \varphi(\text{HC})$. (The filters H_1 and H_2 are not assumed to be mutually generic.)*

Proof. Assume this were not the case. By symmetry, choose $p_1 \in H_1$ such that $p_1 \Vdash \check{A} \in \varphi(\text{HC})$, and choose $p_2 \in H_2$ such that $p_2 \Vdash \check{A} \in \text{HC} \wedge \check{A} \notin \varphi(\text{HC})$. Let G be a $\mathbb{P}_1 \times \mathbb{P}_2$ -generic filter with $(p_1, p_2) \in G$. Write $G = G_1 \times G_2$, hence $\mathbb{V}[G] = \mathbb{V}[G_1][G_2] = \mathbb{V}[G_2][G_1]$. As $p_1 \in G_1$ and $p_2 \in G_2$, we have $\mathbb{V}[G_1] \models A \in \varphi(\text{HC})$ and $\mathbb{V}[G_2] \models A \in \text{HC} \wedge A \notin \varphi(\text{HC})$. But applying the HC-forcing invariance of φ twice, we get that $A \in \varphi(\text{HC})^{\mathbb{V}[G_1]}$ iff $A \in \varphi(\text{HC})^{\mathbb{V}[G_1][G_2]}$ iff $A \in \varphi(\text{HC})^{\mathbb{V}[G_2]}$, a contradiction. \square

Lemma 2.4. *Suppose \mathbb{P}_1 and \mathbb{P}_2 are both notions of forcing in \mathbb{V} . If G is a $\mathbb{P}_1 \times \mathbb{P}_2$ -generic filter and $G = G_1 \times G_2$, then $\mathbb{V} = \mathbb{V}[G_1] \cap \mathbb{V}[G_2]$.*

Proof. This is Corollary 2.3 of [10]. \square

Lemma 2.5. *Let $\theta(x)$ be any formula of set theory, possibly with hidden parameters from \mathbb{V} , and let $\mathbb{V}[G]$ be any forcing extension of \mathbb{V} . Suppose that there is some $b \in \mathbb{V}[G]$ such that for every forcing extension $\mathbb{V}[G][H]$ of $\mathbb{V}[G]$,*

$$\mathbb{V}[G][H] \models \theta(b) \wedge \exists^=1 x \theta(x)$$

Then $b \in \mathbb{V}$.

Proof. Fix $\theta(x), \mathbb{V}[G]$ and b as above. Let $\mathbb{P} \in \mathbb{V}$ be the forcing notion for which G is \mathbb{P} -generic. Let τ be a \mathbb{P} -name such that $b = \text{val}(\tau, G)$. Choose $p \in G$ such that

$$p \Vdash \text{“for all forcing notions } \mathbb{Q}, \Vdash_{\mathbb{Q}} \theta(\check{\tau}) \wedge \exists^=1 x \theta(x)\text{”}$$

Let H be \mathbb{P} -generic over $\mathbb{V}[G]$ with $p \in H$. So $G \times H$ is $\mathbb{P} \times \mathbb{P}$ -generic over \mathbb{V} . Let $i_1, i_2 : \mathbb{P} \rightarrow \mathbb{P} \times \mathbb{P}$ be the canonical injections. Then, since $(p, p) \in G \times H$, we have that

$$\mathbb{V}[G][H] \models \theta(\text{val}(i_1(\tau), G \times H)) \wedge \theta(\text{val}(i_2(\tau), G \times H)) \wedge \exists^=1 x \theta(x).$$

Hence $\mathbb{V}[G][H] \models \text{val}(i_1(\tau), G \times H) = \text{val}(i_2(\tau), G \times H)$ and so by Lemma 2.4, $\text{val}(i_1(\tau), G \times H) \in \mathbb{V}$. But $\text{val}(i_1(\tau), G \times H) = b$ so we are done. \square

Lemma 2.3 lends credence to the following definition.

Definition 2.6. Suppose that $\varphi(x)$ is HC-forcing invariant. Then φ_{ptl} is the class of all sets A such that $A \in \mathbb{V}$ and, for some (equivalently, for every) forcing extension $\mathbb{V}[G]$ of \mathbb{V} with $A \in \text{HC}^{\mathbb{V}[G]}$, we have $A \in \varphi(\text{HC})^{\mathbb{V}[G]}$.

As motivation for the notation used in the definition above, φ_{ptl} describes the class of all $A \in \mathbb{V}$ that are *potentially* in $\varphi(\text{HC})$. We are specifically interested in those HC-forcing invariant φ for which φ_{ptl} is a **set** as opposed to a proper class.

Definition 2.7. An HC-forcing invariant formula $\varphi(x)$ is *short* if φ_{ptl} is a set.

We begin with some easy observations. As notation, if C is a subclass of \mathbb{V} , then define $\mathcal{P}(C)$ to be all sets A in \mathbb{V} such that every element of A is in C . (This definition is only novel when C is a proper class.) Similarly, $\mathcal{P}_{\aleph_1}(C)$ is the class of all sets $A \in \mathcal{P}(C)$ that are countable (in \mathbb{V} !). Let $\delta(x)$ be the formula

$$\delta(x) := \exists h[h : x \rightarrow \omega \text{ is 1-1}]$$

Given a formula $\varphi(x)$, let $\mathcal{P}(\varphi)(y)$ denote the formula $\forall x(x \in y \rightarrow \varphi(x))$ and let $\mathcal{P}_{\aleph_1}(\varphi)$ denote $\mathcal{P}(\varphi)(y) \wedge \delta(y)$.

Lemma 2.8. 1. The *ptl*-operator commutes with boolean combinations, i.e., if φ and ψ are both HC-forcing invariant, then $(\varphi \wedge \psi)_{\text{ptl}} = \varphi_{\text{ptl}} \cap \psi_{\text{ptl}}$ and $(\neg\varphi)_{\text{ptl}} = \mathbb{V} \setminus \varphi_{\text{ptl}}$.

2. $\delta(\text{HC}) = \text{HC}$. In particular $\delta(x)$ is HC-forcing invariant and $\delta_{\text{ptl}} = \mathbb{V}$.

3. If φ is HC-forcing invariant, then so are both $\mathcal{P}(\varphi)$ and $\mathcal{P}_{\aleph_1}(\varphi)$. Moreover, $\mathcal{P}(\varphi)(\text{HC}) = \mathcal{P}_{\aleph_1}(\varphi)(\text{HC})$ and $\mathcal{P}(\varphi_{\text{ptl}}) = (\mathcal{P}(\varphi))_{\text{ptl}} = (\mathcal{P}_{\aleph_1}(\varphi))_{\text{ptl}}$.

4. Suppose $s : \omega \rightarrow \text{HC}$ is any map such that for each n , $\varphi(x, s(n))$ is HC-forcing invariant. (Recall that HC-forcing invariant formulas are permitted to have a parameter from HC.) Then $\psi(x) := \exists n(n \in \omega \wedge \varphi(x, s(n)))$ is HC-forcing invariant and $\psi_{\text{ptl}} = \bigcup_{n \in \omega} \varphi(x, s(n))_{\text{ptl}}$.

Proof. The verification of (1) and (2) is immediate, simply by unpacking definitions.

(3) That $\mathcal{P}(\varphi)(y)$ and $\mathcal{P}_{\aleph_1}(\varphi)(y)$ are HC-forcing invariant is routine. Since $\delta(\text{HC}) = \text{HC}$ we have $\mathcal{P}(\varphi)(\text{HC}) = \mathcal{P}_{\aleph_1}(\varphi)(\text{HC})$ and so $(\mathcal{P}(\varphi))_{\text{ptl}} = (\mathcal{P}_{\aleph_1}(\varphi))_{\text{ptl}}$.

The only other thing to check is that $\mathcal{P}(\varphi_{\text{ptl}}) = (\mathcal{P}(\varphi))_{\text{ptl}}$. Begin by choosing any $A \in \mathcal{P}(\varphi)_{\text{ptl}}$. We must show that every element $a \in A$ is in φ_{ptl} . Fix any element $a \in A$. Choose any forcing extension $\mathbb{V}[G]$ of \mathbb{V} with $A \in \text{HC}^{\mathbb{V}[G]}$. Then

$$\text{HC}^{\mathbb{V}[G]} \models \forall x(x \in A \rightarrow \varphi(x))$$

Since $A \in \text{HC}^{\mathbb{V}[G]}$, $a \in \text{HC}^{\mathbb{V}[G]}$ as well. Thus, $\text{HC}^{\mathbb{V}[G]} \models \varphi(a)$, so $a \in \varphi_{\text{ptl}}$.

Conversely, suppose $A \in \mathcal{P}(\varphi_{\text{ptl}})$. This means that $A \in \mathbb{V}$ and every element of A is in φ_{ptl} . Choose any forcing extension $\mathbb{V}[G]$ of \mathbb{V} such that $A \in \text{HC}^{\mathbb{V}[G]}$. As $\text{HC}^{\mathbb{V}[G]}$ is transitive, every element $a \in A$ is also

an element of $\text{HC}^{\mathbb{V}[G]}$. Thus, $\text{HC}^{\mathbb{V}[G]} \models \varphi(a)$ for every $a \in A$. That is, $\text{HC}^{\mathbb{V}[G]} \models \forall x(x \in A \rightarrow \varphi(x))$, so $A \in (\mathcal{P}(\varphi))_{\text{ptl}}$.

(4) Note that $s \in \text{HC}$ so can be used as a hidden parameter for $\psi(x)$. Choose $A \in \text{HC}$. It is immediate from the definition of ψ that $\text{HC} \models \psi(A)$ if and only if $\text{HC} \models \varphi(A, s_n)$ for some $n \in \omega$. As this equivalence relativizes to any $A \in \text{HC}^{\mathbb{V}[G]}$, both statements follow. \square

2.2 Strongly definable families

In the previous subsection, we described a restricted vocabulary of formulas. Here, we discuss parameterized families of classes of $\text{HC}^{\mathbb{V}[G]}$ that are describable in this vocabulary.

Definition 2.9. Let φ be HC-forcing invariant. A family $X = (X^{\mathbb{V}[G]} : \mathbb{V}[G] \text{ a forcing extension of } \mathbb{V})$ is *strongly definable via* φ if $X^{\mathbb{V}[G]} = \varphi(\text{HC})^{\mathbb{V}[G]}$ always. We say the family X is a *strongly definable family*, or just *strongly definable*, if it is strongly definable via some HC-forcing invariant formula φ .

We say that two HC-forcing invariant formulas φ and ψ are *persistently equivalent* if $\varphi(\text{HC})^{\mathbb{V}[G]} = \psi(\text{HC})^{\mathbb{V}[G]}$ for every forcing extension $\mathbb{V}[G]$. Persistently equivalent formulas give rise to the same strongly definable family, and if X is strongly definable via both φ and ψ , then φ and ψ are persistently equivalent.

We note that the strongly definable families can also be defined as the HC-forcing invariant formulas modulo persistent equivalence; using this one can verify that all the results of this paper are really theorems of ZFC.

Definition 2.10. If X is strongly definable, define X_{ptl} to be the class of all sets $A \in \mathbb{V}$ such that $A \in X^{\mathbb{V}[G]}$ for some forcing extension $\mathbb{V}[G]$ of \mathbb{V} . We call X *short* if X_{ptl} is a set as opposed to a proper class.

Note that if X is strongly definable via φ , then $X_{\text{ptl}} = \varphi_{\text{ptl}}$ and X is short if and only if φ is short.

We can define operations on the collection of strongly definable families. Of particular interest is the countable power set. That is, in the notation preceding Lemma 2.8, given any HC-forcing invariant $\varphi(x)$, the two HC-invariant formulas $\mathcal{P}(\varphi)$ and $\mathcal{P}_{\aleph_1}(\varphi)$ are persistently equal. Thus, if X is strongly definable via φ , then $\mathcal{P}(\varphi)$ and $\mathcal{P}_{\aleph_1}(\varphi)$ give rise to the same family, which we denote by $\mathcal{P}_{\aleph_1}(X)$.

We begin by enumerating several examples and easy observations that help establish our notation.

Example 2.11. 1. ω is strongly definable via the HC-forcing invariant formula $\varphi_\omega(x) := "x \text{ is a natural number.}"$ Here, $\omega_{\text{ptl}} = \omega$. In particular, ω is short.

2. ω_1 is strongly definable via the HC-forcing invariant formula " x is an ordinal." Here, $(\omega_1)_{\text{ptl}} = ON$, the class of all ordinals. Thus, ω_1 is not short.¹

¹The symbol ω_1 , depending on context, refers to either the least uncountable ordinal $(\omega_1)^\mathbb{V}$, or else the parametrized family $((\omega_1)^{\mathbb{V}[G]} : \mathbb{V}[G] \text{ a forcing extension of } \mathbb{V})$. It is in the latter sense that it makes sense to say that ' ω_1 is strongly definable'. This kind of ambiguity will not create problems, since in practice, given a subset $X \subseteq \text{HC}$, there is only one natural definition of X . Thus there is only one natural way of viewing X as a family parametrized by forcing extensions.

3. The set of reals, $\mathbb{R} = \mathcal{P}_{\aleph_1}(\omega)$, is strongly definable via $\varphi_{\mathbb{R}}(x) := \mathcal{P}(\varphi_\omega)$. By Lemma 2.8(3) $\mathbb{R}_{\text{ptl}} = \mathcal{P}(\omega) = \mathbb{R}$, hence \mathbb{R} is short.
4. $\mathcal{P}_{\aleph_1}(\mathbb{R})$, the set of countable sets of reals, is strongly definable either by $\mathcal{P}_{\aleph_1}(\varphi_{\mathbb{R}})$ or by $\mathcal{P}(\varphi_{\mathbb{R}})$. Via either definition, by Lemma 2.8(3), $(\mathcal{P}_{\aleph_1}(\mathbb{R}))_{\text{ptl}} = \mathcal{P}(\mathcal{P}(\omega))$, hence is short.
5. More generally, if X is short, then it follows from Lemma 2.8(3) that $\mathcal{P}_{\aleph_1}(X)_{\text{ptl}} = \mathcal{P}(X_{\text{ptl}})$ and so $\mathcal{P}_{\aleph_1}(X)$ is short.
6. For any $\alpha < \omega_1$, let HC_α denote the sets in HC whose rank is less than α . Then each HC_α is strongly definable (since the formula “ $\text{rank}(x) < \alpha$ ” is HC -forcing invariant). Also, each $(\text{HC}_\alpha)_{\text{ptl}} = \mathbb{V}_\alpha$, so each HC_α is short.

Notation 2.12. Suppose that X_1, \dots, X_n are each strongly definable families. We say $\psi(X_1, \dots, X_n)$ *holds persistently* if, for every forcing extension $\mathbb{V}[G]$, we have

$$\mathbb{V}[G] \models \psi(X_1^{\mathbb{V}[G]}, \dots, X_n^{\mathbb{V}[G]}).$$

We list three examples of this usage in the definition below.

Definition 2.13. Suppose that f , X and Y are each strongly definable families.

- The notation $f : X \rightarrow Y$ *persistently* means that $f^{\mathbb{V}[G]} : X^{\mathbb{V}[G]} \rightarrow Y^{\mathbb{V}[G]}$ for all forcing extensions $\mathbb{V}[G]$ of \mathbb{V} .
- The notation $f : X \rightarrow Y$ *is persistently injective* means that $f : X \rightarrow Y$ persistently and additionally, for all forcing extensions $\mathbb{V}[G]$ of \mathbb{V} , $f^{\mathbb{V}[G]} : X^{\mathbb{V}[G]} \rightarrow Y^{\mathbb{V}[G]}$ is 1-1.
- The notion $f : X \rightarrow Y$ *is persistently bijective* means $f^{-1} : X \rightarrow Y$ is strongly definable and both $f : X \rightarrow Y$ and $f^{-1} : Y \rightarrow X$ are persistently injective.

The reader is cautioned that when $f : X \rightarrow Y$ persistently (or is persistently injective), the image of f need not be strongly definable. Indeed the “image” of a strongly definable function is not well-behaved in many respects, including the lack of a surjectivity statement in the following proposition.

Proposition 2.14. *Suppose that f , X , and Y are each strongly definable.*

1. *Suppose $f : X \rightarrow Y$ persistently. Then $f_{\text{ptl}} : X_{\text{ptl}} \rightarrow Y_{\text{ptl}}$, i.e., f_{ptl} is a class function with domain X_{ptl} and image contained in Y_{ptl} .*
2. *If $f : X \rightarrow Y$ is persistently injective, then $f_{\text{ptl}} : X_{\text{ptl}} \rightarrow Y_{\text{ptl}}$ is injective as well.*
3. *If $f : X \rightarrow Y$ is persistently bijective, then $f_{\text{ptl}} : X_{\text{ptl}} \rightarrow Y_{\text{ptl}}$ is bijective.*

Proof. (1) It is obvious that f_{ptl} is a class of ordered pairs and is single-valued. As well, by definition of f_{ptl} , if $(a, b) \in f_{\text{ptl}}$, then for any forcing extension $\mathbb{V}[G]$ of \mathbb{V} with $(a, b) \in \text{HC}^{\mathbb{V}[G]}$ we have $(a, b) \in f^{\mathbb{V}[G]}$. Thus $\text{dom}(f_{\text{ptl}}) \subseteq X_{\text{ptl}}$ and $\text{im}(f_{\text{ptl}}) \subseteq Y_{\text{ptl}}$. To see that $\text{dom}(f_{\text{ptl}})$ is equal to X_{ptl} requires Lemma 2.5. Choose any $a \in X_{\text{ptl}}$ and look at the formula $\theta(a, y)^{\text{HC}}$, where $\theta(x, y)$ defines f , i.e., f is strongly definable via θ . Let $\mathbb{V}[G]$ be any forcing extension of \mathbb{V} in which $a \in X^{\mathbb{V}[G]}$ (in particular $a \in \text{HC}^{\mathbb{V}[G]}$). Let $b := f^{\mathbb{V}[G]}(a)$. The definition of persistence tells us that the hypotheses of Lemma 2.5 apply, hence $b \in \mathbb{V}$. Thus, $(a, b) \in f_{\text{ptl}}$ and $a \in \text{dom}(f_{\text{ptl}})$.

(2) Choose $a, b, c \in \mathbb{V}$ such that $(a, c), (b, c) \in f_{\text{ptl}}$. Choose a forcing extension $\mathbb{V}[G]$ of \mathbb{V} such that $a, b, c \in \text{HC}^{\mathbb{V}[G]}$. Thus, $(a, c), (b, c) \in f^{\mathbb{V}[G]}$. As f is persistently injective, it follows that $a = b$ holds in $\mathbb{V}[G]$ and hence in \mathbb{V} . So f_{ptl} is injective.

(3) From (2) we get that both f_{ptl} and $(f^{-1})_{\text{ptl}}$ are injective. It follows that f_{ptl} is bijective. \square

We close this subsection with a characterization of surjectivity. Its proof is simply an unpacking of the definitions. However, $f : X \rightarrow Y$ being persistently surjective need not imply that the induced map $f_{\text{ptl}} : X_{\text{ptl}} \rightarrow Y_{\text{ptl}}$ is surjective.

Lemma 2.15. *Suppose that f , X , and Y are each strongly definable via the HC-forcing invariant formulas $\theta(x, y)$, $\varphi(x)$, and $\gamma(y)$, respectively and that $f : X \rightarrow Y$ persistently. Then $f : X \rightarrow Y$ is persistently surjective if and only if the formula $\rho(y) := \exists x \theta(x, y)$ is HC-forcing invariant and persistently equivalent to $\gamma(y)$.*

2.3 Potential Cardinality

Definition 2.16. Suppose X and Y are strongly definable. We say that X is *HC-reducible to Y* , written $X \leq_{\text{HC}} Y$, if there is a strongly definable f such that $f : X \rightarrow Y$ is persistently injective. As notation, we write $X <_{\text{HC}} Y$ if $X \leq_{\text{HC}} Y$ but $Y \not\leq_{\text{HC}} X$. We also write $X \sim_{\text{HC}} Y$ if $X \leq_{\text{HC}} Y$ and $Y \leq_{\text{HC}} X$; this is apparently weaker than X and Y being in persistent bijection.

The following notion will be very useful for our applications, as it can often be computed directly. With this and Proposition 2.18 we can prove otherwise difficult non-embeddability results for \leq_{HC} .

Definition 2.17. Suppose X is strongly definable. The *potential cardinality* of X , denoted $\|X\|$, refers to $|X_{\text{ptl}}|$ if X is short, or ∞ otherwise. By convention we say $\kappa < \infty$ for any cardinal κ .

Proposition 2.18. *Suppose X and Y are both strongly definable.*

1. *If Y is short and $X \leq_{\text{HC}} Y$, then X is short.*
2. *If $X \leq_{\text{HC}} Y$, then $\|X\| \leq \|Y\|$.*
3. *If X is short, then $X <_{\text{HC}} \mathcal{P}_{\aleph_1}(X)$.*

Proof. (1) follows immediately from (2).

(2) Choose a strongly definable f such that $f : X \rightarrow Y$ is persistently injective. Then by Proposition 2.14(2), $f_{\text{ptl}} : X_{\text{ptl}} \rightarrow Y_{\text{ptl}}$ is an injective class function. Thus, $|X_{\text{ptl}}| \leq |Y_{\text{ptl}}|$.

(3) Note that for any strongly definable X , $X \leq_{\text{HC}} \mathcal{P}_{\aleph_1}(X)$ is witnessed by the strongly definable map $x \mapsto \{x\}$. For the other direction, suppose by way of contradiction that X is short, but $\mathcal{P}_{\aleph_1}(X) \leq_{\text{HC}} X$. Also, $(\mathcal{P}_{\aleph_1}(X))_{\text{ptl}} = \mathcal{P}(X_{\text{ptl}})$ by Lemma 2.8(3). Thus, by (2), we would have that $|\mathcal{P}(X_{\text{ptl}})| \leq |X_{\text{ptl}}|$, which contradicts Cantor's theorem since X_{ptl} is a set. \square

Using the fact that $(\text{HC}_\beta)_{\text{ptl}} = \mathbb{V}_\beta$, the following Corollary is immediate.

Corollary 2.19. *If X is strongly definable and $\|X\| \leq \beth_\alpha$ for some $\alpha < \omega_1$, then $\text{HC}_{\omega+\alpha+1} \not\leq_{\text{HC}} X$.*

2.4 Quotients

We begin with the obvious definition.

Definition 2.20. A pair (X, E) is a *strongly definable quotient* if both X and E are strongly definable and persistently, E is an equivalence relation on X .

There is an immediate way to define a reduction of two quotients:

Definition 2.21. Let (X, E) and (Y, F) be strongly definable quotients. Say $(X, E) \leq_{\text{HC}} (Y, F)$ if there is a strongly definable f such that all of the following hold persistently:

- f is a subclass of $X \times Y$.
- The E -saturation of $\text{dom}(f)$ is X . That is, for every $x \in X$, there is an $x' \in X$ and $y' \in Y$ where xEx' holds and $(x', y') \in f$.
- f induces a well-defined injection on equivalence classes. That is, if (x, y) and (x', y') are in f , then xEx' holds if and only if yFy' does.

Define $(X, E) <_{\text{HC}} (Y, F)$ and $(X, E) \sim_{\text{HC}} (Y, F)$ in the natural way.

We wish to define the potential cardinality $\|(X, E)\|$. It turns out that $|X_{\text{ptl}}/E_{\text{ptl}}|$ is too small, typically. For our purposes, we can restrict to more well-behaved quotients.

Definition 2.22. A *representation* of a strongly definable quotient (X, E) is a pair f, Z of strongly definable families such that $f : X \rightarrow Z$ is persistently surjective and persistently,

$$\forall a, b \in X [E(a, b) \Leftrightarrow f(a) = f(b)]$$

We say that (X, E) is *representable* if it has a representation.

In the case that (X, E) is representable, the set of E -classes is strongly definable in a sense – we equate it with the representation. For this reason we will also say Z is a representation of (X, E) . Note that if $f_1 : X \rightarrow Z_1$ and $f_2 : X \rightarrow Z_2$ are two representations of (X, E) , then there is a persistently bijective, strongly definable $h : Z_1 \rightarrow Z_2$. This observation implies the following definition is well-defined.

Definition 2.23. If (X, E) is a representable strongly definable quotient, then define $\|(X, E)\| = \|Z\|$ for some (equivalently, for all) Z such that there is a representation $f : (X, E) \rightarrow Z$.

The following lemma can be proved by a routine composition of maps.

Lemma 2.24. *Suppose (X, E) and (Y, E') are strongly definable quotients, with representations $f : (X, E) \rightarrow Z$ and $g : (Y, E') \rightarrow Z'$. Suppose h is a witness to $(X, E) \leq_{\text{HC}} (Y, E')$. Then the induced function $h^* : Z \rightarrow Z'$ is strongly definable, persistently injective, and witnesses $Z \leq_{\text{HC}} Z'$.*

We close this section with an observation about restrictions of representations.

Lemma 2.25. *Suppose $f : (X, E) \rightarrow Z$ is a representation and $Y \subseteq X$ is strongly definable and persistently E -saturated. Let E' and $g = f|_Y$ be the restrictions of E and f , respectively, to Y . Then the image $g(Y)$ is strongly definable, and so $g : (Y, E') \rightarrow g(Y)$ is a representation.*

Proof. We show that $\varphi(z) := \exists y(y \in Y \wedge g(y) = z)$ is HC-forcing invariant. Fix any $z \in \text{HC}$ and let $\mathbb{V}[G]$ be a forcing extension of \mathbb{V} . We show that $z \in \varphi(\text{HC})$ iff $z \in \varphi(\text{HC})^{\mathbb{V}[G]}$. Left to right is clear. For right to left: choose a witness $y \in Y^{\mathbb{V}[G]}$ such that $g(y) = z$ in $\mathbb{V}[G]$. As $f(y) = z$, we conclude $z \in Z^{\mathbb{V}[G]}$ and hence $z \in Z^{\mathbb{V}}$ as Z is strongly definable. So, choose $y^* \in X^{\mathbb{V}}$ with $f(y^*) = z$. Thus, in $\mathbb{V}[G]$, $E(y, y^*)$ holds. As Y is persistently E -saturated, $y^* \in Y^{\mathbb{V}[G]}$. Since Y is strongly definable, we conclude $y^* \in Y^{\mathbb{V}}$, so y^* witnesses that $z \in (g(Y))^{\mathbb{V}}$.

As this argument relativizes to any forcing extension, we conclude that $\varphi(z)$ is HC-forcing invariant. \square

All of the examples we work with will be representable, where the representations are Scott sentences. Therefore this simple definition of $\|(X, E)\|$ will suffice completely for our purposes. In the absence of a representation, one can still define $\|(X, E)\|$ using the notion of *pins*; see for instance [17] for a thorough discussion.

3 Connecting Potential Cardinality with Borel Reducibility

The standard framework for Borel reducibility of invariant classes is the following. Let L be a countable language and let X_L be the set of L -structures with universe ω . Endow X_L with the usual logic topology; then X_L becomes a Polish space. Moreover, if Φ is a sentence of $L_{\omega_1\omega}$ then $\text{Mod}(\Phi)$ is a Borel subset of X_L ; hence $\text{Mod}(\Phi)$ is a standard Borel space. The relation \cong_{Φ} is the restriction of the isomorphism relation to $\text{Mod}(\Phi) \times \text{Mod}(\Phi)$. When no ambiguity arises we omit the Φ . If L' is another countable language and

Φ' is a sentence of $L'_{\omega_1\omega}$, then a Borel reduction from $(\text{Mod}(\Phi), \cong) \rightarrow (\text{Mod}(\Phi'), \cong)$ is a Borel map $f : \text{Mod}(\Phi) \rightarrow \text{Mod}(\Phi')$ such that, for all $M, N \in \text{Mod}(\Phi)$, $M \cong N$ if and only if $f(M) \cong f(N)$.

We want to apply the machinery of the previous section to this setup. First, recall that we are working entirely in *ZFC*; thus a language L is just a set with an arity function, and an L -structure with universe ω is just a function $f : L \rightarrow \bigcup_n \mathcal{P}(\omega^n)$ respecting the arities. Since our languages are countable we can suppose that they are elements of *HC*. We will presently show that for any sentence Φ of $L_{\omega_1\omega}$, $(\text{Mod}(\Phi), \cong)$ is a strongly definable quotient. We will also show that $(\text{Mod}(\Phi), \cong)$ is representable, and that Borel reductions are in particular HC-reductions.

3.1 Canonical Scott sentences

For what follows, we need the notion of a canonical Scott sentence of any infinite L -structure, regardless of cardinality. The definition below is in both Barwise [1] and Marker [16].

Definition 3.1. Suppose L is countable and M is any infinite L -structure, say of power κ . For each $\alpha < \kappa^+$, define an $L_{\kappa^+, \omega}$ formula $\varphi_\alpha^{\bar{a}}(\bar{x})$ for each finite $\bar{a} \in M^{<\omega}$ as follows:

- $\varphi_0^{\bar{a}}(\bar{x}) := \bigwedge \{\theta(\bar{x}) : \theta \text{ atomic or negated atomic and } M \models \theta(\bar{a})\}$;
- $\varphi_{\alpha+1}^{\bar{a}}(\bar{x}) := \varphi_\alpha^{\bar{a}}(\bar{x}) \wedge \bigwedge \{\exists y \varphi_\alpha^{\bar{a}, b}(\bar{x}, y) : b \in M\} \wedge \forall y \bigvee \{\varphi_\alpha^{\bar{a}, b}(\bar{x}, y) : b \in M\}$;
- For α a non-zero limit, $\varphi_\alpha^{\bar{a}}(\bar{x}) := \bigwedge \{\varphi_\beta^{\bar{a}}(\bar{x}) : \beta < \alpha\}$.

Next, let $\alpha^*(M) < \kappa^+$ be least ordinal α such that for all finite \bar{a} from M ,

$$\forall \bar{x} [\varphi_\alpha^{\bar{a}}(\bar{x}) \rightarrow \varphi_{\alpha+1}^{\bar{a}}(\bar{x})].$$

Finally, put $\text{css}(M) := \varphi_{\alpha^*(M)}^{\emptyset} \wedge \bigwedge \{\forall \bar{x} [\varphi_{\alpha^*(M)}^{\bar{a}}(\bar{x}) \rightarrow \varphi_{\alpha^*(M)+1}^{\bar{a}}(\bar{x})] : \bar{a} \in M^{<\omega}\}$.

For what follows, it is crucial that the choice of $\text{css}(M)$ really is canonical. In particular, in the infinitary clauses forming the definition of $\text{tp}_{\alpha+1}^{\bar{a}}(\bar{x})$, we consider the conjunctions and disjunctions be taken over *sets* of formulas, as opposed to sequences. In particular, we ignore the multiplicity of a formula inside the set. By our conventions about working wholly in *ZFC*, countable languages and sentences of $L_{\infty, \omega}$ are sets, and in particular $\text{css}(M)$ is a set.

We summarize the well-known, classical facts about canonical Scott sentences with the following:

Fact 3.2. *Fix a countable language L .*

1. *For every L -structure M , $M \models \text{css}(M)$; and for all L -structures N , $M \equiv_{\infty, \omega} N$ if and only if $\text{css}(M) = \text{css}(N)$ if and only if $N \models \text{css}(M)$.*
2. *If M is countable, then $\text{css}(M) \in \text{HC}$.*
3. *css is absolute between transitive models of ZFC^- , where ZFC^- is ZFC but with the powerset axiom is deleted. (Recall that $\text{HC} \models \text{ZFC}^-$.)*

4. If M and N are both countable, then $M \cong N$ if and only if $\text{css}(M) = \text{css}(N)$ if and only if $N \models \text{css}(M)$.

Our primary interest in canonical Scott sentences is that they give rise to representations of classes of L -structures. The key to the representability is Karp's Completeness Theorem for sentences of $L_{\omega_1, \omega}$, see e.g., Theorem 3 of Keisler [12], which says that if a sentence σ of $L_{\omega_1, \omega}$ is consistent, then it has a countable model. It quickly follows that if σ is a sentence of $L_{\omega_1, \omega}$, and σ has a model in a forcing extension, then σ already has a countable model in \mathbb{V} .

We begin by considering $\text{CSS}(L)$, the set of all canonical Scott sentences of structures in X_L , the set of L -structures with universe ω .

Lemma 3.3. *Fix a countable language L . Then:*

1. $\text{CSS}(L)$ is strongly definable via the formula $\varphi(y) := \exists M(M \in X_L \wedge \text{css}(M) = y)$;
2. The strongly definable function $\text{css} : X_L \rightarrow \text{CSS}(L)$ is persistently surjective;
3. $\text{css} : X_L \rightarrow \text{CSS}(L)$ is a representation of the strongly definable quotient (X_L, \cong) .

Proof. Note that $L_{\omega_1, \omega}$ is strongly definable, clearly.

(1) We need to verify that $\varphi(y)$ is HC-forcing invariant. Suppose $\sigma \in \text{HC}$ and $\mathbb{V}[G]$ is a forcing extension of \mathbb{V} . If $\sigma \in \varphi(\text{HC})^{\mathbb{V}[G]}$, then there is some $M \in X_L^{\mathbb{V}[G]}$ such that $\text{css}(M) = \sigma$. Hence by the preceding discussion, σ has a countable model $N \in (X_L)^{\mathbb{V}}$. In $\mathbb{V}[G]$, $N \models \sigma$; but $\sigma = \text{css}(M)$. So $\text{css}(N) = \sigma$. As this argument readily relativizes to any forcing extension, φ is HC-forcing invariant.

(2) follows from (1) and Lemma 2.15.

(3): (X_L, \cong) is a strongly definable quotient since css is strongly definable. Thus we conclude by (2) and Fact 3.2(4). \square

In most of our applications, we are interested in strongly definable subclasses of X_L that are closed under isomorphism. For any sentence Φ of $L_{\omega_1, \omega}$, because $\text{Mod}(\Phi)$ is a Borel subset of X_L , it follows from Shoenfield's Absoluteness Theorem that both $\text{Mod}(\Phi)$ and the restriction of css to $\text{Mod}(\Phi)$ (also denoted by css) $\text{css} : \text{Mod}(\Phi) \rightarrow \text{HC}$ are strongly definable.

Definition 3.4. For Φ any sentence of $L_{\omega_1, \omega}$, $\text{CSS}(\Phi) = \{\text{css}(M) : M \in \text{Mod}(\Phi)\} \subseteq \text{HC}$.

Proposition 3.5. *Fix any sentence $\Phi \in L_{\omega_1, \omega}$ in a countable vocabulary. Then $\text{css} : \text{Mod}(\Phi) \rightarrow \text{CSS}(\Phi)$ is a representation of the quotient $(\text{Mod}(\Phi), \cong)$. In particular the latter is strongly definable.*

Proof. As $\text{Mod}(\Phi)$ is strongly definable, this follows immediately from Lemmas 3.3 and 2.25. \square

Definition 3.6. Let Φ be any sentence of $L_{\omega_1, \omega}$ in a countable vocabulary. We say that Φ is *short* if $\text{CSS}(\Phi)_{\text{ptl}}$ is a set (as opposed to a proper class). If Φ is short, let the *potential cardinality* of Φ , denoted $\|\Phi\|$, be the (usual) cardinality of $\text{CSS}(\Phi)_{\text{ptl}}$; otherwise let it be ∞ .

It follows from Proposition 3.5 and Definitions 2.16 and 2.23 that

$$\|\Phi\| = \|(\text{Mod}(\Phi), \cong)\| = \|\text{CSS}(\Phi)\| = |\text{CSS}(\Phi)_{\text{ptl}}|.$$

In order to understand the class $\text{CSS}(\Phi)_{\text{ptl}}$, note that if $\varphi \in \text{CSS}(\Phi)_{\text{ptl}}$, then $\varphi \in \mathbb{V}$ and is a sentence of $L_{\infty, \omega}$. If we choose any forcing extension $\mathbb{V}[G]$ of \mathbb{V} for which $\varphi \in \text{HC}^{\mathbb{V}[G]}$, then $\mathbb{V}[G] \models \text{'}\varphi \in L_{\omega_1, \omega}\text{'}$ and there is some $M \in \text{HC}^{\mathbb{V}[G]}$ such that $V[G] \models \text{'}\varphi \in \text{Mod}(\Phi) \text{ and } \text{css}(M) = \varphi\text{'}$. Thus, we refer to elements of $\text{CSS}(\Phi)_{\text{ptl}}$ as being *potential canonical Scott sentences* of a model of Φ . In particular, every element of $\text{CSS}(\Phi)_{\text{ptl}}$ is *potentially satisfiable* in the sense that it is satisfiable in some forcing extension $\mathbb{V}[G]$ of \mathbb{V} . There is a proof system for sentences of $L_{\infty, \omega}$ for which a sentence is consistent if and only if it is potentially satisfiable as defined above.² When we say ' φ implies ψ ', we mean with respect to this proof system; equivalently, in any forcing extension $\mathbb{V}[G]$, every model of φ is a model of ψ .

We can also ask what is the image of the class function css when restricted to the class of models of Φ . As notation, let $\text{CSS}(\Phi)_{\text{sat}}$ denote the class $\{\text{css}(M) : M \in \mathbb{V} \text{ and } M \models \Phi\}$. This choice of notation is clarified by the following easy lemma.

Lemma 3.7. $\text{CSS}(\Phi)_{\text{sat}} \subseteq \text{CSS}(\Phi)_{\text{ptl}}$.

Proof. Choose any $\varphi \in \text{CSS}(\Phi)_{\text{sat}}$ and choose any $M \in \mathbb{V}$ such that $M \models \Phi$ and $\text{css}(M) = \varphi$. Then clearly, $\varphi \in \mathbb{V}$. As well, choose a forcing extension $\mathbb{V}[G]$ in which M is countable. Then, in $\mathbb{V}[G]$, there is some $M' \in \text{Mod}(\Phi)$ (i.e., where the universe of M' is ω) such that $M' \cong M$. Then $(\text{css}(M'))^{\mathbb{V}[G]} = \varphi$ and so $\varphi \in \text{CSS}(\Phi)_{\text{ptl}}$. \square

To summarize, elements of $\text{CSS}(\Phi)_{\text{ptl}}$ are called *potential canonical Scott sentences*, whereas elements of $\text{CSS}(\Phi)_{\text{sat}}$ are *satisfiable*. This suggests a property of the sentence Φ .

Definition 3.8. A sentence $\Phi \in L_{\omega_1, \omega}$ (or a complete first-order theory T) is *grounded* if $\text{CSS}(\Phi)_{\text{sat}} = \text{CSS}(\Phi)_{\text{ptl}}$, i.e., if every potential canonical Scott sentence is satisfiable.

As a trivial example, if T is \aleph_0 -categorical, then as all models of T are back-and-forth equivalent, $\text{CSS}(T)_{\text{ptl}}$ is a singleton, hence T is grounded. In Section 5 we show that both of the theories $\text{REF}(\text{bin})$ and $\text{REF}(\text{inf})$ are grounded, but in Section 6 we prove that the theory TK is not grounded.

Next, we show that a Borel reduction between invariant classes yields a strongly definable map between the associated canonical Scott sentences.

Fact 3.9. *Suppose Φ and Φ' are sentences of $L_{\omega_1, \omega}$ and $L'_{\omega_1, \omega}$ respectively. If there is a Borel reduction $f : (\text{Mod}(\Phi), \cong) \rightarrow (\text{Mod}(\Phi'), \cong)$ then there is a strongly definable $f^* : \text{CSS}(\Phi) \rightarrow \text{CSS}(\Phi')$ between canonical Scott sentences that is persistently injective. Hence $\text{CSS}(\Phi) \leq_{\text{HC}} \text{CSS}(\Phi')$.*

²In Chapter 4 of [12], Keisler gives a proof system for $L_{\omega_1, \omega}$, and shows in Theorem 3 that it is complete, i.e. if φ is unprovable then $\neg\varphi$ has a model. The natural generalization of this proof system to $L_{\infty, \omega}$ works: the proof of Theorem 3 shows that whenever φ is unprovable, then $\neg\varphi$ lies in a consistency property. Forcing on the consistency property gives a model of $\neg\varphi$.

Proof. It is a standard theorem, see e.g., [11] Proposition 12.4, that the graph of f is Borel. So f is naturally a strongly definable family of Borel sets. By Lemma 2.24, it suffices to show that $f : (\text{Mod}(\Phi), \cong) \leq_{\text{HC}} (\text{Mod}(\Phi'), \cong)$, which amounts to showing that f remains well-defined and injective on isomorphism classes in every forcing extension. But this is a Π_2^1 statement in codes for f, Φ, Φ' , and thus is absolute to forcing extensions by Shoenfield's Absoluteness Theorem. \square

The following Theorem is simply a distillation of our previous results.

Theorem 3.10. *Let Φ and Ψ be sentences of $L_{\omega_1, \omega}$, possibly in different countable vocabularies.*

1. *If Ψ is short, while Φ is not short, then $(\text{Mod}(\Phi), \cong)$ is not Borel reducible to $(\text{Mod}(\Psi), \cong)$.*
2. *If $\|\Psi\| < \|\Phi\|$, then $(\text{Mod}(\Phi), \cong)$ is not Borel reducible to $(\text{Mod}(\Psi), \cong)$.*

Proof. (1) follows immediately from (2).

(2) Suppose $f : (\text{Mod}(\Phi), \cong) \rightarrow (\text{Mod}(\Psi), \cong)$ were a Borel reduction. Then by Fact 3.9 we would obtain a strongly definable $f^* : \text{CSS}(\Phi) \rightarrow \text{CSS}(\Psi)$ that is persistently injective, meaning $\|\text{CSS}(\Phi)\| \leq \|\text{CSS}(\Psi)\|$. This contradicts Proposition 2.18(2). \square

3.2 Consequences of \cong_{Φ} being Borel

Although our primary interest is classes $\text{Mod}(\Phi)$ where \cong_{Φ} is not Borel, in this brief subsection we see the consequences of \cong_{Φ} being Borel.

Theorem 3.11. *The following are equivalent for a sentence $\Phi \in L_{\omega_1, \omega}$ in countable vocabulary.*

1. *The relation of \cong on $\text{Mod}(\Phi)$ is a Borel subset of $\text{Mod}(\Phi) \times \text{Mod}(\Phi)$;*
2. *For some $\alpha < \omega_1$, $\text{CSS}(\Phi) \subseteq \text{HC}_{\alpha}$;*
3. *For some $\alpha < \omega_1$, $\text{CSS}(\Phi)$ is persistently contained in HC_{α} ;*
4. *$\text{CSS}(\Phi)_{\text{ptl}}$ is contained in ∇_{α} for some $\alpha < \omega_1$.*

Proof. To see the equivalence of (1) and (2), first note that in both conditions we are only considering models of Φ with universe ω and the canonical Scott sentence of such objects. In particular, neither condition involves passing to a forcing extension. However, it is a classical result (see for instance [4], Theorem 12.2.4) that \cong is Borel if and only if the Scott ranks of countable models are bounded below ω_1 , which is equivalent to stating that there is a bound on the canonical Scott sentences in the HC_{α} hierarchy.

For (2) implies (3), note that the formula $\exists M : M \models \Phi \wedge \text{css}(M) \notin \text{HC}_{\alpha}$ is a Σ_1 formula in the parameters $\Phi, \alpha \in \text{HC}$ and so is absolute to forcing extensions by Lemma 2.2.

That (3) implies (4) and (4) implies (2) follow directly from Example 2.11(6). \square

We obtain an immediate corollary to this. Let $I_{\infty,\omega}(\Phi)$ denote the cardinality of a maximal set of pairwise $\equiv_{\infty,\omega}$ -inequivalent models $M \in \mathbb{V}$ (of any cardinality) with $M \models \Phi$. If no maximal set exists, we write $I_{\infty,\omega}(\Phi) = \infty$. By Fact 3.2 and Lemma 3.7, $I_{\infty,\omega}(\Phi) = |\text{CSS}(\Phi)_{\text{sat}}| \leq \|\Phi\|$.

Corollary 3.12. *Let Φ be any sentence in $L_{\omega_1,\omega}$ in a countable vocabulary such that \cong is a Borel subset of $\text{Mod}(\Phi) \times \text{Mod}(\Phi)$. Then*

1. Φ is short; and
2. $I_{\infty,\omega}(\Phi) < \beth_{\omega_1}$. (In fact $\|\Phi\| < \beth_{\omega_1}$.)

Proof. Assume that \cong is a Borel subset of $\text{Mod}(\Phi) \times \text{Mod}(\Phi)$. By Theorem 3.11(3), $\text{CSS}(\Phi)_{\text{ptl}} \subseteq \mathbb{V}_\alpha$ for some $\alpha < \omega_1$ and hence is a set. Thus, Φ is short, and $I_{\infty,\omega}(\Phi) = |\text{CSS}(\Phi)_{\text{sat}}| \leq |\text{CSS}(\Phi)_{\text{ptl}}| \leq |\mathbb{V}_\alpha| < \beth_{\omega_1}$. \square

We remark that the implication in Corollary 3.12 does not reverse. In Sections 5 and 6 we show that both of the complete theories $\text{REF}(\text{bin})$ and K are short, but on countable models of either theory, \cong is not Borel.

3.3 Maximal Complexity

In this subsection, we recall two definitions of maximality. The first, Borel completeness, is from Friedman-Stanley [3].

Definition 3.13. Let $\Phi \in L_{\omega_1,\omega}$. The quotient $(\text{Mod}(\Phi), \cong)$ is *Borel complete* if every $(\text{Mod}(\Psi), \cong)$ is Borel reducible to $(\text{Mod}(\Phi), \cong)$.

Corollary 3.14. *If $(\text{Mod}(\Phi), \cong)$ is Borel complete, then Φ is not short.*

Proof. Let $L = \{\leq\}$ and let Ψ assert that \leq is a linear ordering. As $(\text{Mod}(\Phi), \cong)$ is Borel complete, there is a Borel reduction $f : (\text{Mod}(\Psi), \cong) \rightarrow (\text{Mod}(\Phi), \cong)$. However, it is easily proved that for distinct ordinals $\alpha \neq \beta$, the L -structures (α, \leq) and (β, \leq) are $\equiv_{\infty,\omega}$ -inequivalent models of Ψ , hence have distinct canonical Scott sentences. Thus, $\text{CSS}(\Psi)_{\text{sat}}$ is a proper class, and hence so is $\text{CSS}(\Psi)_{\text{ptl}}$ by Lemma 3.7. So Φ cannot be short by Theorem 3.10(1). \square

If one is only interested in classes of countable models, then the Borel complete classes are clearly maximal with respect to Borel reducibility. As any invariant class of countable structures has a natural extension to a class of uncountable structures, one can ask for more. The following definitions from [15] generalize Borel completeness to larger cardinals λ . To see that it is a generalization, recall that among countable structures, isomorphism is equivalent to back-and-forth equivalence, and that for structures of size λ , $\equiv_{\lambda^+,\omega}$ -equivalence is also equivalent to back-and-forth equivalence. Consequently, ‘Borel complete’ in the sense of Definition 3.13 is equivalent to ‘ \aleph_0 -Borel complete’ in Definition 3.15. So, ‘ Φ is λ -Borel complete for all infinite λ ’ implies Φ Borel complete. However, in Section 6 we will see that the theory TK is Borel complete, but is not λ -Borel complete for large λ .

Definition 3.15. Let Φ be a sentence of $L_{\omega_1, \omega}$.

- For $\lambda \geq \aleph_0$, let $\text{Mod}_\lambda(\Phi)$ denote the class of models of Φ with universe λ .
- Topologize $\text{Mod}_\lambda(\Phi)$ by declaring that $\mathcal{B} := \{U_{\theta(\bar{\alpha})} : \theta(\bar{x}) \text{ is quantifier free and } \bar{\alpha} \in \lambda^{<\omega}\}$ is a sub-basis, where $U_{\theta(\bar{\alpha})} = \{M \in \text{Mod}_\lambda(\Phi) : M \models \theta(\bar{\alpha})\}$.
- A set is λ -Borel if it is in the λ^+ -algebra generated by the sub-basis \mathcal{B} .
- A function $f : \text{Mod}_\lambda(\Phi) \rightarrow \text{Mod}_\lambda(\Psi)$ is a λ -Borel embedding if
 - the inverse image of every (sub)-basic open set is λ -Borel; and
 - For $M, N \in \text{Mod}_\lambda(\Phi)$, $M \equiv_{\infty, \omega} N$ if and only if $f(M) \equiv_{\infty, \omega} f(N)$.
- $(\text{Mod}_\lambda(\Phi), \equiv_{\infty, \omega})$ is λ -Borel reducible to $(\text{Mod}_\lambda(\Psi), \equiv_{\infty, \omega})$ if there exists a λ -Borel embedding $f : \text{Mod}_\lambda(\Phi) \rightarrow \text{Mod}_\lambda(\Psi)$.
- Φ is λ -Borel complete if every $(\text{Mod}_\lambda(\Psi), \equiv_{\infty, \omega})$ is λ -Borel reducible to $(\text{Mod}_\lambda(\Phi), \equiv_{\infty, \omega})$.

For example, the class of graphs (directed or undirected) is λ -Borel complete for all infinite λ . This is a standard coding argument. Although we are not aware of any direct reference, Theorem 5.5.1 of [7] states that graphs can interpret any theory. It is easily checked that the map constructed in the proof of Theorem 5.5.1 is in fact a λ -Borel reduction for every λ .

Also, in [15] it is proved that the class of subtrees of $\lambda^{<\omega}$ is λ -Borel complete, and more recently the second author has proved that the class of linear orders is λ -Borel complete for all λ .

3.4 Jumps and products

In this subsection we recall two procedures – the *jump* and the *product* – and use them to define a sequence $\langle T_\alpha : \alpha \in \omega_1 \rangle$ of complete, first order theories for which the potential cardinality is strictly increasing.

Definition 3.16. Suppose L is a countable relational language and $\Phi \in L_{\omega_1, \omega}$. The *jump of Φ* , written $J(\Phi)$, is a sentence of $L'_{\omega_1, \omega}$ defined as follows, where $L' = L \cup \{E\}$ is obtained by adding a new binary relation symbol E to L . Namely $J(\Phi)$ states that E is an equivalence relation with infinitely many classes, each of which is a model of Φ . If $R \in L$ and \bar{x} is a tuple not all from the same E -class, then $R(\bar{x})$ is defined to be false, so that the models are independent.

There is a corresponding notion of jump that can be defined directly on equivalence relations: Given an equivalence relation E on X , its jump is the equivalence relation $J(E)$ on X^ω , defined by setting $(x_n : n \in \omega) J(E) (y_n : n \in \omega)$ if there is some $\sigma \in S_\infty$ with $x_{\sigma(n)} E y_n$ for all $n \in \omega$. Then the previous definition of the jump can be viewed as the special case where (X, E) is $(\text{Mod}(\Phi), \cong)$.

The notion of a jump was investigated in [3], where it was shown that if E is a Borel equivalence relation on a Polish space X with more than one class, then $E <_B J(E)$. We give a partial generalization of this in

Proposition 3.17(3) – if $\Phi \in L_{\omega_1, \omega}$ is short, then $\|\Phi\| < \|J(\Phi)\|$, so $\Phi <_B J(\Phi)$. Using the theory of pins [17], one can use essentially the same proof to give a true generalization: if (X, E) is strongly definable, short, and has more than one E -class, then $\|(X, E)\| < \|(X^\omega, J(E))\|$, so in particular $E <_B J(E)$.

The following Proposition lists the basic properties of the jump operation.

Proposition 3.17. *Let Φ and Ψ be $L_{\omega_1, \omega}$ -sentences in countable relational languages.*

1. *If Φ is a complete first order theory, so is $J(\Phi)$.*
2. *If Φ is grounded, so is $J(\Phi)$.*
3. *If Φ is short, then $J(\Phi)$ is also short. More precisely, if $\|\Phi\|$ is infinite, $\|J(\Phi)\| = 2^{\|\Phi\|}$. If $2 \leq \|\Phi\| < \aleph_0$, $\|J(\Phi)\| = \aleph_0$. If $\|\Phi\| = 1$, then $\|J(\Phi)\| = 1$.*
4. *The jump is monotone: if $\Phi \leq_B \Psi$, then $J(\Phi) \leq_B J(\Psi)$.*
5. *It is always true that $\Phi \leq_B J(\Phi)$; if Φ is short and not \aleph_0 -categorical, then $\Phi <_B J(\Phi)$.*

Note that here and throughout, we use $\Phi \leq_B \Psi$ as a shorthand for $(\text{Mod}(\Phi), \cong) \leq_B (\text{Mod}(\Psi), \cong)$, and similarly with $<_B$ and \sim_B .

Proof. (1) That the jump is first-order is clear. Completeness follows from a standard Ehrenfeucht-Fraïssé argument.

(2) Let $\Psi \in \text{CSS}(J(\Phi))_{\text{ptl}}$, and let $\mathbb{V}[G]$ be some forcing extension in which Ψ is hereditarily countable. Let $M \models \Psi$ be the unique countable model of Ψ in $\mathbb{V}[G]$. Let X be the set of E -classes in M , and for each $x \in X$, let Ψ_x be the canonical Scott sentence of x , viewed as an L -structure. Let m_x be the number of E -classes of M which are isomorphic to x as L -structures, if this number is finite; if infinite, let $m_x = \omega$.

The set of pairs $S = \{(\Psi_x, m_x) : x \in X\}$ depends only on the isomorphism type of M , so is uniquely definable from Ψ . By Lemma 2.5, $S \in \mathbb{V}$, although it may no longer be countable. Since Φ is grounded, each Ψ_x has a model in \mathbb{V} , so let N be the model coded from S – for each pair (Ψ', m) in S , give m distinct E -classes, each of which is a model of Ψ' . Then $M \cong N$ in any sufficiently large forcing extension of $\mathbb{V}[G]$, so $N \models \Psi$, as desired.

(3) We assume $\|\Phi\|$ is infinite; the finite cases are similar and trivial, respectively. First, let $X \subset \text{CSS}(\Phi)_{\text{ptl}}$ be arbitrary. Let $\mathbb{V}[G]$ be a forcing extension in which X is hereditarily countable, and let M_X be a countable model of $J(\Phi)$ such that each E -class of M_X is a model of some $\Psi \in X$, and each $\Psi \in X$ is represented infinitely often. Since each Ψ is \aleph_0 -categorical, M_X is determined up to isomorphism by these constraints. Therefore, $\Psi_X = \text{css}(M_X)$ is determined entirely by X , so $\Psi_X \in \mathbb{V}$ by Lemma 2.5. Furthermore, if $X \neq Y$, then $M_X \not\cong M_Y$, so $\Psi_X \neq \Psi_Y$, so $\|J(\Phi)\| \geq 2^{\|\Phi\|}$.

The other direction is similar to part (2). Let $\Psi \in \text{CSS}(J(\Phi))_{\text{ptl}}$, let $\mathbb{V}[G]$ be a forcing extension in which Ψ is hereditarily countable, and let $M_\Psi \models \Psi$ be the unique countable model. Let X be the set of E -classes of M_Ψ , and for each $x \in X$, let Ψ_x be $\text{css}(x)$, where we consider the E -class x as an L -structure and a

model of Ψ . For each x , let $m(x)$ be the number of equivalence classes of M which are isomorphic to x as \mathcal{L} -structures, or ω if there are infinitely many. Since M_Ψ is determined up to isomorphism by Ψ , the set $S_\Psi = \{(\Psi_x, m(x)) : x \in X\}$ is determined entirely by Ψ . Therefore, $S_\Psi \in \mathbb{V}$ by Lemma 2.5. If $\Psi \neq \Psi'$, $M_\Psi \not\cong M_{\Psi'}$, so $S_\Psi \neq S_{\Psi'}$. Thus we see that $\|J(\Phi)\| \leq |(\omega + 1)^{\|\Phi\|}|$. Since $\|\Phi\|$ is infinite, this last is equal to $2^{\|\Phi\|}$, completing the proof.

(4) follows from the fact that each equivalence class is in canonical bijection with ω , allowing us to apply f to each class and reindex. (5) follows from (3) and the fact that if $\|\Phi\| > \|\Psi\|$ then $\Phi \not\leq_B \Psi$. \square

Another important operation is the *product*:

Definition 3.18. Suppose I is a countable set and for each i , Φ_i is a sentence of $L_{\omega_1, \omega}$ in the countable relational language L_i . The *product of the Φ_i* , denoted $\prod_i \Phi_i$, is a sentence of $L_{\omega_1, \omega}$, where $L = \{U_i : i \in I\} \cup \bigcup_i L_i$ is the disjoint union of the L_i 's together with new unary predicates $\{U_i : i \in I\}$.

Namely $\prod_i \Phi_i$ states that the U_i are disjoint, that the elements of U_i form a model of Φ_i when viewed as an L_i -structure, and that if $R \in L_i$ and \bar{x} is not all from U_i , then $R(\bar{x})$ is false, so that the models are independent. If I is finite, we also require that each element is in some U_i . This becomes more convenient

The proofs of the corresponding facts for products are quite similar to those for the jump, so we omit them:

Proposition 3.19. *Let $\{\Phi_i : i \in I\}$ and $\{\Psi_j : j \in J\}$ be countable sets of $L_{\omega_1, \omega}$ -sentences in countable relational languages.*

1. *If each Φ_i is complete and first order, so is $\prod_i \Phi_i$.*
2. *If each Φ_i is grounded, so is $\prod_i \Phi_i$.*
3. *If each Φ_i is short then $\prod_i \Phi_i$ is short. More precisely, $\|\prod_i \Phi_i\| = \kappa \cdot \prod_i \|\Phi_i\|$, where κ is 1 if I is finite or \aleph_0 if I is infinite.*
4. *The product is monotone: if $f : I \rightarrow J$ is an injection and $\Phi_i \leq_B \Psi_{f(i)}$ for all $i \in I$, then $\prod_i \Phi_i \leq_B \prod_j \Psi_j$.*
5. *It is always true that for all $i \in I$, $\Phi_i \leq_B \prod_i \Phi_i$. If additionally, for all $i \in I$, there is an $i' \in I$ where $\Phi_i <_B \Phi_{i'}$, then for all $i \in I$, $\Phi_i <_B \prod_i \Phi_i$.*

Note that if we use $L_{\omega_1, \omega}$ to add to our definition of product that the U_i are exhaustive, the product of first-order theories may not be first order, but the statement of (3) improves and we can lose reference to κ . This can also be achieved by working in a multisorted first-order logic, which preserves (1) and fixes (3). This is the approach we prefer, but for the time being we remain flexible.

We close this section by defining a set of concrete benchmarks: Note that these are essentially the same as the \mathcal{I}_α in [3], the \cong_α in [6], the $=^\alpha$ in [4], and the T_α in [13].

Definition 3.20. T_0 is the theory of (\mathbb{Z}, S) , where S is the graph of the successor function.

For each countable α , $T_{\alpha+1} = J(T_\alpha)$. If α is a limit ordinal, $T_\alpha = \prod_{\beta < \alpha} T_\beta$.

We quickly summarize the properties which are relevant to us. First, they all have Borel isomorphism relation; that is, \cong_{T_α} is Borel for every α . (This is because having Borel isomorphism relation is preserved under jumps and products, as can be easily checked.) Second, they are cofinal among such theories: if \cong_Φ is Borel, then $\Phi \leq_B T_\alpha$ for some α , see [4] Corollary 12.2.8. Therefore, characterizing the relationship between some Φ and the T_α is a reasonable way to gauge the complexity of $(\text{Mod}(\Phi), \cong)$. Finally:

Corollary 3.21. *Each T_α is a complete first-order theory which is short and grounded. For finite α , $\|T_\alpha\| = \beth_\alpha$; for infinite α , $\|T_\alpha\| = \beth_{\alpha+1}$. In particular, $T_\alpha <_B T_\beta$ for all $\alpha < \beta < \omega_1$.*

Proof. We first check $\alpha = 0$. Clearly T_0 has exactly \aleph_0 models up to back-and-forth equivalence. If $M \models T_0$ in some $\mathbb{V}[G]$, then either M has finite dimension n , so is isomorphic to a model in \mathbb{V} , or M has infinite dimension, so is back-and-forth equivalent to the countable model $(\omega \times \mathbb{Z}, S)$ in \mathbb{V} . Either way, this simultaneously shows that T_0 is grounded and $\|T_0\| = \aleph_0$, completing the proof.

The rest of the proof goes immediately by induction, using Propositions 3.17 and 3.19. Groundedness follows from part (2). Size counting (and therefore shortness) follows from part (3), which is immediate at successor stages but the limits require an argument. So let α be a limit ordinal. Then $\|T_\alpha\| = \prod_\beta \|T_\beta\|$ (regardless of whether there are unsorted elements). By the inductive hypothesis and the rules of cardinal multiplication, this is equal to $\prod_\beta \beth_\beta = \prod_\beta 2^{\beth_\beta}$, which in turn is equal to $2^{(\sum_\beta \beth_\beta)} = 2^{\beth_\alpha} = \beth_{\alpha+1}$, completing the proof.

The strictness of the ascending chain follows from induction and part (5). □

4 Compact group actions

In this brief section we use the technology of canonical Scott sentences and representability to analyze the effect of a continuous action of a compact group on a Polish space X . In particular, we show that the quotient of $\mathcal{P}_{\aleph_1}(X)$ by the diagonal action of G is representable. We also show that if the group is abelian, we can bound the potential cardinality of the representation. In Section 6 we use these results to analyze the models of the theory K and to contrast K with TK .

Suppose we have a Polish group G acting on a Polish space X . To apply our machinery to this situation we need to say what it means for the objects involved to be strongly definable:

- Definition 4.1.**
- A strongly definable Polish space is a sequence (X, d, D, i) of strongly definable families, where persistently: d is a complete metric on X , $D \subset X$ is dense and $i : \omega \rightarrow D$ is a bijection.
 - A strongly definable Polish group is a sequence (G, d', D', i', \times) where (G, d', D', i') is a strongly definable Polish space and persistently, \times is a compatible group operation on G .

- Suppose (G, d', D', i', \times) is a strongly definable Polish group, (X, d, D, i) is a strongly definable Polish space. Then a strongly definable continuous action of G on X is a strongly definable family \cdot such that persistently, $\cdot \subset G \times X \times X$ is a continuous action of G on X .

Throughout this subsection, we fix a strongly definable Polish space (X, d, D, i) , a strongly definable, persistently compact Polish group (G, d', D', i', \times) , and a strongly definable continuous action \cdot of G on X .

We also fix strongly definable families

$$\mathcal{B}_n = \{U_i^n : i \in \omega\}$$

such that persistently, each \mathcal{B}_n is a basis for the topology on X^n . (For instance, take \mathcal{B}_1 to be the balls with rational radius and center in D , using the enumeration of D given by i .)

The action of G on X naturally gives diagonal actions on both X^n and $\mathcal{P}(X)$ defined by $g \cdot \bar{a} = \langle g \cdot a : a \in \bar{a} \rangle$ and $g \cdot A = \{g \cdot a : a \in A\}$, respectively. Clearly, the diagonal action of G takes countable subsets of X to countable subsets. For all of these spaces, let \sim^G be the equivalence relation induced by G .

In order to understand the quotient $(\mathcal{P}_{\aleph_1}(X), \sim^G)$, we begin with one easy lemma that uses the fact that G is compact. This lemma is the motivation for the language we define below.

Lemma 4.2. *If $A, B \in \mathcal{P}_{\aleph_1}(X)$, then $A \sim^G B$ if and only if there is a bijection $\sigma : A \rightarrow B$ satisfying $\bar{a} \sim^G h(\bar{a})$ for all $\bar{a} \in A^{<\omega}$.*

Proof. If $g \cdot A = B$, then $\sigma := g \upharpoonright_A$ is as desired. For the converse, fix such a σ ; we will show there is $g \in G$ inducing σ . Let $\{a_n : n \in \omega\}$ be an enumeration of A , and for each n , let \bar{a}_n be the tuple $a_0 \dots a_{n-1}$ and let $C_n \subseteq G$ be the set of all $g \in G$ with $g \cdot \bar{a}_n = \sigma(\bar{a}_n)$. C_n is closed since the action is continuous and C_n is nonempty by hypothesis. Since G is compact, $C = \bigcap_n C_n$ is nonempty, and clearly any $g \in C$ has $g \cdot A = B$. \square

We define a language L and a class of L -structures that encode this information. Put $L := \{R_i^n : i \in \omega, n \geq 1\}$, where each R_i^n is an n -ary relation. Let M_X be the L -structure with universe all of X , with each R_i^n interpreted by

$$M_X \models R_i^n(\bar{a}) \quad \text{if and only if} \quad G \cdot \bar{a} \cap U_i^n = \emptyset$$

As notation, let $\text{qf}_n(\bar{a})$ denote the quantifier-free type of $\bar{a} \in X^n$. It is easily seen that to specify $\text{qf}_n(\bar{a})$ it is enough to specify the set of $i \in \omega$ such that $M_X \models R_i^n(\bar{a})$. Also,

$$\text{qf}_n(\bar{a}) = \text{qf}_n(\bar{b}) \quad \text{if and only if} \quad G \cdot \bar{a} = G \cdot \bar{b}$$

As well, note that every $g \in G$ induces an L -automorphism of M_X given by $a \mapsto g \cdot a$. These two observations imply that M_X has a certain homogeneity – For $\bar{a}, \bar{b} \in X^n$, $\text{qf}_n(\bar{a}) = \text{qf}_n(\bar{b})$ if and only if there is an automorphism of M_X taking \bar{a} to \bar{b} .

For $\bar{a}, \bar{b} \in X^n$ the relation $\bar{a} \sim^G \bar{b}$ is absolute between \mathbb{V} and any forcing extension $\mathbb{V}[H]$. To see this, note that it suffices to check that qf_n is absolute; and in turn it suffices to check that each R_i^n is absolute.

But $\bar{a} \in R_i^n$ iff for some or any sequence $(\bar{d}_m : m \in \omega)$ from D^n converging to \bar{a} , we have that for large enough m , $D' \cdot \bar{d}_m \cap U_i^n = \emptyset$.

It is not hard to check that the range of qf_n is analytic ($\Gamma(\bar{x})$ is in the range of qf_n iff there is a convergent sequence $(\bar{d}_m : m < \omega)$ from D^n satisfying various Borel properties). Hence by Shoenfield Absoluteness, the range of qf_n is absolute.

As notation, call an L -structure $N \in HC$ nice if it is isomorphic to a substructure of M_X . Let \mathcal{W} consist of all nice L -structures.

Lemma 4.3. *An L -structure $N \in HC$ is nice if and only if for every $n \geq 1$, every quantifier-free n -type realized in N is realized in M_X .*

Proof. Left to right is obvious. For the converse, choose any $N \in HC$ for which every quantifier-free n type realized in N is realized in M_X . We construct an L -embedding of N into M_X via a “forth” construction using the homogeneity of M_X . Enumerate the universe of $N = \{a_n : n \in \omega\}$ and let \bar{a}_n denote $\langle a_i : i < n \rangle$. Assuming $f_n : \bar{a}_n \rightarrow M_X$ has been defined, choose any $\bar{b} \in X^{n+1}$ such that $\text{qf}_{n+1}(\bar{a}_{n+1}) = \text{qf}_{n+1}(\bar{b})$. Write \bar{b} as $\bar{b}_n b^*$. As $\text{qf}_n(\bar{b}_n) = \text{qf}_n(f_n(\bar{a}_n))$, there is an automorphism σ of M_X with $\sigma(\bar{b}_n) = f_n(\bar{a}_n)$. Then define f_{n+1} to extend f_n and satisfy $f_{n+1}(a_n) = \sigma(b^*)$. \square

Define a map $f : \mathcal{P}_{\aleph_1}(X) \rightarrow \mathcal{W}$ by $A \mapsto M_A$, the substructure of M_X with universe A .

Our first goal is the following Theorem.

Theorem 4.4. *Suppose (X, d, D, i) is a strongly definable Polish space, (G, d', D', i', \times) is a strongly definable, persistently compact Polish group, and \cdot is a strongly definable continuous action of G on X . Then:*

1. Both \mathcal{W} and f are strongly definable;
2. Persistently, for all $A, B \in \mathcal{P}_{\aleph_1}(X)$, $A \sim^G B$ if and only if $f(A) \cong f(B)$;
3. The canonical Scott sentence map $\text{css} : (\mathcal{W}, \cong) \rightarrow \text{CSS}(\mathcal{W})$ is a representation, where $\text{CSS}(\mathcal{W}) = \{\text{css}(N) : N \in \mathcal{W}\}$;
4. The quotient $(\mathcal{P}_{\aleph_1}(X), \sim^G)$ is representable via the composition map $\text{css} \circ f$ that takes $A \mapsto \text{css}(M_A)$.

Proof. (1) It is obvious that f is strongly definable.

That \mathcal{W} is strongly definable follows from Lemma 4.3 and the absoluteness results mentioned above. In particular, for any L -structure $N \in HC$ that is not nice, there is some n and $\bar{a} \in N^n$ such that $\Gamma := \text{qf}_n(\bar{a})$ is not realized in M_X . But then, in any forcing extension $\mathbb{V}[H]$, $(M_X)^{\mathbb{V}[H]}$ does not realize Γ , so N is not nice in $\mathbb{V}[H]$. As this argument relativizes to any forcing extension, \mathcal{W} is strongly definable.

For (2), if $A \sim^G B$, then any $g \in G$ that satisfies $g \cdot A = B$ induces a bijection between A and B such that $\bar{a} \sim^G g \cdot \bar{a}$ for all $\bar{a} \in A^{<\omega}$. As this implies $G \cdot \bar{a} = G \cdot (g \cdot \bar{a})$, $\text{qf}_n(\bar{a})$ in M_A is equal to $\text{qf}_n(g \cdot \bar{a})$ in M_B . Thus, the action by g induces an isomorphism of the L -structures M_A and M_B . Conversely, suppose $\sigma : M_A \rightarrow M_B$ is an L -isomorphism. Then $\sigma(\bar{a}) \sim^G \bar{a}$ for every $\bar{a} \in A^{<\omega}$, so $A \sim^G B$ by Lemma 4.2.

(3) As \mathcal{W} is strongly definable, this follows immediately from Lemmas 3.3 and 2.25.

(4) follows immediately from (1), (2), and (3). \square

As a consequence of Theorem 4.4, $\|(\mathcal{P}_{\aleph_1}(X), \sim^G)\|$ is defined. For an arbitrary compact group action, this quotient need not be short. Indeed, Theorem 6.8 gives an example where it is not. However, if we additionally assume that G is abelian, then we will see below that $\|(\mathcal{P}_{\aleph_1}(X), \sim^G)\| \leq \beth_2$. The reason for this stark discrepancy is due to the comparative simplicity of abelian group actions. In particular, if an abelian group G acts transitively on a set S , then S is essentially an affine copy of $G/\text{Stab}(a)$, where $\text{Stab}(a)$ is the subgroup of G stabilizing any particular $a \in S$. The following Lemma is really a restatement of this observation.

Lemma 4.5. *Suppose that an abelian group G acts on a set X . Then for every $n \geq 1$, if three n -tuples $\bar{a}, \bar{b}, \bar{c} \in X^n$ satisfy $\bar{a} \sim^G \bar{b} \sim^G \bar{c}$ and $\overline{ab} \sim^G \overline{ac}$, then $\bar{b} = \bar{c}$.*

Proof. Let $g \in G$ be such that $g\bar{a} = \bar{b}$. Choose $h \in G$ such that $h(\overline{ab}) = \overline{ac}$. Then in particular, $h\bar{a} = \bar{a}$ and $h\bar{b} = \bar{c}$. From this, $gh\bar{a} = \bar{b}$ and $hg\bar{a} = \bar{c}$. But $gh = hg$, so $\bar{b} = \bar{c}$, as desired. \square

We will show that $\|(\mathcal{W}, \cong)\| \leq \beth_2$ by showing that each Scott sentence in the representation is from $L_{\beth_1^+ \omega}$, and then using the fact that there are at most \beth_2 such sentences. In the case that the Scott sentence is satisfiable, this means the model has size at most \beth_1 , so this is perhaps unsurprising. We will accomplish this complexity bound by a type-counting argument; here is the notion of type we will use.

If φ is a canonical Scott sentence – that is, $\varphi \in \text{CSS}(L)_{\text{ptl}}$ – then let $S_n^\infty(\varphi)$ be the set of all canonical Scott sentences in the language $L' = L \cup \{c_0, \dots, c_{n-1}\}$ which imply φ . We will refer to elements of $S_n^\infty(\varphi)$ as types – infinitary formulas with free variables x_0, \dots, x_{n-1} , resulting from replacing each c_i with a new variable x_i not otherwise appearing in the formula. It is equivalent to define $S_n^\infty(\varphi)$ by forcing – if $\mathbb{V}[H]$ makes φ hereditarily countable and $M \in \mathbb{V}[H]$ is the unique countable model of φ , then $S_n^\infty(\varphi)$ is the set $\{\text{css}(M, \bar{a}) : \bar{a} \in M^n\}$. Evidently this set depends only on the isomorphism class of M , so by the usual argument with Lemma 2.5, this set is in \mathbb{V} .

Proposition 4.6. *Suppose φ is a canonical Scott sentence in a language of size at most κ , and for all n , $|S_n^\infty(\varphi)| \leq \kappa$, where κ is an infinite cardinal. Then φ is a sentence of $L_{\kappa+\omega}$.*

Proof. We use the precise syntactic definition of Scott formulas from Definition 3.1. For a moment, pass to a forcing extension $\mathbb{V}[G]$ in which φ is hereditarily countable, and let M be its unique countable model. For each ordinal α , let $S_n^\alpha(\varphi)$ be the set $\{\varphi_\alpha^\alpha(\bar{x}) : \bar{x} \in M^n\}$. Since M is (persistently above $\mathbb{V}[G]$) unique up to isomorphism, and since this set is unchanged by passing to an isomorphic image of M , $S_n^\alpha(\varphi)$ is in \mathbb{V} and depends only on φ . Moreover, there is a natural surjection $\pi_n^\alpha : S_n^\infty(\varphi) \rightarrow S_n^\alpha(\varphi)$ taking $\text{css}(M, \bar{a})$ to $\varphi_\alpha^\alpha(\bar{x})$; each π_n^α is in \mathbb{V} .

Let α^* be the Scott rank of M . Again, this is invariant under isomorphism, so depends only on φ . For any two sentences $\psi, \tau \in S_n^\infty(\varphi)$, let $d(\psi, \tau)$ be the least $\alpha < \alpha^*$ where $\pi_n^{\alpha+1}(\psi) \neq \pi_n^{\alpha+1}(\tau)$; if there is no

such α , then $\psi = \tau$, so say $d(\psi, \tau) = \alpha^*$. It is immediate from the construction of Scott rank that if $\alpha \leq \alpha^*$, there are Scott sentences ψ and τ of some arity where $d(\psi, \tau) = \alpha$; hence $d : \bigcup_n (S_n^\infty(\varphi))^2 \rightarrow \alpha^* + 1$ is surjective. Further, d depends only on φ , so by Lemma 2.5, $d \in \mathbb{V}$.

The rest of the proof takes place in \mathbb{V} . Because π_n^α is surjective and $|S_n^\infty(\varphi)| \leq \kappa$, $|S_n^\alpha(\varphi)| \leq \kappa$ for all α . Similarly, $|\bigcup_n (S_n^\infty(\varphi))^2| \leq \kappa$ and d is surjective, so $|\alpha^* + 1| \leq \kappa$. By induction we show that for all $\alpha \leq \alpha^* + 1$, $S_n^\alpha(\varphi) \subseteq L_{\kappa+\omega}$. The base case is trivial, since there are only κ atomic formulas. The step follows from the fact that $|S_n^\alpha(\varphi)| \leq \kappa$, and the limit follows from the fact that $\alpha^* < \kappa^+$, so in both cases we need only take conjunctions and disjunctions of κ formulas at a time.

Observe that φ is precisely the following:

$$\pi_0^{\alpha^*}(\varphi) \wedge \bigwedge \left\{ \forall \bar{x} \left(\pi_n^{\alpha^*}(\varphi^*)(\bar{x}) \rightarrow \pi_n^{\alpha^*+1}(\varphi^*)(\bar{x}) \right) : n \in \omega, \varphi^* \in S_n^\infty(\varphi) \right\}$$

Since $S_n^\alpha(\varphi) \subseteq L_{\kappa+\omega}$ for all α and n , and since they all have size at most κ , φ is in $L_{\kappa+\omega}$, as desired. \square

The following holds by a straightforward induction on the complexity of formulas:

Lemma 4.7. *For all infinite cardinals κ and languages L of size at most κ , there are exactly 2^κ different $L_{\kappa+\omega}$ formulas.*

Now we can prove our theorem. Recall that \sim_G is the diagonal equivalence relation on $\mathcal{P}_{\aleph_1}(X)$, induced by the diagonal action of G .

Theorem 4.8. *Let X be a strongly definable Polish space, let G be a strongly definable, persistently compact abelian group, and suppose \cdot is a strongly definable continuous action of G on X . Suppose all this holds persistently. Then $\|(\mathcal{P}_{\aleph_1}(X), \sim_G)\| \leq \beth_2$.*

Proof. We use Proposition 4.6 to show that $\text{CSS}(\mathcal{W})_{\text{ptl}} \subseteq L_{\beth_1^+}$. Given this, then by Lemma 4.7, we conclude that $|\text{CSS}(\mathcal{W})_{\text{ptl}}| \leq \beth_2$, as desired. So let $\varphi \in \text{CSS}(\mathcal{W})_{\text{ptl}}$ be arbitrary; it is enough to show that $|S_n^\infty(\varphi)| \leq \beth_1$.

For each n , let $\text{qf}_n(\varphi)$ be the set of quantifier-free n -types which are consistent with φ . We have a surjective map $\pi_n : S_n^\infty(\varphi) \rightarrow \text{qf}_n(\varphi)$ sending $\psi(\bar{x})$ to the set of quantifier-free formulas in \bar{x} which it implies. For any $p \in \text{qf}_n(\varphi)$, let $S_n^\infty(\varphi, p)$ be $\pi_n^{-1}(p)$, the set of $\psi \in S_n^\infty(\varphi)$ where $\pi_n(\psi) = p$. (All of these definitions have taken place in \mathbb{V} .) Since the language is countable, $|\text{qf}_n(\varphi)| \leq \beth_1$. Thus it is sufficient to show that for all p , $|S_n^\infty(\varphi, p)| \leq \beth_1$.

Now we take advantage of the fact that G is abelian:

Claim: Suppose $p^* \in \text{qf}_{2n}(\varphi)$ is such that $p^* \upharpoonright_{[0,n]} = p^* \upharpoonright_{[n,2n]} = p$. Suppose that $\psi^*, \tau^* \in S_{2n}^\infty(\varphi)$ both complete p^* . Further, suppose $\psi^* \upharpoonright_{[0,n]} = \tau^* \upharpoonright_{[0,n]}$. Then $\psi^* = \tau^*$.

Proof: Pass to a forcing extension $\mathbb{V}[H]$ in which φ is hereditarily countable, and let M be its unique countable model. By Theorem 4.4, we may assume $M = M_A$ for some $A \in \mathcal{P}_{\aleph_1}(X)^{\mathbb{V}[H]}$. Choose some tuples $(\bar{a}_0 \bar{a}_1)$ and $(\bar{b}_0 \bar{b}_1)$ from M^{2n} where $\text{css}(M, \bar{a}_0 \bar{a}_1) = \psi^*$ and $\text{css}(M, \bar{b}_0 \bar{b}_1) = \tau^*$. By assumption $\text{css}(M, \bar{a}_0) = \text{css}(M, \bar{b}_0)$, so we may assume $\bar{a}_0 = \bar{b}_0$. Since all of the tuples \bar{b}_0 , \bar{a}_1 , and \bar{b}_1 have the same quantifier-free type, they

are in the same G -orbit, and similarly with $\bar{b}_0\bar{a}_1$ and $\bar{b}_0\bar{b}_1$. Thus Lemma 4.5 applies directly to the triple $(\bar{b}_0, \bar{b}_1, \bar{a}_1)$, so in particular $\bar{b}_1 = \bar{a}_1$. Thus $\psi^* = \text{css}(M, \overline{b_0a_1}) = \text{css}(M, \overline{b_0b_1}) = \tau^*$, as desired. \square

Fix some $\psi \in S_n^\infty(\varphi, p)$, and define $\Gamma(\psi)$ to be the set of all $p^* \in \text{qf}_{2n}(\varphi)$ such that $p^* \upharpoonright_{[0,n]} = p^* \upharpoonright_{[n,2n]} = p$ and such that for some $\psi^* \in S_{2n}^\infty(\varphi, p^*)$, $\psi^* \upharpoonright_{[0,n]} = \psi$.

By the Claim, if $p^* \in \Gamma(\psi)$, there is a unique $\psi^* \in S_{2n}^\infty(\varphi)$ where $\pi_{2n}(\psi^*) = p^*$ and $\psi^* \upharpoonright_{[0,n]} = \psi$. So define $F(p^*)$ to be $\psi^* \upharpoonright_{[n,2n]}$. Evidently $|\Gamma(\psi)| \leq \beth_1$, so it is enough to show that $F : \Gamma(\psi) \rightarrow S_n^\infty(\varphi, p)$ is surjective.

But this is almost immediate. Fix any $\tau \in S_n^\infty(\varphi, p)$ and let $\mathbb{V}[H]$ be a forcing extension in which φ is hereditarily countable, and let M be its unique countable model; as before, we may assume $M = M_A$ for some $A \in \mathcal{P}_{\aleph_1}(X)^{\mathbb{V}[H]}$. Choose any $\bar{a} \in A^n$ where $\text{css}(M, \bar{a}) = \psi$ and any $\bar{b} \in A^n$ where $\text{css}(M, \bar{b}) = \tau$. Finally, let p^* be the quantifier-free type of \overline{ab} in M . Clearly $p^* \in \Gamma(\psi)$ and $F(p^*) = \tau$. \square

This theorem will be crucial in Section 6.

5 Refining Equivalence Relations

We begin by defining an incomplete first-order theory REF. Its language is $L = \{E_n : n \in \omega\}$ and its axioms posit:

- Each E_n is an equivalence relation;
- E_0 has a single equivalence class; that is, we consider $x E_0 y$ to be universally true;
- For all n , E_{n+1} refines E_n ; that is, every E_n -class is a union of E_{n+1} -classes.

The theory REF is very weak, which makes the generality of the following proposition surprising.

Proposition 5.1. *REF is grounded.*

Proof. We begin with an analysis of an arbitrary model M of REF. As notation, for any $a \in M$ and $n \in \omega$, let $[a]_n$ denote the equivalence class of a , i.e., $\{b \in M : M \models E_n(a, b)\}$. As the equivalence relations refine each other, the classes $T(M) = \{[a]_n : a \in M, n \in \omega\}$ form an ω -tree, ordered by $[a]_n \leq [b]_m$ if and only if $n \leq m$ and $[b]_m \subseteq [a]_n$. Next, let E_∞ be the equivalence relation given by $E_\infty(a, b)$ if and only if $E_n(a, b)$ for every $n \in \omega$. Let $[a]_\infty$ be the E_∞ -class of a . Then M/E_∞ can be construed as a subset of the branches $[T(M)]$ of $T(M)$. As we are interested in determining models up to back-and-forth equivalence (as opposed to isomorphism), the following definition is natural.

For each $a \in M$, let the *color of a* , $c(a) \in (\omega + 1) \setminus \{0\}$ be given by

$$c(a) = \begin{cases} |[a]_\infty| & \text{if } [a]_\infty \text{ is finite} \\ \omega & \text{if } [a]_\infty \text{ is infinite} \end{cases}$$

Next, we describe some expansions of M to larger languages. For each $n \in \omega$, let $L_n = L \cup \{U_i : i \leq n\}$, where the U_i 's are distinct unary predicates. Given any $M \models \text{REF}$, $n \in \omega$, and $a \in M$, let $M_n(a)$ denote the L_n -structure $(M, [a]_0, \dots, [a]_n)$, i.e., where each predicate U_i is interpreted as $[a]_i$.

We now exhibit some invariants, which we term the *data of M* , written $D(M)$ which we will see only depend on the $\equiv_{\infty, \omega}$ -equivalence class of M .

For each $n \in \omega$, let

$$I_n(M) = \{\text{css}(M_n(a)) : a \in M\}.$$

We combine the sets $I_n(M)$ into a tree $(I(M), \leq)$ where $I(M) = \bigcup_{n \in \omega} I_n(M)$ and, for $\sigma_n \in I_n(M)$ and $\psi_m \in I_m(M)$, we say $\sigma_n \leq \psi_m$ if and only if $n \leq m$ and $\psi_m \vdash \sigma_n$. That is, if in any forcing extension the reduct of any model of ψ_m to L_n is a model of σ_n . Then clearly $(I(M), \leq)$ is an ω -tree.

Continuing, for each $n > 0$ and $\sigma_n \in I_n(M)$, let the multiplicity of σ_n , $\text{mult}_M(\sigma_n) \in (\omega + 1) \setminus \{0\}$, be given by: $\text{mult}_M(\sigma_n) = k < \omega$ if k is maximal such that there are elements $\{b_i : i < k\} \subseteq M$ such that

$$\bigwedge_{i < j < k} \left[E_{n-1}(b_i, b_j) \wedge \neg E_n(b_i, b_j) \right] \wedge \bigwedge_{i < k} \text{css}(M_n(b_i)) = \sigma_n$$

and let $\text{mult}_M(\sigma_n) = \omega$ if there is an infinite family $\{b_i : i < \omega\}$ as above.

Now, each $a \in M$ induces a canonical sequence $\text{Seq}_M(a) := \langle \text{css}(M_n(a)) : n \in \omega \rangle$, which is clearly a branch through the tree $I(M)$, and depends only on $\text{css}(M, a)$. Let $\text{Seq}(M) = \{\text{Seq}_M(a) : a \in M\}$. So $\text{Seq}(M) \subseteq [I(M)]$, the set of branches of $I(M)$. Finally, for any $s \in \text{Seq}(M)$, we define the color spectrum of s as $\text{Sp}_M(s) := \{c(a) : \text{Seq}_M(a) = s\}$. Thus, each $\text{Sp}_M(s)$ is a non-empty subset of $(\omega + 1) \setminus \{0\}$.

Define the *data of M* , $D(M) := \langle (I(M), \leq), \text{mult}_M, \text{Seq}(M), \text{Sp}_M \rangle$.

Claim 1: For any $M, N \models \text{REF}$, $M \equiv_{\infty, \omega} N$ if and only if $D(M) = D(N)$.

Proof: First, note that if $D(M) = D(N)$, then as the trees $(I(M), \leq)$ and $(I(N), \leq)$ are equal, they have the same root, so $\text{css}(M_0(a)) = \text{css}(N_0(b))$ for some/every $a \in M, b \in N$. So $M \equiv_{\infty, \omega} N$.

For the forward direction, it is easy to check that $D(M)$ only depends on the isomorphism type of M , and also that D is absolute to forcing extensions. Hence if $M \equiv_{\infty, \omega} N$, then pass to a forcing extension $\mathbb{V}[G]$ in which $M \cong N$; then we get $(D(M))^\mathbb{V} = (D(M))^{\mathbb{V}[G]} = (D(N))^{\mathbb{V}[G]} = (D(N))^\mathbb{V}$. \square

To begin the proof of groundedness, choose any $\sigma \in \text{CSS}(\text{REF})_{\text{ptl}}$. Choose any forcing extension $\mathbb{V}[G]$ of \mathbb{V} in which $\sigma \in \text{HC}^{\mathbb{V}[G]}$ and hence $\sigma \in \text{CSS}(\text{REF})^{\mathbb{V}[G]}$. Choose any model $M \in \mathbb{V}[G]$ with $\text{css}(M) = \sigma$. Working in $\mathbb{V}[G]$, compute $D(M)$, the data of M . However, in light of Claim 1, $D(M)$ only depends on σ , and so by Lemma 2.5 $D(M) \in \mathbb{V}$. As σ is fixed, for the remainder of the argument we write

$$D = \langle (I, \leq), \text{mult}, \text{Seq}, \text{Sp} \rangle.$$

To complete the proof of the Proposition, we work in \mathbb{V} and ‘unpack’ the data D to construct an L -structure $N \in \mathbb{V}$ such that in $\mathbb{V}[G]$, $M \equiv_{\infty, \omega} N$. Once we have this, as $\sigma = \text{css}(M)$, it follows that $N \models \sigma$ and

so N witnesses that $\sigma \in \text{CSS}(\text{REF})_{\text{sat}}$. That is, the proof of groundedness will be finished once we establish the following Claim.

Claim 2: There is an L -structure $N \in \mathbb{V}$ such that $\mathbb{V}[G] \models 'N \equiv_{\infty, \omega} M'$.

Proof: Before beginning the ‘unpacking’ of D , we note some connections between M and D that are not part of the data. First, there is a surjective tree homomorphism $h : T(M) \cup M/E_\infty \rightarrow I \cup \text{Seq}$ given by $h([a]_n) = \text{css}(M_n(a))$ for $n \in \omega$ and $h([a]_\infty) = \langle \text{css}(M_n(a)) : n \in \omega \rangle$. Note that for each $s \in \text{Seq}$ and each $k \in \text{Sp}(s)$, $\{[a]_\infty : h(a) = s \text{ and } c(a) = k\}$ is dense in $h^{-1}(s)$. The following relationship between M and h follows quickly:

$(\star)_{M,h}$: For every $n \geq 1$, $s \in \text{Seq}$, $k \in \text{Sp}(s)$, and $a \in M$ such that $h([a]_{n-1}) = s(n-1)$, there are pairwise E_n -inequivalent $\{d_i : i < \text{mult}(s(n))\} \subseteq M$ such that

$$\bigwedge_{i < \text{mult}(s(n))} \left[E_{n-1}(d_i, a) \wedge h([d_i]_\infty) = s \wedge c(d_i) = k \right]$$

We also identify two species of elements of Seq . Call $s \in \text{Seq}$ of *isolated type* if there is $n \in \omega$ such that $\text{mult}(s(m)) = 1$ for every $m \geq n$ and of *perfect type* otherwise. The latter name is apt, as $h^{-1}(s)$ is perfect (has no isolated points) whenever s is not of isolated type. We argue that if $s \in \text{Seq}$ is of isolated type, then $\text{Sp}(s)$ is a singleton. Indeed, choose n such that $\text{mult}(s(m)) = 1$ for every $m \geq n$ and choose $a, b \in M$ such that $h([a]_\infty) = h([b]_\infty) = s$. We will show that $c(a) = c(b)$. To see this, by applying $(\star)_{M,h}$ at level $n+1$ with $k = c(b)$, get $d \in M$ such that $E_n(a, d)$, $h([d]_\infty) = s$, and $c(d) = c(b)$. But now, as $h([a]_\infty) = h([d]_\infty) = s$, the choice of n implies that $E_\infty(a, d)$. Thus, $c(a) = c(d) = c(b)$ as required.

We begin ‘unpacking’ D by inductively constructing an ω -tree (J, \leq) and a surjective tree homomorphism $h' : (J, \leq) \rightarrow (I, \leq)$. Begin the construction of $J = \bigcup_{n \in \omega} J_n$ by taking $J_0 = \{\rho_0\}$ to be a singleton and defining $h'(\rho_0) = \sigma$. Suppose the n th level J_n has been defined, together with $h' : \bigcup_{j \leq n} J_j \rightarrow \bigcup_{j \leq n} I_j$. For each $\rho_n \in J_n$, we define its immediate successors $\text{Succ}_J(\rho_n)$ as follows. Look at $\text{Succ}_I(h'(\rho_n)) \subseteq I_{n+1}$. For each $\sigma_{n+1} \in \text{Succ}_I(h'(\rho_n))$, choose a set $A_{n+1}(\sigma_{n+1})$ of cardinality $\text{mult}(\sigma_{n+1}) \in (\omega + 1) \setminus \{0\}$ such that the sets $A_{n+1}(\sigma_{n+1})$ are pairwise disjoint. Let

$$\text{Succ}_J(\rho_n) := \bigcup \{A_{n+1}(\sigma_{n+1}) : \sigma_{n+1} \in \text{Succ}_I(h'(\rho_n))\}$$

and put $J_{n+1} := \bigcup \{\text{Succ}_J(\rho_n) : \rho_n \in J_n\}$. We extend h' by $h'(\rho) = \sigma_{n+1}$ for every $\rho \in A_{n+1}(\sigma_{n+1})$.

Now, having completed the construction of (J, \leq) and the tree homomorphism $h' : (J, \leq) \rightarrow (I, \leq)$, there is a unique extension (which we also call h') $h' : [J] \rightarrow [I]$ from the branches of J to the branches of I such that $h'(\eta) = s$ if and only if $h'(\eta \upharpoonright_n) = s(n-1)$ for every $n \in \omega$.

The universe of the L -structure N we are building will be a subset of $(h')^{-1}(\text{Seq}) \times (\omega + 1)$ and for $(\eta, i), (\nu, j) \in N$, we will interpret E_n by

$$E_n((\eta, i), (\nu, j)) \quad \text{if and only if} \quad \eta \upharpoonright_n = \nu \upharpoonright_n.$$

In particular, we will have $[(\eta, i)]_\infty = \{(\eta, j) : (\eta, j) \in N\}$. To finish our description of N , we must assign a ‘color’ to each element of $(h')^{-1}(Seq)$. Fix s ; we assign colors to $(h')^{-1}(s)$. First, if s is of isolated type, then from above, we know that $Sp(s) = \{k\}$ for a single color $k \leq \omega$. Accordingly, put elements $\{(\eta, i) : i < k\}$ into the universe of N for every η satisfying $h'(\eta) = s$. For each $s \in Seq$ that is not of isolated type, note that $(h')^{-1}(s)$ has no isolated points. Thus, we can choose a partition $(h')^{-1}(s) = \bigcup D_k(s)$ into disjoint dense subsets indexed by colors $k \in Sp(s)$. Then, for each $\eta \in D_k(s)$ put elements $\{(\eta, i) : i < k\}$ into the universe of N . This completes our construction of the L -structure $N \in \mathbb{V}$, and it is easily verified that this construction entails $(\star)_{N, h'}$.

We now work in $\mathbb{V}[G]$ and demonstrate that $M \equiv_{\infty, \omega} N$. Indeed, all that we need for this is that in $\mathbb{V}[G]$, both $(\star)_{M, h}$ and $(\star)_{N, h'}$ hold. Let \mathcal{F} consist of all (\bar{a}, \bar{b}) such that $\text{lg}(\bar{a}) = \text{lg}(\bar{b})$, \bar{a} from M , and \bar{b} from N that satisfy for each $i < \text{lg}(\bar{a})$, $c(a_i) = c(b_i)$ and $h([a_i]_\infty) = h'([b_i]_\infty)$; and for each $n \in \omega$, $i < j < \text{lg}(\bar{a})$, $M \models E_n(a_i, a_j)$ if and only if $N \models E_n(b_i, b_j)$ and $a_i = a_j$ if and only if $b_i = b_j$.

To see that \mathcal{F} is a back-and-forth system, choose any $(\bar{a}, \bar{b}) \in \mathcal{F}$ and choose any $a^* \in M$. We will find $b^* \in N$ such that $(\bar{a}a^*, \bar{b}b^*) \in \mathcal{F}$, and the argument in the other direction is symmetric. If $\text{lg}(\bar{a}) = 0$, or if $a^* \in \bar{a}$, it is obvious what to do, so assume this is not the case. If $E_\infty(a^*, a_i)$ for some i , then as $c(a_i) = c(b_i)$, we can find $b^* \notin \bar{b}$ such that $E_\infty(b^*, b_i)$ which suffices.

Now assume that $\neg E_\infty(a^*, a_i)$ holds for each i . Let $k = c(a^*)$ and $s = h([a^*]_\infty)$. Let $n > 0$ be least such that $\neg E_n(a^*, a_i)$ for all i . Let $A_1 = \{a_i : E_{n-1}(a^*, a_i)\}$ and let B_1 be the associated subset of \bar{b} . By the axioms of REF it suffices to find $b^* \in N$ such that $c(b^*) = k$, $h'([b^*]_\infty) = s$, $E_{n-1}(b^*, b)$ for some/every $b \in B_1$, but $\neg E_n(b^*, b)$ for every $b \in B_1$. To find such an element, let

$$A_2 = \{a \in A_1 : \text{there is some } a' \in [a]_n \text{ such that } c(a') = k \text{ and } h([a']_\infty) = s\}$$

Let $A_3 \subseteq A_2$ be any maximal, pairwise E_n -inequivalent subset of A_2 and let $\ell = |A_3|$. The set $\{a^*\} \cup A_3$ witnesses that $\text{mult}(s(n)) > \ell$. [More precisely, for each $a \in A_3$, choose $a' \in [a]_n$ with $c(a') = k$ and $h([a']_\infty) = s$. Then $\{a^*\} \cup \{a' : a \in A_3\}$ witnesses $\text{mult}(s(n)) > \ell$.] Let B_3 be the associated subset of \bar{b} ; so $|B_3| = \ell$.

Choose $a_i \in A_1$. Then by $(\star)_{N, h'}$, applied at b_i (noting that $[b_i]_{n-1} = s(n-1)$), choose a family $\{d_i : i < \text{mult}(s(n))\}$ as there. By pigeon-hole choose an $i^* < \text{mult}(s(n))$ such that $\neg E_n(d_{i^*}, b)$ holds for all $b \in B_3$. It is easily checked that d_{i^*} is a possible choice for b^* . As noted above, this completes the proof of the Claim. \square

In particular, $N \models \sigma$, establishing groundedness. \square

We now turn our attention to two classical complete theories extending REF. These are often given as first examples in stability theory. We denote them by REF(inf) and REF(bin), respectively. REF(bin) is the extension of REF asserting that for every n , E_{n+1} partitions each E_n -class into two E_{n+1} -classes, while REF(inf) asserts that for all n , E_{n+1} partitions each E_n -class into infinitely many E_{n+1} -classes.

The following facts are well known.

Fact 5.2. *Both $REF(bin)$ and $REF(inf)$ are complete theories that admit quantifier elimination.*

- $REF(bin)$ is superstable but not ω -stable; and
- $REF(inf)$ is stable, but not superstable.

We will see below that these theories have extremely different countable model theory. Both are similar in that the isomorphism relation \cong is not Borel. However, it turns out that $REF(inf)$ is Borel complete, and indeed, is λ -Borel complete for every λ . That is, $REF(inf)$ the class of models is maximally complicated, both at the countable level as well as at every uncountable level. On the other hand, $REF(bin)$ is far from being Borel complete. In fact, its class of countable models embeds T_2 but not T_3 .

5.1 Finite Branching

In this subsection we show that $T_2 \leq_B REF(bin)$ but $T_3 \not\leq_B REF(bin)$, and that the isomorphism relation of $REF(bin)$ is not Borel.

For the following, it would be inconvenient to work with T_2 directly. Instead, let F_2 be the equivalence relation on $(2^\omega)^\omega$ defined by: $(x_n : n \in \omega) F_2 (y_n : n \in \omega)$ iff $\{x_n : n \in \omega\} = \{y_n : n \in \omega\}$. Then the quotient $(2^\omega)^\omega / F_2$ is in natural bijection with $\mathcal{P}_{\aleph_1}(2^\omega) \setminus \{\emptyset\}$, so we think of F_2 as representing countable sets of reals. It is not hard to check that $(Mod(T_2), \cong)$ is Borel bireducible with $((2^\omega)^\omega, F_2)$. So for T a theory, showing that $T_2 \leq_B T$ is the same thing as showing $F_2 \leq_B T$.

Theorem 5.3. $T_2 \leq_B REF(bin)$.

Proof. Begin by building a special countable model M of $REF(bin)$. Let S be the set of sequences from 2^ω which are eventually zero, and fix a bijection $c : S \rightarrow \mathbb{N}$. Let M be the set of all (η, n) where $\eta \in S$ and $n < c(\eta)$. As usual, say $(\eta_1, n_1) E_m (\eta_2, n_2)$ holds if and only if η_1 and η_2 agree on the first m places. Clearly M is a model of $REF(bin)$, and the color of (η, n) is exactly $c(\eta)$; observe that no element has color \aleph_0 . (Recall that the color of a is the cardinality of $[a]_\infty$.) We will construct our models as superstructures of M , whose new elements all have color \aleph_0 and are not E_∞ -equivalent to any element of M .

Let $X \subseteq (2^\omega)^\omega$ be the set of all $(x_n : n \in \omega)$ such that each $x_n \notin S$. Then $(X, F_2 \upharpoonright_X) \cong_B ((2^\omega)^\omega, F_2)$, via any Borel bijection between 2^ω and $2^\omega \setminus S$. (By Corollary 13.4 and Theorem 4.6 of [11], any two uncountable Borel sets are in Borel bijection.) So it suffices to show that $(X, F_2 \upharpoonright_X) \leq_B REF(bin)$. Given $I \subseteq 2^\omega \setminus S$ countable, let M_I be the L -structure extending M with universe $M \cup (I \times \omega)$, where again, $(\eta_1, n_1) E_m (\eta_2, n_2)$ holds if and only if $\eta_1 \upharpoonright_m = \eta_2 \upharpoonright_m$.

It is not hard to check that one can define a Borel map $f : X \rightarrow Mod(REF(bin))$, such that for all $\bar{x} = (x_n : n \in \omega) \in X$, $f(\bar{x}) \cong M_{\{x_n : n \in \omega\}}$. Given that, it suffices to show that for all distinct $I, J \subseteq 2^\omega \setminus S$ countable, $M_I \not\cong M_J$.

So suppose $M_I \cong M_J$, say via $g : M_I \rightarrow M_J$. I aim to show that for all $(\eta, n) \in M_I$, $g(\eta, n) = (\eta, n')$ for some $n' < \omega$. This suffices to show $I = J$, since then $I = \{\eta : (\eta, n) \in M_I \text{ for all } n\} = \{\eta : (\eta, n) \in$

M_J for all $n\} = J$. So let $(\eta, n) \in M_I$; write $g(\eta, n) = (\tau, n')$. I show for each $m < \omega$ that $\eta \upharpoonright_m = \tau \upharpoonright_m$. Indeed, pick $\nu \in S$ such that $\nu \upharpoonright_m = \eta \upharpoonright_m$. Then $g(\nu, 0) = (\nu, k)$ for some $k < c(\tau)$, since $g([\nu, 0]_\infty)$ is the unique E_∞ -class of M_J of size $c(\nu)$. Then since $((\eta, n)E_m(\tau, 0))^{M_I}$, we have $((\tau, n')E_m(\nu, k))^{M_J}$. Hence $\tau \upharpoonright_m = \nu \upharpoonright_m = \eta \upharpoonright_m$. \square

To show that T_3 does not embed into $\text{REF}(\text{bin})$, we clarify $\equiv_{\infty\omega}$ -equivalence on a slightly wider class of L -structures. Let $\text{REF}(\text{fin})$ denote the sentence of $L_{\omega_1, \omega}$ extending REF asserting that every E_{n+1} -class partitions every E_n -class into finitely many E_{n+1} -classes.

Lemma 5.4. *Every model M of $\text{REF}(\text{fin})$ has an $\equiv_{\infty\omega}$ -equivalent submodel $N \subseteq M$ of size at most \beth_1 .*

Proof. For each E_∞ -class $[a]_\infty \subseteq M$, let

$$B([a]_\infty) = \begin{cases} [a]_\infty & \text{if } [a]_\infty \text{ is countable} \\ \text{any countably infinite subset of } [a]_\infty & \text{if } [a]_\infty \text{ is uncountable} \end{cases}$$

and let N be the substructure of M with universe $\bigcup\{B([a]_\infty) : a \in M\}$. It is easily seen that $N \equiv_{\infty, \omega} M$. That N has size at most continuum follows from the finite splitting at each level. \square

Combined with groundedness, this gives us the nonembedding result we wanted:

Theorem 5.5. $\|\text{REF}(\text{bin})\| = \|\text{REF}(\text{fin})\| = I_{\infty, \omega}(\text{REF}(\text{bin})) = I_{\infty, \omega}(\text{REF}(\text{fin})) = \beth_2$. *In particular, both $\text{REF}(\text{bin})$ and $\text{REF}(\text{fin})$ are short and $T_3 \not\leq_B \text{REF}(\text{bin}), \text{REF}(\text{fin})$.*

Proof. Recall by Corollary 3.21 that $\|T_2\| = \beth_2$ and $\|T_3\| = \beth_3$. Since $T_2 \leq_B \text{REF}(\text{bin})$, $\beth_2 = \|T_2\| \leq \|\text{REF}(\text{bin})\|$. On the other hand, since REF is grounded, $\|\text{REF}(\text{fin})\| = I_{\infty, \omega}(\text{REF}(\text{fin}))$ but the latter cardinal is bounded above by \beth_2 by Lemma 5.4. Thus, all four cardinals are equal to \beth_2 . So, by definition, both $\text{REF}(\text{bin})$ and $\text{REF}(\text{fin})$ are short. As $\|T_3\| = \beth_3$, the nonembeddability of T_3 into either class follows from Theorem 3.10(2). \square

Finally, we show that isomorphism for $\text{REF}(\text{bin})$ is not Borel.

Theorem 5.6. *Isomorphism on $\text{REF}(\text{bin})$ is not Borel.*

Proof. It is commonly known – see for example Theorem 12.2.4 of [4] – that isomorphism is Borel if and only if, for some α , \equiv_α is sufficient to decide isomorphism. Since \equiv_0 is implied by \equiv and $\text{REF}(\text{bin})$ is a complete theory with more than one model, \equiv_0 does not decide isomorphism. We proceed by induction with a combined step and limit induction step. So suppose $\alpha_0 \leq \alpha_1 \leq \dots$ are such that each \equiv_{α_n} does not decide isomorphism. That is, for each n , there is a pair A_n, B_n of countable models of $\text{REF}(\text{bin})$ which are nonisomorphic but where $A_n \equiv_{\alpha_n} B_n$. Let $\alpha = \sup\{\alpha_n + 1 : n \in \omega\}$. We will construct a pair (indeed, a large family) of countable models of $\text{REF}(\text{bin})$ that are pairwise \equiv_α -equivalent but not isomorphic. This is sufficient.

Recall that among countable models M of $\text{REF}(\text{bin})$, the *color* of an element $a \in M$ is the size of its E_∞ -class $[a]_\infty$. By adding an element to each finite E_∞ class occurring in A_n, B_n , respectively, we can suppose the color “1” does not occur in any of the A_n 's, B_n 's.

Let $\mathfrak{C} \models \text{REF}(\text{bin})$ be the model with universe $2^\omega \times \omega$, where as usual $(\eta, n)E_k(\tau, m)$ iff $\eta \upharpoonright_k = \tau \upharpoonright_k$. \mathfrak{C} will serve as a ‘monster model’ of sorts; in particular we can suppose each A_n, B_n are (elementary) substructures of \mathfrak{C} .

We begin by forming a single countable model $M \preceq \mathfrak{C}$ that encodes all of the complexity of the models A_n, B_n . For $s \in 2^{<\omega}$, let $(A_n)^s$ be a ‘shift’ of A_n by s . Formally, $(A_n)^s = \{(s \frown \eta, j) : (\eta, k) \in A_n\}$ and we define $(B_n)^s$ analogously. Whereas the substructures A_n and $(A_n)^s$ of \mathfrak{C} are certainly not elementarily equivalent, the relationships between A_n and B_n are maintained. That is, if $\text{lg}(s) = \text{lg}(t)$, then for any n , $(A_n)^s \equiv_{\alpha_n} (B_n)^t$, but because $A_n \not\cong B_n$, there is no elementary bijection $f : (A_n)^s \rightarrow (B_n)^t$.

As notation, for $i \in \{0, 1\}$ and $n \in \omega$, let S_n^i be the subset of 2^{2n+2} satisfying

- $s(j) = 0$ for every odd $j < 2n$;
- $s(2n) = i$; and
- $s(2n + 1) = 1$.

Note that not only are the sets S_n^i disjoint, but in fact, $S^* := \bigcup\{S_n^i : i \in \{0, 1\}, n \in \omega\}$ is an antichain on $2^{<\omega}$. Let

$$M := \bigcup_{n < \omega, s \in S_n^0} (A_n)^s \cup \bigcup_{n < \omega, s \in S_n^1} (B_n)^s$$

It is readily checked that $M \preceq \mathfrak{C}$. Because every element of every A_n, B_n has color distinct from 1, no element of M has color 1 either. As notation, we refer to the subsets $(A_n)^s$ and $(B_n)^s$ as the *s-bubbles* of M . Obviously, for a specific choice of s , M contains only one of $(A_n)^s$ or $(B_n)^s$. We write $M(s)$ for this *s-bubble*.

For each $x \in 2^\omega$, let $x^* \in 2^\omega$ be defined by $x^*(j) = 0$ if j is odd and $x^*(j) = x(j/2)$ if j is even. For each countable, dense subset $X \subseteq 2^\omega$, let

$$M_X = M \cup \{(x^*, 0) : x \in X\} \quad \text{and let} \quad S_X^* = S^* \cup \{x^* : x \in X\}$$

Clearly, $M \preceq M_X \preceq \mathfrak{C}$ and an element $c \in M_X$ has color 1 if and only if $c \notin M$. Write $M_X(s) = M(s)$ for $s \in S_n^i$.

Claim 1: Let $X, Y \subset 2^\omega$ be countable and dense. Then $M_X \cong M_Y$ if and only if $X = Y$.

Proof: If $X = Y$ then $M_X = M_Y$. On the other hand, suppose X and Y are dense and $f : M_X \cong M_Y$. We claim that for all $\eta \in X$, $f(\eta^*, 0) = (\eta^*, 0)$. This suffices, since then $X = \{\eta : (\eta^*, n) \in M_X \text{ iff } n = 0\} \subseteq \{\eta : (\eta^*, n) \in M_Y \text{ iff } n = 0\} = Y$ and by symmetry $Y \subseteq X$.

So fix $\eta \in X$ and write $f(\eta^*, 0) = (\tau^*, 0)$ where $\tau \in Y$ ($f(\eta^*, 0)$ must be of this form since it is of color 1 in M_Y). Suppose towards a contradiction that $\eta \neq \tau$; let n be least such that $\eta(n) \neq \tau(n)$. Let $s = \eta^* \upharpoonright_{2n+1} \frown (1)$

and let $t = \tau^* \upharpoonright_{2n+1} \frown (1)$. Then our purported isomorphism f would induce an elementary permutation between $(A_n)^s$ and $(B_n)^t$ (or between $(A_n)^t$ and $(B_n)^s$), which is impossible since $A_n \not\cong B_n$. □

By contrast, we have:

Claim 2: Let $X, Y \subset 2^\omega$ be countable and dense. Then $M_X \equiv_\alpha M_Y$.

Proof: We recall that $M_X \equiv_\alpha M_Y$ iff Player II has a winning strategy in the following game $\mathcal{G}(M_X, M_Y, \alpha)$:

Players I and II alternate moves. On Player I's n -th turn, he either plays a pair (a_n, β_n) where $a_n \in M_X$ and β_n is an ordinal with $\alpha > \beta_0 > \dots > \beta_n$, or else he plays a pair (b_n, β_n) , where $b_n \in M_Y$ and β_n is an ordinal with $\alpha > \beta_0 > \dots > \beta_n$. (Really Player I should also specify which of M_X, M_Y he is playing in, but no harm is caused by suppressing this). On Player II's n -th turn, she plays either $b_n \in M_Y$ or $a_n \in M_X$, depending on Player I's move; she is required to make sure that $(a_0, \dots, a_n) \mapsto (b_0, \dots, b_n)$ is partial elementary from M_X to M_Y . This specifies the game, since Player I cannot survive indefinitely.

Now for each $n < \omega$, we are assuming that $A_n \equiv_{\alpha_n} B_n$, where $(\alpha_n : n < \omega)$ is increasing (possibly not strictly), such that $\alpha = \sup\{\alpha_n + 1 : n < \omega\}$. Fix a winning strategy Γ_n for Player II in the game $\mathcal{G}(A_n, B_n, \alpha_n)$. Given $s \in S_n^0, t \in S_n^1$, let $\Gamma_{s,t} = \Gamma_{t,s}$ be the corresponding strategy for the game $\mathcal{G}((A_n)^s, (B_n)^t, \alpha_n)$. For $s, t \in S_n^0, (A_n)^s \cong (A_n)^t$; use this to get $\Gamma_{s,t}$, a winning strategy for Player II in the game $\mathcal{G}((A_n)^s, (A_n)^t, \infty)$. Similarly define $\Gamma_{s,t}$ for $s, t \in S_n^1$.

We now describe a winning strategy Γ for Player II in the game $\mathcal{G}(M_X, M_Y, \alpha)$.

Case 1: suppose Player I plays (a_0, β_0) where $a_0 = (\eta^*, 0)$ for some $\eta \in X$. Choose n large enough so that $\alpha_n \geq \beta_0$. Using the density of Y , choose $\tau \in Y$ such that $\tau \upharpoonright_n = \eta \upharpoonright_n$. By back-and-forth, we can choose a tree isomorphism $F : (2^{<\omega} \cup X, \subseteq) \cong (2^{<\omega} \cup Y, \subseteq)$ such that F is the identity on 2^n . This map F induces a tree isomorphism $F^* : S_X^* \rightarrow S_Y^*$ defined by $F^*(s^*) = F(s)^*$. On the first move, Player II plays $(\tau^*, 0)$.

On subsequent moves:

If Player I plays $((\nu^*, 0), \beta)$ where $\nu \in X$, then Player II plays $(F^*(\nu^*), 0)$.

If Player I plays $((\nu^*, 0), \beta)$ where $\nu \in Y$, then Player II plays $((F^*)^{-1}(\nu^*), 0)$.

If Player I plays $((\nu, k), \beta)$, where $(\nu, k) \in M_X(s)$ for some $s \in S_m^0 \cup S_m^1$, then Player II plays according to $\Gamma_{s, F^*(s)}$, where we take as input all the previous moves that took place in $M_X(s)$ and $M_Y(F^*(s))$. This will be valid, since either $m \leq n$, in which case $\Gamma_{s, F^*(s)}$ actually describes an isomorphism, or else $m > n$, and so the ordinals involved in the relevant previous moves will all be less than $\beta_0 \leq \alpha_n$.

If Player I plays $((\nu, k), \beta)$, where $(\nu, k) \in M_Y(s)$ for some $s \in S_n^0 \cup S_n^1$, then Player II plays according to $\Gamma_{(F^*)^{-1}(s), s}$, where we take as input all the previous moves that took place in $M_X((F^*)^{-1}(s))$ and $M_Y(s)$.

Case 2: Suppose Player I plays (a_0, β_0) where $a_0 \in M_X(s)$ for some $s \in S_N^0 \cup S_N^1$. Choose $n \geq N$ such that $\alpha_n \geq \beta_0$. By back-and-forth, we can choose a tree isomorphism $F : (2^{<\omega} \cup X, \subseteq) \cong (2^{<\omega} \cup Y, \subseteq)$ such that F is the identity on 2^n . From F we obtain $F^* : S_X^* \rightarrow S_Y^*$ as in Case 1. On the first move, Player II plays according to $\Gamma_{s, s}$, and afterwards plays as in Case 1.

The remaining cases where Player I starts in M_Y are the same, just interchange the roles of X and Y . □

With both claims finished, let $X \subset 2^\omega$ be the set of sequences which are eventually zero, and $Y \subset 2^\omega$ be the set of sequences which are eventually one. Then $M_X \equiv_\alpha M_Y$ and $M_X \not\cong M_Y$. This completes the induction and the proof. \square

This gives the first known example of the following behavior:

Corollary 5.7. *There is a complete first-order theory for whom isomorphism is neither Borel nor Borel complete.*

Here, the example is $\text{REF}(\text{bin})$, the paradigmatic example of a superstable, non- ω -stable theory. Thus we might informally expect this behavior to be extremely common for such theories. Since isomorphism is not Borel, we cannot truly consider $\text{REF}(\text{bin})$ to be especially tame. However, the theory is relatively simple in the sense that it cannot code much infinitary behavior. We end with the following class of examples which follow naturally from this one:

Corollary 5.8. *For any α with $2 \leq \alpha < \omega_1$, there is a complete first-order theory S_α whose isomorphism relation is not Borel, and where $T_\beta \leq_B S_\alpha$ if and only if $\beta \leq \alpha$.*

Each of these theories is grounded, superstable, but not ω -stable.

Proof. Take $S_2 := \text{REF}(\text{bin})$. We construct $S_{\alpha+1}$ as $J(S_\alpha)$, and for limit α , construct S_α as $\prod_{\beta} S_\beta$ as β varies below α . That $T_\alpha \leq_B S_\alpha$ follows from induction and Propositions 3.17 and 3.19, part (4). That $T_{\alpha+1} \not\leq_B S_\alpha$ follows from the fact that $\|S_\alpha\| = \|T_\alpha\| < \|T_{\alpha+1}\|$, which follows from part (3) of those Propositions.

Groundedness follows from part (2), and the place in the stability spectrum is a standard type-counting argument, beginning with the fact that $\text{REF}(\text{bin})$ has the desired properties. \square

We end this subsection with an open question:

Question 5.9. Let α be 0 or 1. Is there a first-order theory S_α whose isomorphism relation is not Borel, and where $T_\beta \leq_B S_\alpha$ if and only if $\beta \leq \alpha$?

Note that the instance of the above question for $\alpha = 0$ is precisely Vaught's conjecture for first-order theories. (A theory T has a perfect set of nonisomorphic models if and only if $T_1 \leq_B T$.) For $\alpha = 1$, abelian p -groups are an infinitary counterexample; we would like a first-order counterexample.

5.2 Infinite Branching

We now turn our attention to $\text{REF}(\text{inf})$ specifically, and prove the following theorem:

Theorem 5.10. *$\text{REF}(\text{inf})$ is Borel complete. Indeed, for each infinite cardinal λ , $\text{REF}(\text{inf})$ is λ -Borel complete.*

Proof. Let Φ be the $L_{\omega_1\omega}$ sentence in the language $\{\leq\}$ describing ω -trees. By Theorem 3.11 of [15], Φ is λ -Borel complete for each λ , so it is enough to produce a λ -Borel reduction f from $\text{Mod}_\lambda(\Phi)$ to $\text{Mod}_\lambda(\text{REF}(\text{inf}))$.

Call a subtree $S \subset \lambda^{<\omega}$ is *reasonable* if for every element $s \in S$, $\{\alpha < \lambda : s \frown (\alpha) \notin S\}$ is infinite. We describe an operation $S \mapsto M_S$ sending reasonable subtrees of $\lambda^{<\omega}$ of size λ , to models of $\text{REF}(\text{inf})$ of size λ , such that $S \equiv_{\infty\omega} S'$ iff $M_S \equiv_{\infty\omega} M_{S'}$. It will then be routine to define a λ -Borel map $f : \text{Mod}_\lambda(\Phi) \rightarrow \text{Mod}_\lambda(\text{REF}(\text{inf}))$, such that given $S' \in \text{Mod}_\lambda(\Phi)$ there is some subtree $S \subset \lambda^{<\omega}$ reasonable with $S \cong S'$ and $f(S') \cong M_S$. Then f will be the desired reduction.

Let $I \subset \lambda^\omega$ be the set of all ω -sequences from λ which are eventually zero. For any set M satisfying

$$I \times \{0\} \subseteq M \subseteq I \times \{0, 1\}$$

if we construe M as an $L = \{E_n : n \in \omega\}$ -structure by the rule $E_n((\eta, i), (\nu, j))$ if and only if $\eta \upharpoonright_n = \nu \upharpoonright_n$, then M is a model of $\text{REF}(\text{inf})$.

So, given a reasonable subtree $S \subset \lambda^{<\omega}$ of size λ , let M_S be the L -structure whose universe is

$$(I \times \{0\}) \cup \{(\eta, 1) : t \frown (1) \subset \eta \text{ for some } t \in \lambda^{<\omega} \setminus S\}$$

We check that the mapping $S \mapsto M_S$ works.

To see this, we describe an inverse operation. Given any L -structure M whose universe satisfies $I \times \{0\} \subseteq M \subseteq I \times \{0, 1\}$, let

$$\text{Tr}(M) = \{s \in \lambda^{<\omega} : \forall \alpha < \lambda \exists \eta \in \lambda^\omega [s \frown (\alpha) \subset \eta \text{ and } (\eta, 1) \notin M]\}$$

We first argue that for any subtree $S \subseteq \lambda^{<\omega}$, we have $\text{Tr}(M_S) = S$. Indeed, suppose $s \in S$. Choose $\alpha \in \lambda$ arbitrarily. Then the element $\eta := s \frown (\alpha) \frown \bar{0}$ of I witnesses that $s \in \text{Tr}(M_S)$. Conversely, if $s \notin S$ then as $(\eta, 1) \in M_S$ for every $\eta \supset s \frown (1)$, $s \notin \text{Tr}(M_S)$.

Thus, in particular, $\text{Tr}(M_S)$ is a subtree of $\lambda^{<\omega}$ whenever S is.

Claim: For any subtrees S, T of $\lambda^{<\omega}$, if the L -structures $M_S \equiv_{\infty, \omega} M_T$, then $(S, \subseteq) \equiv_{\infty, \omega} (T, \subseteq)$.

Proof: Assume $M_S \equiv_{\infty, \omega} M_T$. Pass to a forcing extension $\mathbb{V}[G]$ in which $\lambda^\mathbb{V}$ is countable. Choose an L -isomorphism $f : M_S \rightarrow M_T$. This induces a tree isomorphism $f^* : (\text{Tr}(M_S), \subseteq) \rightarrow (\text{Tr}(M_T), \subseteq)$. Combined with the computation above, (S, \subseteq) and (T, \subseteq) are isomorphic in $\mathbb{V}[G]$, so they are back-and-forth equivalent in \mathbb{V} . \square

To complete the proof, suppose two reasonable subtrees satisfy $(S, \subseteq) \equiv_{\infty\omega} (T, \subseteq)$. Pass to a forcing extension wherein λ is countable, so that $S \cong T$. Then, since S and T are reasonable, we can choose a tree automorphism $f : (\lambda^{<\omega}, \subseteq) \cong (\lambda^{<\omega}, \subseteq)$ that carries S to T . Then clearly f induces an L -isomorphism from M_S to M_T . This implies that the L -structures M_S and M_T are back-and-forth equivalent in the ground model. \square

The following Corollary follows immediately from Corollary 3.14, Proposition 5.1, and Theorem 5.10.

Corollary 5.11. *REF(inf) is not short. Indeed, REF(inf) has class-many $\equiv_{\infty\omega}$ -inequivalent models in \mathbb{V} .*

6 ω -Stable Examples

Here we discuss two more first-order theories whose isomorphism relations are not Borel, but where one is Borel complete, and the other does not embed T_3 . Interestingly, both are extremely similar model-theoretically. Both are ω -stable with quantifier elimination, and have ENI-NDOP and eni-depth 2, which together give a strong structure theorem in terms of stability theory.³ Indeed, as we will show, both have exactly \beth_2 models up to $\equiv_{\infty\omega}$, meaning that at a “macro” level, they are extremely similar. Yet at a “micro” level, they are quite different. Because of the similarity of the examples, we are able to highlight exactly why one is relatively simple, while the other is not.

Let us define the theories. The first, K , is due to Koerwien and constructed in [14]. The language has unary sorts U , V_i , and C_i , as well as unary functions S_i and π_i^j for $i \in \omega$ and $j \leq i + 1$. The axioms are as follows:

- The sorts U , V_i , and C_i are all disjoint. U and each of the V_i are infinite, but each C_i has size 2.
- π_i^{i+1} is a function from V_i to U ; π_i^j is a function from V_i to C_j when $j \leq i$.
- For each tuple $\bar{c} = (c_0, \dots, c_i)$ and each $u \in U$, $\pi_i^{-1}(\bar{c}, u)$ is nonempty. Here π_i refers to the product map $\pi_i^0 \times \dots \times \pi_i^{i+1} : V_i \rightarrow C_0 \times \dots \times C_i \times U$.
- S_i is a unary successor function from V_i to itself, and $\pi_i \circ S_i = \pi_i$.

We have a few remarks. Typically we will drop the subscript on π_i and S_i if it is clear from context. There is a slight ambiguity about the sorts, whether one works in traditional first-order logic (and thus there may be “unsorted” elements) or in multisorted logic (where there will not be). Since the unsorted elements never have any effect other than to complicate notation, we work in multisorted logic.

The properties of K have been well studied by Koerwien in [14]; we summarize his findings here:

Theorem 6.1. *K is complete with quantifier elimination. It is ω -stable, has ENI-NDOP, and is eni-shallow of eni-depth 2. Furthermore, the isomorphism relation for K is not Borel.*

Our other theory is a tweak of K , so we call it TK . The language is slightly different; we have unary sorts U , V_i , and C_i as before, but have unary functions S_i , π_i^0 , π_i^1 , and τ_{i+1} for $i \in \omega$. The axioms are as follows:

- The sorts U , V_i , and C_i are all disjoint. U and each of the V_i are infinite, but each C_i has size 2^i .
- τ_{i+1} is a surjection from C_{i+1} to C_i where, for all $c \in C_i$, $|\tau_{i+1}^{-1}(c)| = 2$.
- π_i^1 is a function from V_i to U ; π_i^0 is a function from V_i to C_i .

³In [15], an attempt is made to characterize which first-order ω -stable theories are Borel complete, using the dividing lines: ENI-DOP vs ENI-NDOP, and eni-deep vs eni-shallow. In particular, it is shown that any ω -stable theory which either has ENI-DOP or is eni-deep is Borel complete; and if an ω -stable theory has both ENI-NDOP and is eni-shallow, then it has fewer than \beth_{ω_1} -many models up to back-and-forth equivalence.

- For each tuple $c \in C_i$ and each $u \in U$, $\pi_i^{-1}(c, u)$ is nonempty. Here π_i refers to the product map $\pi_i^0 \times \pi_i^1 : V_i \rightarrow C_i \times U$.
- S_i is a unary successor function from V_i to itself, and $\pi_i \circ S_i = \pi_i$.

The preceding notes also apply to K. The behavior is extremely similar, and essentially the same proofs of basic properties of K apply to TK. We summarize this now:

Theorem 6.2. *TK is complete with quantifier elimination. It is ω -stable, has ENI-NDOP, and is eni-shallow of eni-depth 2.*

We can easily see that both K and TK have relatively few models up to back-and-forth equivalence:

Proposition 6.3. $I_{\infty, \omega}(K) = I_{\infty, \omega}(TK) = \beth_2$.

Indeed, every model M of either theory has a submodel N where $M \equiv_{\infty \omega} N$ and $|N| \leq \beth_1$.

Proof. Let T be either K or TK. For the proof of the proposition we can restrict attention to models of T with a fixed algebraic closure of the empty set $\bigcup_i C_i$. If $T = K$, then let C be all finite sequences $(a_j : j < i)$ with $i > 0$ and with each $a_j \in C_j$; if $T = TK$ then let $C = \bigcup_i C_i$.

We first show $I_{\infty, \omega}(T) \geq \beth_2$. For each $\eta \in 2^\omega$, let u_η be some element which will eventually be part of U in some model of T . For any $n \in \omega$ and any $c \in C$ where $\pi_n^{-1}(c, u_\eta)$ is nonempty, we insist the S_n -dimension of $\pi_n^{-1}(c, u_\eta)$ be infinite if $\eta(n) = 1$, or equal to one otherwise. (If $T = K$, then $\pi_n^{-1}(c, u_\eta)$ is nonempty if and only if $\text{lg}(c) = n$; if $T = TK$ then $\pi_n^{-1}(c, u_\eta)$ is nonempty if and only if $c \in C_n$.) For any infinite $X \subseteq 2^\omega$, define M_X to have $U^{M_X} = \{u_\eta : \eta \in X\}$ with the described behavior of the V_i and S_i . Evidently if $Y \subseteq 2^\omega$ is infinite and $X \neq Y$, then for any $\eta \in X \setminus Y$, there is no $\nu \in Y$ where $(M_X, u_\eta) \equiv_{\infty \omega} (M_Y, u_\nu)$, and symmetrically. Thus, $M_X \not\equiv_{\infty \omega} M_Y$. Since there are \beth_2 infinite subsets of 2^ω , $I_{\infty \omega}(T) \geq \beth_2$.

That $I_{\infty \omega}(T) \leq \beth_2$ follows immediately from the second claim. So let M be some model of T , of any particular cardinality. We begin by stripping down the V_i . For every $u \in U$ and $c \in C$, if $\pi^{-1}(c, u)$ is uncountable, drop all but a countable S -closed subset of infinite S -dimension. Do this for all pairs (c, u) . The result is $\equiv_{\infty \omega}$ -equivalent to the original by an easy argument, and $\pi^{-1}(c, u)$ is now always countable.

Next we need only ensure that U has size at most continuum. So put an equivalence relation E on U , where we say uEu' holds if and only if, for all $c \in C$, the dimensions of $\pi^{-1}(c, u)$ and $\pi^{-1}(c, u')$ are equal. If any E -class is uncountable, drop all but a countably infinite subset; the resulting structure is $\equiv_{\infty \omega}$ -equivalent to the original again. Further, each E -class is now countable, and there are only $|C^\omega| = \beth_1$ possible E -classes, so the structure now has size at most \beth_1 . This completes the proof. \square

Any additional complexity of either theory comes from elementary permutations of the algebraic closure of the empty set. In any model M of either K or TK, $\text{acl}_M(\emptyset) = \bigcup_{i \in \omega} C_i(M)$. In models M of TK, the projection functions $\{\tau_i\}$ naturally induce a tree structure, so we think of $\text{acl}_M(\emptyset)$ as being a copy of $(2^{<\omega}, \leq)$. In models M of K, as each $C_i(M)$ has exactly two elements, so one can think of $\text{acl}_M(\emptyset)$ as being indexed by $2 \times \omega$.

Note, however, there is some freedom with all this; for our purposes, $\text{acl}_M(\emptyset)$ could equally well be viewed as any subset of $\text{acl}_{M^{eq}}(\emptyset)$ whose definable closure contains $\text{acl}_M(\emptyset)$ (here M^{eq} is the result of eliminating imaginaries from M). In the case $M \models K$ it is most convenient to say that $\text{acl}_M(\emptyset)$ is all finite sequences $\langle a_j : j < i \rangle$, where each $a_j \in C_j(M)$. These finite sequences, when ordered by initial segment, also give a natural correspondence of $\text{acl}_M(\emptyset)$ with the tree $(2^{<\omega}, \leq)$. Henceforth, when discussing models M of either K or TK, we will view $\text{acl}_M(\emptyset)$ as being indexed by the tree $(2^{<\omega}, \leq)$.

Next, we discuss the group G of elementary permutations of $\text{acl}_M(\emptyset)$ (which only depends on the theory). For K, the relevant group is $(2^\omega, \oplus)$, the direct product of ω copies of the two-element group. Indeed, in any model of K, any elementary permutation of $\text{acl}_M(\emptyset)$ is determined by the sequence of permutations of $C_i(M)$. In TK, as elementary permutations just have to respect the τ_i structure, the relevant group of elementary permutations is $\text{Aut}(2^{<\omega}, \leq)$. Both of these groups are compact (in fact this is true for all first order theories), but only the group for K is abelian. It turns out that being abelian is enough to produce relative simplicity, while being nonabelian leaves enough room to allow TK to be Borel complete.

For the next proposition we need some setup.

Let X be the Polish space of all $f : 2^{<\omega} \rightarrow (\omega + 1 \setminus \{\emptyset\})$. Let T be either K or TK, and let G be either $(2^\omega, \oplus)$ or $\text{Aut}(2^{<\omega}, \leq)$, respectively. G acts on $2^{<\omega}$ naturally: if $G = (2^\omega, \oplus)$, then $g \cdot \sigma = g \upharpoonright_{|\sigma|} \oplus \sigma$. If $G = \text{Aut}(2^{<\omega}, \leq)$, then $g \cdot \sigma = g(\sigma)$. From this we get an action of G on X : namely for $f \in X, g \in G$, $(g \cdot f)(\sigma) = f(g^{-1} \cdot \sigma)$. This is a strongly definable, continuous action in the sense of Definition 4.1. Let E_G be the equivalence relation on X induced by the action, as well as the equivalence relation on $\mathcal{P}_{\aleph_1}(X)$ induced by the diagonal action.

Now G acts diagonally on X^ω also; this action commutes with the permutation action of S_∞ on X^ω . So $G \times S_\infty$ acts naturally on X^ω ; let $E_{G \times S_\infty}$ be the equivalence relation induced by this action.

Proposition 6.4. *Let T be either K or TK. Then:*

- $(\text{Mod}(T), \cong) \sim_B (X^\omega, E_{G \times S_\infty})$.
- $(\text{Mod}(T), \cong) \sim_{HC} (\mathcal{P}_{\aleph_1}(X), E_G)$.

Proof. For the various codings below, fix a pairing function $\langle \cdot, \cdot \rangle : (\omega + 1 \setminus \emptyset)^2 \rightarrow (\omega + 1 \setminus \emptyset)$. Note that one difference between K and TK that frequently affects the coding is: $\pi^{-1}(\emptyset, u)$ is only nonempty for models of TK.

To show $(\text{Mod}(T), \cong) \leq_B (X^\omega, E_{G \times S_\infty})$, first let $M \in \text{Mod}(T)$ be arbitrary. We may choose an indexing of $\text{acl}_M(\emptyset)$ by $2^{<\omega}$, and of U^M by ω , using the original indexing of the universe of M by ω . Then each element $u \in U^M$ induces a function $c_u \in X$, where $c_u(\sigma)$ is the S -dimension of $\pi^{-1}(\sigma, u)$. (In the case of $T = K$, define $c_u(\emptyset) = 1$.) Then take M to the sequence $(c_{u_n} : n \in \omega)$, where u_n is the n -th element of U . It is clear that this works.

To show $(X^\omega, E_{G \times S_\infty}) \leq_B (\text{Mod}(T), \cong)$, fix a sequence $\bar{x} = (x_n : n \in \omega)$. We describe the case for the theory K . We define $M_{\bar{x}}$ to have $U^{M_{\bar{x}}} = \{4n : n \in \omega\}$, and have $C_i^{M_{\bar{x}}} = \{4i+1, 4i+2\}$. Then, using the infinitely many remaining elements, we arrange that for each $\sigma \in 2^{<\omega} \setminus \emptyset$ and each $n < \omega$, the S -dimension of $\pi^{-1}(\sigma_*, 4n)$ is $\langle x_n(\sigma), x_n(\emptyset) \rangle$, where $\sigma_* = (4i+1 + \sigma(i) : i < \text{lg}(\sigma))$. The case for TK is similar.

We have shown that $(\text{Mod}(T), \cong) \sim_B (X^\omega, E_{G \times S_\infty})$. It follows that they are \leq_{HC} -biembeddable; so to conclude the proof of the proposition it suffices to show that $(X^\omega, E_{G \times S_\infty}) \sim_{\text{HC}} (\mathcal{P}_{\aleph_1}(X), E_G)$.

To show $(X^\omega, E_{G \times S_\infty}) \leq_{\text{HC}} (\mathcal{P}_{\aleph_1}(X), E_G)$, we just need to handle multiplicities. So fix $\bar{x} = (x_n : n \in \omega) \in X^\omega$, and for each n , define $m_{\bar{x}}(n)$ to be $|\{m : x_m = x_n\}|$. For each $n \in \omega$ let $y_n \in X$ be defined by $y_n(\sigma) = \langle x_n(\sigma), m_{\bar{x}}(n) \rangle$. Then $\bar{x} \mapsto \{y_n : n < \omega\}$ works.

We define a reverse embedding $f : (\mathcal{P}_{\aleph_1}(X), E_G) \leq_{\text{HC}} (X^\omega, E_{G \times S_\infty})$ as follows (where recall that we do not require f to be single-valued). Namely, given $A \subset X$ countable and given $\bar{x} \in X^\omega$, put $(A, \bar{x}) \in f$ whenever \bar{x} is an infinite-to-one enumeration of A . Also put $(\emptyset, \bar{x}) \in f$ for some fixed injective $\bar{x} \in X^\omega$. \square

6.1 Koerwien's Example

For this subsection, we want to prove that K is not Borel complete, and further, to characterize exactly which T_α embed in K . To do this, we show directly that $T_2 \leq_B K$, and then show that $\|K\|_{\text{ptl}} = \beth_2$. The former is quite straightforward:

Proposition 6.5. $T_2 \leq_B K$.

Proof. Let $X \subset 2^\omega$ be countable; we describe a model $M_X \models K$ from which X can be easily recovered. Let U be the set $A \cup X$, where A is some countable infinite set which is disjoint from X . Let $C_i = \{c_0^i, c_1^i\}$. For each tuple (a, \bar{c}) with $a \in A$, arrange that $\pi^{-1}(a, \bar{c})$ has S -dimension 1. For each tuple (x, \bar{c}) with $x \in X$, arrange that $\pi^{-1}(x, \bar{c})$ has S -dimension $x(|\bar{c}|) + 2$. Clearly $M_X \cong M_Y$ iff $X = Y$.

Now it is not hard, given $\bar{x} = (x_n : n \in \omega) \in (2^\omega)^\omega$, to produce in a Borel fashion a model $M_{\bar{x}} \models K$ with universe ω , such that $M_{\bar{x}} \cong M_{\{x_n : n \in \omega\}}$. This gives a Borel reduction from $((2^\omega)^\omega, F_2)$ to $(\text{Mod}(K), \cong)$, which suffices (see the discussion preceding Theorem 5.3). \square

Having accomplished this, we can state everything we need about K :

Theorem 6.6. $\|K\| = \beth_2$. Therefore, K is not Borel complete; indeed, there is no Borel embedding of T_3 into $\text{Mod}(K)$.

Proof. That $\|K\| \geq \beth_2$ follows immediately from Proposition 6.3. Since $G = (2^\omega, \oplus)$ is compact and abelian, $\|(\mathcal{P}_{\aleph_1}(X), E_G)\| \leq \beth_2$ by Theorem 4.8. As $\|K\| = \|(\mathcal{P}_{\aleph_1}(X), E_G)\|$ by Proposition 6.4, we conclude that $\|K\| = \beth_2$. That there is no Borel embedding of T_3 into $\text{Mod}(K)$ is immediate from Theorem 3.10(2) and Corollary 3.21. \square

Once we have one such example, we can apply the usual constructions to get a large class of ω -stable examples:

Corollary 6.7. *For each ordinal α with $2 \leq \alpha < \omega_1$, there is an ω -stable theory S_α whose isomorphism relation is not Borel and where $T_\beta \leq_B S_\alpha$ if and only if $\beta \leq \alpha$.*

Proof. Let S_2 be K . Then proceed inductively as in Corollary 5.8. □

There is no such example when $\alpha = 0$, since Vaught's Conjecture holds for ω -stable theories. (T has a perfect set of nonisomorphic models iff $T_1 \leq_B T$, and so whenever T is ω -stable, either $T \leq_B T_0$ or $T_1 \leq_B T$.) It is unknown if there is an example when $\alpha = 1$.

6.2 A New ω -Stable Theory

We now consider TK , with the aim of showing it is Borel complete. Indeed with Proposition 6.4 we have already shown $(\text{Mod}(TK), \cong)$ is Borel equivalent to $(X^\omega, E_{G \times S_\infty})$, where X is the space of all $c : 2^{<\omega} \rightarrow \omega$. (We are replacing $\omega + 1 \setminus \emptyset$ with ω , which is harmless.) Recall that $G = \text{Aut}(2^{<\omega}, \leq)$ acts on X by permuting the fibers; that is, for any $c : 2^{<\omega} \rightarrow \omega$, any $g \in G$, and any $\sigma \in 2^{<\omega}$, $(g \cdot c)(\sigma) = c(g^{-1} \cdot \sigma)$. Then G acts on X^ω diagonally, while S_∞ acts on X^ω by permuting the fibers, so these actions commute with one another and induce an action of the product group $G \times S_\infty$.

Thus, to show TK is Borel complete, it is enough to show $(X^\omega, E_{G \times S_\infty})$ is Borel complete, which we do directly.

Theorem 6.8. $(\text{Graphs}, \cong) \leq_B (X^\omega, E_{G \times S_\infty})$.

Proof. To simplify notation, for the whole of this proof we write E in place of $E_{G \times S_\infty}$. We need some setup first. Observe that G naturally acts on 2^ω , the set of branches of $(2^{<\omega}, \leq)$, by $g \cdot \sigma = \bigcup_n g \cdot \sigma \upharpoonright_n$; this is a well-defined sequence precisely because g is a tree automorphism. Let $\{D_i : i \in \omega\}$ be a countable set of dense, disjoint, countable subsets of 2^ω , and let $D = \bigcup_i D_i$. We need one claim, where we use the relative complexity of G (it would not go through if we replaced TK with K):

Claim: For any $\sigma \in S_\infty$, there is a $g \in G$ where for all $i \in \omega$, $g \cdot D_i = D_{\sigma(i)}$ as sets.

Proof: We construct such a g by a back-and-forth argument. So let \mathcal{F} be the set of finite partial functions from D to itself, satisfying all the following:

- For each $f \in \mathcal{F}$ and each $\eta \in \text{dom}(f)$, if $\eta \in D_i$, then $f(\eta) \in D_{\sigma(i)}$.
- For each $f \in \mathcal{F}$ and each $\eta, \nu \in \text{dom}(f)$, $\text{lg}(\eta \wedge \nu) = \text{lg}(f(\eta) \wedge f(\nu))$, where $\eta \wedge \nu$ denotes the longest common initial segment of η and ν .
- The previous conditions, but with f^{-1} and σ^{-1} instead of f and σ .

Once we establish that \mathcal{F} is a back-and-forth system, then \mathcal{F} defines a $g \in G$ with the desired property. For choose a bijection $f : D \rightarrow D$ such that the finite restrictions of f all lie in \mathcal{F} . If $s \in 2^n$, let $g(s)$ be $f(\eta) \upharpoonright_n$ for any η extending s ; because of the consistency properties of \mathcal{F} , and because D is dense, this is well-defined. Then clearly $g \in G$ has the desired property with respect to σ .

So we need only show that \mathcal{F} is a back-and-forth system. Of course the empty function is in \mathcal{F} . So say $f \in \mathcal{F}$ and $\eta \in 2^\omega$; we want $f' \supset f$ in \mathcal{F} with $\eta \in \text{dom}(f')$. The case where f is empty is easy, so suppose f is nonempty. We also assume $\eta \notin \text{dom}(f)$ already. Let n be maximal among $\{\text{lg}(\eta \wedge \nu) : \nu \in \text{dom}(f)\}$, and let $\nu \in \text{dom}(f)$ be such that $\text{lg}(\eta \wedge \nu) = n$. We then pick an element $f(\eta)$ of 2^ω which extends $f(\nu) \upharpoonright_n \frown (1 - f(\nu)(n))$. That is, $f(\eta)$ agrees with $f(\nu)$ before stage n , but disagrees with it at n . If $\eta \in D_i$, choose this element from $D_{\sigma(i)}$, which is possible by density. This clearly satisfies the desired properties, and the other direction is symmetric, proving the claim. \square

Given $\eta, \tau \in 2^\omega$ and $k \in \{1, 2, 3\}$, let $c_{\eta, \tau}^k : 2^{<\omega} \rightarrow \omega$ be the coloring which sends $s \in 2^{<\omega}$ to k , if $s \subset \eta$ or $s \subset \tau$, or 0 otherwise. Also, fix a bijection $\rho : \omega \rightarrow \bigcup_{i \leq j} D_i \times D_j$. We have now fixed enough notation and can describe our map $f : \text{Graphs} \rightarrow X^\omega$.

Let R be a graph on ω – that is, R is a binary relation on ω which is symmetric and irreflexive. For each $n \in \omega$, $\rho(n)$ is a pair $(\eta, \tau) \in D_i \times D_j$ for some $i \leq j$. If $i = j$, define $c_n = c_{\eta, \tau}^1$. If $i < j$ and $\{i, j\} \in R$, then let $c_n = c_{\eta, \tau}^2$. Otherwise let $c_n = c_{\eta, \tau}^3$. Then put $f(R) := (c_n : n \in \omega)$. $f(R)$ is a visibly element of X^ω , and clearly f is Borel. Note also that f is injective.

Suppose $\sigma : (\omega, R) \cong (\omega, R')$ is a graph isomorphism. We show that $f(R)Ef(R')$. By the claim, there is a $g \in G$ where for all $i \in \omega$, $g \cdot D_i = D_{\sigma(i)}$. Let A be the range of $f(R)$ and let A' be the range of $f(R')$. We show that $g \cdot A = A'$ setwise. First suppose $c_{\eta, \tau}^1 \in A$. Let i be such that $\eta, \tau \in D_i$, so $g(\eta), g(\tau) \in D_{\sigma(i)}$. Then $g \cdot c_{\eta, \tau}^1 = c_{g(\eta), g(\tau)}^1 \in A'$. Similarly if $c_{\eta, \tau}^2 \in A$, there is some $i < j$ where $\eta \in D_i$, $\tau \in D_j$, and $\{i, j\} \in R$. Since $\sigma : (\omega, R) \rightarrow (\omega, R')$ is a graph isomorphism, $\{\sigma(i), \sigma(j)\} \in R'$, so $c_{g(\eta), g(\tau)}^2 \in A'$ (this uses $c_{g(\eta), g(\tau)}^2 = c_{g(\tau), g(\eta)}^2$). The case $c_{\eta, \tau}^3 \in A$ is the same. Thus $g \cdot A \subset A'$; by a symmetric argument $g \cdot A = A'$. Since $g \cdot f(R)$ and $f(R')$ are both injective and they have the same range, some permutation of $g \cdot f(R)$ is equal to $f(R')$, i.e. $f(R)Ef(R')$.

It only remains to show that if $f(R)Ef(R')$, then $(\omega, R) \cong (\omega, R')$. So suppose $f(R)Ef(R')$. Let A be the range of $f(R)$ and let A' be the range of $f(R')$, and choose $g \in G$ such that $g \cdot A = A'$. Let $i < \omega$; then since for all $\eta, \tau \in D_i$, $c_{g(\eta), g(\tau)}^1 \in A'$, we have that $g \cdot D_i = D_{\sigma(i)}$ for some $\sigma(i) < \omega$. I claim that $\sigma : (\omega, R) \cong (\omega, R')$. Indeed, for $i < j$, $(i, j) \in R$ iff there are $\eta \in D_i$, $\tau \in D_j$ with $c_{\eta, \tau}^2 \in A$, which is the case iff there are $\eta \in D_{\sigma(i)}$, $\tau \in D_{\sigma(j)}$ with $c_{\eta, \tau}^2 \in A'$, which is the case iff $(\sigma(i), \sigma(j)) \in R'$. \square

We have now shown:

Theorem 6.9. *TK is Borel complete.*

Proof. By Theorem 6.8, together with the fact that graphs are Borel complete. \square

This resolves a few open questions, raised in [15]:

Corollary 6.10. *The ω -stable theory TK is Borel complete, but does not have ENI-DOP and is not eni-deep. Indeed TK is not λ -Borel complete for any λ with $2^\lambda > \beth_2$, as $|CSS(TK)_{\text{sat}}| = \beth_2$.*

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