MATH/CMSC 456 (Washington) Exam 2 May 2, 2006

1. (a) Given any $x < 123456$, we have $H(x) = x$, so H is not preimage resistant. Also, $H(123457) = H(1)$, so H is not strongly collision free.
(b) The line through $(1,3)$ and $(-2,0)$ has equation $y \equiv x + 2$. Intersecting with the curve yields $(x+2)^2 \equiv x^3 + 8$, so $0 \equiv x^3 - x^2 + \cdots$. The sum of the roots is $1 + (-2) + x \equiv 1$, so $x \equiv 2$. Then $y \equiv x + 2 \equiv 4$. Reflecting yields the answer $(2, -4)$, or $(2, 7)$.

2. The remaining steps are:

4. Victor randomly chooses $i = 1$ or $i = 2$ and asks Peggy for r_i.
5. Peggy sends r_i.
6. Victor checks that $r_i^e \equiv x_i$.
7. They repeat steps 1 through 6 seven times.

If Peggy does not know m, the probability is $1/2$ that Peggy can correctly supply r_i in a given round. Therefore, the probability is $(1/2)^k$ that Peggy can succeed for k rounds if she doesn't know m. When $k = 7$, this probability is less than .01, so if Peggy succeeds for seven rounds then the probability is more than 99% that she knows m.

3. Collisions can be found by a birthday attack. Make a list of $H(x)$ for around 2^{30} (maybe a little more) random values of x. Since the length of the list is approximately $\sqrt{2^{60}}$, there is a good chance that two hash values are the same. This yields a collision.

4. Bob switches C_0 and C_1, so he inputs C_1, C_0 into the machine. it outputs C_0 as the left half and $C_1 \oplus f(C_0)$ as the right half. But $C_0 = R$ and $C_1 \oplus f(C_0) = (L \oplus f(R)) \oplus f(R) = L$. Therefore, the output is RL. Switch the two halves to get LR.

5. (a) Look for a match between the two lists. If there is, then $\beta \equiv \alpha^{j+kN}$, so $x = j + kN$ solves the discrete log problem. This always works because we can write $x = x_0 + x_1 N$. When $j = x_0$ and $k = x_1$, we have a match. This means that there is a match between the two lists.
(b) Make two lists:

1. jA for $0 \le j < N$
2. $B - kNA$ for $0 \le k < N$. Look for a match between the two lists. When there is a match, we have $B = (j + kN)A$.

6. (a) We need m so that $\alpha^m \equiv \beta^r (r)^s \pmod{p}$. This simplifies to $\alpha^m \equiv \beta^r (\alpha\beta)^{-r} \equiv \alpha^{-r} \pmod{p}$, so we can take $m \equiv -r \pmod{p-1}$.
(b) We need $H(m) \equiv -\alpha\beta$. But $H(m)$ is assumed to be preimage resistant, so it is hard to find m satisfying this property.