Lecture 1: From Data to Graphs, Weighted Graphs and Graph Laplacian

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From Data to Graphs

Datasets Diversity



- Social Networks: Set of individuals ("agents", "actors") interacting with each other (e.g., Facebook, Twitter, joint paper authorship, etc.)
- Communication Networks: Devices (phones, laptops, cars) communicating with each other (emails, spectrum occupancy)
- Transportation networks: Road networks, Air/Sea travel network
- Biological Networks: Macroscale: How animals interact with each other; Microscale: How proteins interact with each other.
- Databases of signals: speech, images, movies; graph relationship tends to reflect signal similarity: the higher the similarity, the larger the weight.
- Chemical networks: chemical compunds; each node is an atom.
- Point Clouds: scanners of 3D objects; aerial LIDAR images.

Databases of Graphs

Public Datasets

Here are several public databases:

- Duke: https://dnac.ssri.duke.edu/datasets.php
- 2 Stanford: https://snap.stanford.edu/data/
- Oni. Koblenz: http://konect.uni-koblenz.de/
- M. Newman (U. Michigan): http://www-personal.umich.edu/ mejn/netdata/
- S A.L. Barabasi (U. Notre Dame): http://www3.nd.edu/ networks/resources.htm
- OUCI: https://networkdata.ics.uci.edu/resources.php
- Google/YouTube: https://research.google.com/youtube8m/
- Ohemical Compounds: http://quantum-machine.org/datasets/
- 9 Princeton Shape Benchmark: https://shape.cs.princeton.edu/benchmark/
- Road networks: http://networkrepository.com/road.php
- USGS: https://catalog.data.gov/dataset/usgs-national-transportation-dataset-ntddownloadable-data-collectionde7d2
- FSU: https://people.sc.fsu.edu/~jburkardt/datasets/cities/cities.html

Graphs

Typology

Two classes of graphs:

- **(**) Geometric graphs: vertices are points in a vector space, e.g. \mathbb{R}^3 ;
- Relational (or abstract) graphs: edges represents similarity or dependency between vertices;

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- **①** Geometric graphs: vertices are points in a vector space, e.g. \mathbb{R}^3 ;
- Relational (or abstract) graphs: edges represents similarity or dependency between vertices;
- **1** Unweighted graphs: edges carry same weight (1).
- Weighted undirected graphs: edges may have different weights; the weight matrix is symmetric.
- Oirected graphs: edges are oriented and the weight matrix is not symmetric.

Weighted Graphs

The main goal of this lecture is to introduce basic concepts of weighted and undirected graphs, its associated graph Laplacian, and geometric graphs.

Graphs (and weights) reflect either similarity between nodes, or functional dependency, or interaction strength.

- SIMILARITY: Distance, similarity between nodes \Rightarrow weight $w_{i,j}$
- PREDICTIVE/DEPENDENCY: How node *i* is predicted by its nighbor node *j* ⇒ weight *w_{i,j}*
- INTERACTION: Force, or interaction energy between atom i and atom j

Definitions $G = (\mathcal{V}, \mathcal{E})$ and $G = (\mathcal{V}, \mathcal{E}, w)$

An undirected graph G is given by two pieces of information: a set of vertices \mathcal{V} and a set of edges \mathcal{E} , $G = (\mathcal{V}, \mathcal{E})$.

A weighted graph has three pieces of information: $G = (\mathcal{V}, \mathcal{E}, w)$, the set of vertices \mathcal{V} , the set of edges \mathcal{E} , and a weight function $w : \mathcal{E} \to \mathbb{R}$.

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 $\begin{aligned} \mathcal{V} &= \{1, 2, 3, 4, 5, 6, 7, 8, 9\} \\ \mathcal{E} &= \{(1, 2), (2, 4), (4, 7), \\ (6, 7), (1, 5), (5, 6), (5, 7), \\ (2, 8), (8, 9)\} \\ 9 \ \textit{vertices} \ , 9 \ \textit{edges} \\ \text{Undirected graph, edges} \\ \text{are not oriented. Thus} \\ (1, 2) &\sim (2, 1). \end{aligned}$

Definitions
$$G = (\mathcal{V}, \mathcal{E})$$

A weighted graph $G = (\mathcal{V}, \mathcal{E}, w)$ can be directed or undirected depending whether w(i, j) = w(j, i). Symmetric weights == Undirected graphs

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Example of a Weighted Graph UCINET IV Datasets: Bernard & Killworth Office

Available online at:

http://vlado.fmf.uni-lj.si/pub/networks/data/ucinet/ucidata.htm Content: Two 40 \times 40 matrices: symmetric (B) and non-symmetric (C) Bernard & Killworth, later with the help of Sailer, collected five sets of data on human interactions in bounded groups and on the actors' ability to recall those interactions. In each study they obtained measures of social interaction among all actors, and ranking data based on the subjects' memory of those interactions. The names of all cognitive (recall) matrices end in C, those of the behavioral measures in B.

These data concern interactions in a small business office, again recorded by an "unobtrusive" observer. Observations were made as the observer patrolled a fixed route through the office every fifteen minutes during two four-day periods. BKOFFB contains the observed frequency of interactions; BKOFFC contains rankings of interaction frequency as recalled by the employees over the two week period.

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Example of a Weighted Graph

```
bkoff.dat (similar for bkfrat.dat)
DL
N=40 NM=2
FORMAT = FULLMATRIX DIAGONAL PRESENT
LEVEL LABELS:
BKOFFB
BKOFFC
DATA:
000010000210
00480330110030002110022000
01002901100300
0402101400010011100200100110
00040101810021
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Example of a Weighted Graph

UCINET IV Datasets: Bernard & Killworth Office

```
...

0 0 1 3 0 0 2 0 0 0 1 0 5 0 0 0 0 0 0 5 0 0 0 0

0 1 1 0 0 0 2 4 5 0 0 0 0 0

0 27 3 36 23 34 14 19 13 9 3 26 21 25 1 8 22 12 11 4 2 37 35 17 5 20

7 33 32 39 38 16 28 30 29 24 6 10 18 31

...

29 38 17 4 31 37 6 35 36 22 17 24 39 20 19 26 12 30 32 28 25 1 18 14 33

34

27 8 9 21 11 10 5 3 2 15 23 16 13 0
```

For bkfrat.dat, a similar structure: Lines 8 to 181 contain the symmetric weight matrix. Note, that each row in W takes three lines in the .dat file (due to file formatting).

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Example of a Geometric Graph



Dataset: Point Clouds from Princeton Shape Benchmark (PSB) https://shape.cs.princeton.edu/benchmark/ File: 2; Ant n=3495 vertices (6667 faces; 61 connected components)

Figure: m2

Movie: See the link



Figure: Graph Embedding

Project 2.1: Euclidian Embeddings and Image Registration Project on Geometric Graphs

Projects deal with optimization problems related to geometric graphs. Input Data: (i) a $n \times n$ symmetric matrix of (noisy) pairwise distances between n vertices; (ii) Data set of n-node graphs. Objectives:

- Find a 3D embedding of the graph (i.e., find coordinates of these points);
- Find the optimal rigid transformation to each of shapes in the data set, compute the residual errors and select the "closest" shape from dataset;
- S Create a movie of the continuous transformation.
- Repeat first two questions when only a subset of the n(n-1)/2 pairwise distances are known.

(see movie)

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Paths

Concept: A path is a sequence of edges where the right vertex of one edge coincides with the left vertex of the following edge. Example:



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$$\{(1,2),(2,4),(4,7),(7,5)\} =$$

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Graph Attributes

Graph Attributes (Properties):

• Connected Graphs: Graphs where any two distinct vertices can be connected through a path.

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A complete graph with *n* vertices has $m = \begin{pmatrix} n \\ 2 \end{pmatrix} = \frac{n(n-1)}{2}$ edges.

Graph Attributes

Example:



Graph Attributes

Example:



- This graph is not connected.
- It is not complete.
- It is the union of two connected graphs.

Metric

Distance between vertices: For two vertices x, y, the distance d(x, y) is the length of the shortest path connecting x and y. If x = y then d(x, x) = 0.

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1 If
$$\forall x, y \in \mathcal{V}$$
, $d(x, y) < \infty$ then $(\mathcal{V}, \mathcal{E})$ is connected.

2 If $\forall x \neq y \in \mathcal{V}$, d(x, y) = 1 then $(\mathcal{V}, \mathcal{E})$ is complete.

Metric

Graph diameter: The diameter of a graph $G = (\mathcal{V}, \mathcal{E})$ is the largest distance between two vertices of the graph:

$$D(G) = \max_{x,y \in \mathcal{V}} d(x,y)$$

Image: A matrix and a matrix

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Example: 1 5 2 3 Radu Balan (UMD) Graphs 1

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$$A^T = A$$

For directed graphs the adjacency matrix may not be symmetric.

The Adjacency Matrix

For undirected graphs the adjacency matrix is always symmetric:

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For directed graphs the adjacency matrix may not be symmetric. For weighted graphs $G = (\mathcal{V}, \mathcal{E}, W)$, the weight matrix W is simply given by

$$W_{i,j} = \left\{ egin{array}{cc} w_{i,j} & \textit{if} & (i,j) \in \mathcal{E} \ 0 & \textit{otherwise} \end{array}
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Degree Matrix d(v) and D

For an undirected graph $G = (\mathcal{V}, \mathcal{E})$, let $d_v = d(v)$ denote the number of edges at vertex $v \in \mathcal{V}$. The number d(v) is called the degree (or valency) of vertex v. The *n*-vector $d = (d_1, \dots, d_n)^T$ can be computed by

$$d = A \cdot 1$$

where A denotes the adjacency matrix, and 1 is the vector of 1's, ones(n,1).

Let *D* denote the diagonal matrix formed from the degree vector *d*: $D_{k,k} = d_k$, $k = 1, 2, \dots, n$. *D* is called the *degree matrix*.

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Vertex Degree Matrix D

For an undirected graph $G = (\mathcal{V}, \mathcal{E})$ of *n* vertices, we denote by *D* the $n \times n$ diagonal matrix of degrees: $D_{i,i} = d(i)$.

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For a graph $G = (\mathcal{V}, \mathcal{E})$ the graph Laplacian is the $n \times n$ symmetric matrix Δ defined by:

$$\Delta = D - A$$





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Example:



Intuition



Assume $x = [x_1, x_2, x_3, x_4, x_5]^T$ is a signal of five components defined over the graph. The *Dirichlet* energy *E*, is defined as

$$E = \sum_{\substack{(i,j) \in \mathcal{E} \ i < j}} (x_i - x_j)^2 =$$

Intuition



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$$E = \sum_{\substack{(i,j) \in \mathcal{E} \\ i < j}} (x_i - x_j)^2 = (x_2 - x_1)^2 + (x_3 - x_2)^2$$

$$+(x_4-x_3)^2+(x_5-x_4)^2+(x_1-x_5)^2.$$

Intuition



By regrouping the terms we obtain:

$$E = \langle \Delta x, x \rangle = x^T \Delta x = x^T (D - A) x$$

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Graph Laplacian

Example



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Graph Laplacian

Example



Normalized Laplacians $\tilde{\Delta}$

Normalized Laplacian: (using pseudo-inverses)

$$\begin{split} \tilde{\Delta} &= D^{\dagger/2} \Delta D^{\dagger/2} = \tilde{I} - D^{\dagger/2} A D^{\dagger/2} \\ \tilde{\Delta}_{i,j} &= \begin{cases} 1 & \text{if} \quad i = j \text{ and } d_i > 0 \text{ (non-isolated vertex)} \\ -\frac{1}{\sqrt{d(i)d(j)}} & \text{if} & (i,j) \in \mathcal{E} \\ 0 & \text{otherwise} \end{cases} \end{split}$$

where \tilde{l} denotes the diagonal matrix: $\tilde{l}_{k,k} = 1$ if d(k) > 0 and $\tilde{l}_{k,k} = 0$ otherwise. Thus \tilde{l} is the identity matrix if and only if the graph has no isolated vertices.

 D^{\dagger} denotes the pseudo-inverse, which is the diagonal matrix with elements:

$$D_{k,k}^{\dagger} = \left\{egin{array}{cc} rac{1}{d(k)} & if & d(k) > 0 \ 0 & if & d(k) = 0 \end{array}
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Graphs

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Normalized Laplacians

Normalized Asymmetric Laplacian:

$$L = D^{\dagger}\Delta = \tilde{I} - D^{\dagger}A$$

$$L_{i,j} = \begin{cases} 1 & \text{if } i = j \text{ and } d_i > 0 \text{ (non - isolated vertex)} \\ -\frac{1}{d(i)} & \text{if } (i,j) \in \mathcal{E} \\ 0 & \text{otherwise} \end{cases}$$

Graphs

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$$\Delta D^{\dagger} = \tilde{I} - AD^{\dagger} = L^{T}$$

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Image: A matrix and a matrix

Normalized Laplacians

Example



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Normalized Laplacians

Example



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Laplacian and Normalized Laplacian for Weighted Graphs

In the case of a weighted graph, $G = (\mathcal{V}, \mathcal{E}, w)$, the weight matrix W replaces the adjacency matrix A.

Laplacian and Normalized Laplacian for Weighted Graphs

In the case of a weighted graph, $G = (\mathcal{V}, \mathcal{E}, w)$, the weight matrix W replaces the adjacency matrix A. The other matrices:

$$D= ext{diag}(W\cdot 1) \hspace{0.2cm}, \hspace{0.2cm} D_{k,k}=\sum_{j\in\mathcal{V}}W_{k,j} \hspace{0.2cm}.$$

 $\Delta=D-W$, $\dim\,\ker(D-W)=$ number connected components $\tilde{\Delta}=D^{\dagger/2}\Delta D^{\dagger/2}$ $L=D^\dagger\Delta$

where $D^{\dagger/2}$ and D^{\dagger} denote the diagonal matrices:

$$(D^{\dagger/2})_{k,k} = \begin{cases} \frac{1}{\sqrt{D_{k,k}}} & \text{if } D_{k,k} > 0\\ 0 & \text{if } D_{k,k} = 0 \end{cases}, \ (D^{\dagger})_{k,k} = \begin{cases} \frac{1}{D_{k,k}} & \text{if } D_{k,k} > 0\\ 0 & \text{if } D_{k,k} = 0 \end{cases}$$

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Laplacian and Normalized Laplacian for Weighted Graphs Dirichlet Energy

For symmetric (i.e., undirected) weighted graphs, the Dirichlet energy is defined as (note edges contribute two terms in the sum)

$$E = \frac{1}{2} \sum_{i,j \in \mathcal{V}} w_{i,j} |x_i - x_j|^2$$

Expanding the square and grouping the terms together, the expression simplifies to

$$\sum_{i\in\mathcal{V}}|x_i|^2\sum_j w_{ij}-\sum_{i,j\in\mathcal{V}}w_{i,j}x_ix_j=\langle Dx,x\rangle-\langle Wx,x\rangle=x^{\mathsf{T}}(D-W)x.$$

Hence:

$$E = \frac{1}{2} \sum_{i,j \in \mathcal{V}} w_{i,j} |x_i - x_j|^2 = x^T \Delta x$$

where $\Delta = D - W$ is the weighted graph Laplacian.

Eigenvalues and Eigenvectors

Recall the eigenvalues of a matrix T are the zeros of the characteristic polynomial:

$$p_T(z) = det(zI - T) = 0.$$

There are exactly *n* eigenvalues (including multiplicities) for a $n \times n$ matrix *T*. The set of eigenvalues is calles its *spectrum*.

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If λ is an eigenvalue of T, then an associated eigenvector is the non-zero *n*-vector *x* such that $Tx = \lambda x$.

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There are exactly *n* eigenvalues (including multiplicities) for a $n \times n$ matrix *T*. The set of eigenvalues is calles its *spectrum*.

If λ is an eigenvalue of T, then an associated eigenvector is the non-zero *n*-vector x such that $Tx = \lambda x$.

Recall: If $T = T^T$ then T is called a *symmetric matrix*. Furthermore:

- Every eigenvalue of T is real.
- There is a set of *n* eigenvectors $\{e_1, \dots, e_n\}$ normalized so that the matrix $U = [e_1 | \dots | e_n]$ is orthogonal $(UU^T = U^T U = I_n)$ and $T = U \wedge U^T$, where \wedge is the diagonal matrix of eigenvalues.

Eigenvalues and Eigenvectors

Recall the eigenvalues of a matrix T are the zeros of the characteristic polynomial:

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Recall: If $T = T^T$ then T is called a *symmetric matrix*. Furthermore:

- Every eigenvalue of T is real.
- There is a set of *n* eigenvectors {e₁, ..., e_n} normalized so that the matrix U = [e₁|...|e_n] is orthogonal (UU^T = U^TU = I_n) and T = UΛU^T, where Λ is the diagonal matrix of eigenvalues.
 Remark. Since det(A₁A₂) = det(A₁)det(A₂) and L = D^{-1/2} Δ̃D^{1/2} it follows that eigs(Δ̃) = eigs(L) = eigs(L^T).

UCINET IV Database: Bernard & Killworth Office Dataset

For the Bernard & Killworth Office dataset (bkoff.dat) dataset we obtained the following results:

The graph is connected. $rank(\Delta) = rank(\tilde{\Delta}) = rank(L) = 39$.



Symmetric Matrices

Recall the following result:

Theorem

Assume T is a real symmetric $n \times n$ matrix. Then:

• All eigenvalues of T are real numbers.

- 2 There are n eigenvectors that can be normalized to form an orthonormal basis for ℝⁿ.
- **③** The largest eigenvalue λ_{max} and the smallest eigenvalue λ_{min} satisfy

$$\lambda_{max} = \max_{x \neq 0} \frac{\langle Tx, x \rangle}{\langle x, x \rangle} \ , \ \lambda_{min} = \min_{x \neq 0} \frac{\langle Tx, x \rangle}{\langle x, x \rangle}$$

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Spectral Analysis

Rayleigh Quotient

For two symmetric matrices T, S we say $T \leq S$ if $\langle Tx, x \rangle \leq \langle Sx, x \rangle$ for all $x \in \mathbb{R}^n$.

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Spectral Analysis

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 $\lambda_{min}I < T < \lambda_{max}I$

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Rayleigh Quotient

For two symmetric matrices T, S we say $T \leq S$ if $\langle Tx, x \rangle \leq \langle Sx, x \rangle$ for all $x \in \mathbb{R}^n$. Consequence 3 can be rewritten:

$$\lambda_{\min} I \leq T \leq \lambda_{\max} I$$

In particular we say T is positive semidefinite $T \ge 0$ if $\langle Tx, x \rangle \ge 0$ for every x.

It follows that T is positive semidefinite if and only if every eigenvalue of T is positive semidefinite (i.e. non-negative).

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