

MULTIDIMENSIONAL SOFIC SHIFTS WITHOUT SEPARATION AND THEIR FACTORS

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ABSTRACT. For $d \geq 2$ we exhibit mixing \mathbb{Z}^d shifts of finite type and sofic shifts with large entropy but poorly separated subsystems (in the sofic examples, the only minimal subsystem is a single point). These examples consequently have very constrained factors; in particular, no non-trivial full shift is a factor. We also provide examples to distinguish certain mixing conditions, and develop the natural class of “block gluing” shifts. In particular, we show that block gluing shifts factor onto all full shifts of strictly smaller entropy.

1. INTRODUCTION

The \mathbb{Z}^d shifts of finite type (SFTs) and sofic shifts ($d \geq 2$) are a drastically more problematic and diverse class of systems than their relatives in only one dimension [24, 26, 28]. The most fundamental theorems in the \mathbb{Z} case do not carry over to the \mathbb{Z}^d setting for $d \geq 2$ and nearly all structural results [4, 5, 9, 12, 14, 16, 17, 18, 24, 27] hold only for certain subclasses of \mathbb{Z}^d SFTs or \mathbb{Z}^d sofic shifts (for a rare exception, see [8]). Nontrivial mixing \mathbb{Z}^d SFTs are qualitatively homogeneous in the case $d = 1$, but quite heterogeneous for $d \geq 2$.

As part of the exploration of this landscape of possibilities, we consider subsystems and factors. Every mixing \mathbb{Z} SFT or sofic shift of positive entropy contains a vast collection of pairwise disjoint subsystems (4.5). For $d \geq 2$, Desai has shown that a general \mathbb{Z}^d sofic shift X of positive entropy still contains many subsystems: it contains subshifts achieving all entropies in the interval $[0, h(X)]$. What we emphasize in this paper is that, in contrast to the \mathbb{Z} case, these subsystems may be very poorly separated. For example, we exhibit mixing \mathbb{Z}^d sofic shifts of arbitrarily large entropy whose only minimal subsystem is a fixed point (so no pair of subsystems can be disjoint). Similarly, for $d \geq 2$ we exhibit \mathbb{Z}^d SFTs of arbitrarily large entropy all of whose minimal subsystems are contained in a subSFT of zero entropy.

Under a factor map, the preimages of disjoint subsystems are disjoint. Consequently systems with poorly separated subsystems cannot factor onto systems with a rich collection of pairwise disjoint subsystems. In particular, our examples cannot factor onto any non-trivial block gluing SFT, and in particular cannot factor onto a non-trivial full shift. This provides a negative answer to the question of Johnson

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and Madden studied in [14, 9, 5, 11] and is another stark contrast to the case $d = 1$ (4.6). The sofic examples cannot factor onto or contain a non-trivial SFT, again in stark contrast to the \mathbb{Z} case (4.6). In some cases, our examples can be arranged to have topologically completely positive entropy, but at the same time to have no factor admitting an invariant measure such that as a measurable system it is of completely positive entropy. The proof (6.1) of the latter property appeals to a \mathbb{Z}^d disjointness theorem of Glasner, Thouvenot and Weiss [10].

The paper is organized as follows. In Section 2, we provide symbolic dynamics background. In Section 3, we define a natural notion of “block gluing” (mixing on blocks given a uniform separation) and show for $d \geq 2$ that any block gluing \mathbb{Z}^d shift X has many factors, including full shifts of any smaller entropy and a set of SFTs with entropies dense in $[0, h(X)]$. This generalizes the corresponding result of Johnson-Madden-Desai for corner gluing shifts [14, 9]. In Section 4, we state our main results on systems with restricted subsystems and factors, and contrast them with the \mathbb{Z} case. The proofs are carried out over Sections 5-10, as described in Section 4. The constructions are different for the cases $d = 2$ and $d \geq 3$, and rely heavily on two papers of Mike Hochman [11, 12].

In the appendices, we establish some properties of block gluing shifts, discuss a range of mixing conditions on \mathbb{Z}^d shifts ($d \geq 2$), and provide some examples. We show the uniform filling property is not implied by corner gluing; block gluing does not imply corner gluing; and a block gluing (or even corner gluing) \mathbb{Z}^d SFT need not be entropy minimal.

2. BACKGROUND

We assume a basic familiarity with symbolic dynamics, nevertheless we recall some notations in this section. Formal definitions of some relevant mixing notions so far introduced to the study of \mathbb{Z}^d shifts by other authors can be found in Appendix A, where we also discuss their mutual relations.

A d -dimensional full shift ($d \in \mathbb{N}$) on some finite alphabet \mathcal{A} consists of the set $\mathcal{A}^{\mathbb{Z}^d}$ together with the \mathbb{Z}^d shift action $\sigma^{\vec{i}} : \mathcal{A}^{\mathbb{Z}^d} \rightarrow \mathcal{A}^{\mathbb{Z}^d}$ ($\vec{i} \in \mathbb{Z}^d$) defined at each coordinate $\vec{j} \in \mathbb{Z}^d$ as $(\sigma^{\vec{i}}(x))_{\vec{j}} := x_{\vec{i}+\vec{j}}$.

Closed shift-invariant subsets of some d -dimensional full shift are called \mathbb{Z}^d (sub)shifts. A subsystem of some \mathbb{Z}^d shift X is a closed shift-invariant subset of X , together with the restriction of the \mathbb{Z}^d shift action to this set.

The language of a \mathbb{Z}^d shift X is the set of finite patterns which appear as subwords of elements of X . It is denoted by $\mathcal{L}(X)$. The subset of patterns in the language of X with a certain (finite) shape $F \subsetneq \mathbb{Z}^d$ we denote by $\mathcal{L}_F(X) := \{x|_F \mid x \in X\} \subsetneq \mathcal{L}(X)$.

For any \mathbb{Z}^d shift X , pattern $P \in \mathcal{L}_F(X)$, and point $x \in X$, the frequency of P in x (if it exists) is given as

$$\lim_{n \rightarrow \infty} \frac{|\{\vec{i} \in F_n \mid x|_{F+\vec{i}} = P\}|}{|F_n|}$$

where $F_n := \{\vec{i} \in \mathbb{Z}^d \mid \|\vec{i}\|_\infty \leq n\}$.

A d -dimensional shift of finite type (\mathbb{Z}^d SFT) is a \mathbb{Z}^d shift $X \subseteq \mathcal{A}^{\mathbb{Z}^d}$ defined using a finite set $\mathcal{P} \subseteq \mathcal{A}^F$ of patterns on a finite non-empty shape $F \subsetneq \mathbb{Z}^d$, such that $X = \{x \in \mathcal{A}^{\mathbb{Z}^d} \mid \forall \vec{i} \in \mathbb{Z}^d : x|_{\vec{i}+F} \in \mathcal{P}\}$. If the diameter of F with respect to the maximum-norm $\|\cdot\|_\infty$ on \mathbb{Z}^d is $m \in \mathbb{N}$, then X has order m .

Recall that the *topological entropy* of a \mathbb{Z}^d shift X is by definition

$$h(X) := \lim_{n \rightarrow \infty} \frac{\log |\mathcal{L}_{F_n}(X)|}{|F_n|}$$

where $F_n := \{\vec{v} \in \mathbb{Z}^d \mid \|\vec{v}\|_\infty \leq n\}$. (Note that by [1] this limit $h(X)$ is the same for any sequence of finite convex shapes $(F_n \subsetneq \mathbb{Z}^d)_{n \in \mathbb{N}}$, such that the inradii of F_n diverge to infinity with n .) The *measure-theoretic entropy* $h_\mu(X)$ of X with respect to a shift-invariant Borel probability $\mu \in \mathcal{M}(X)$ is generalized analogously from the \mathbb{Z} case.

Unless qualified, by a dynamical system we mean in this paper a homeomorphism of a compact metric space. We call a point x of a dynamical system (Ξ, ξ) *transitive* if its ξ -orbit is dense in Ξ . A dynamical system (Ξ, ξ) is called *minimal* if all points are transitive, or in other words if all ξ -orbits are dense in Ξ , meaning that Ξ itself is the only non-empty ξ -invariant closed subset of Ξ . A system (Ξ, ξ) is called *trivial* if it consists of only one point.

A *factor map* between two \mathbb{Z}^d shifts means a surjective map intertwining the shift actions and the image of a subshift X under such a map is referred to as a *factor* of X . A factor map is called *topological* if it is in addition continuous and it is called *measurable* if it is a measure preserving Borel map. In this paper, unless indicated otherwise, factor maps are in the topological category. It is well known that factor maps never increase entropy.

The class of \mathbb{Z}^d *sofic shifts* constitutes the factors of \mathbb{Z}^d SFTs. Obviously it is closed under factor maps, i.e. factors of sofic shifts are again sofic.

By a *measurable dynamical system* we mean in this paper an invertible measure preserving transformation of a Borel probability space. We say a measurable dynamical system has (*measurably*) *completely positive entropy* (c.p.e.) if every non-trivial measurable factor has positive entropy. Likewise we say a topological dynamical system has (*topologically*) *completely positive entropy* (top.c.p.e.) if every non-trivial topological factor has positive entropy.

3. FACTORS OF BLOCK GLUING \mathbb{Z}^d SHIFTS ($d \geq 1$)

In this section we introduce the uniform mixing property “block gluing” for \mathbb{Z}^d shifts, and we prove that block gluing \mathbb{Z}^d shifts factor onto full shifts of smaller entropy (and more generally onto any lower entropy \mathbb{Z}^d SFT containing a safe symbol) as well as onto a family of strongly irreducible \mathbb{Z}^d SFTs realizing a dense set of values in the possible range of entropies. We discuss block gluing shifts in detail in Appendix B.

In the sequel, δ_∞ will represent the maximum-metric on \mathbb{Z}^d , i.e. $\forall \vec{v}, \vec{w} \in \mathbb{Z}^d$ we define $\delta_\infty(\vec{v}, \vec{w}) := \|\vec{v} - \vec{w}\|_\infty = \max_{1 \leq k \leq d} |v_k - w_k|$. δ_∞ can be extended to a non-negative symmetric function on non-empty subsets $V, W \subseteq \mathbb{Z}^d$ in the usual way: $\delta_\infty(V, W) := \min_{\vec{v} \in V, \vec{w} \in W} \delta_\infty(\vec{v}, \vec{w})$.

Let $\vec{v}, \vec{w} \in \mathbb{Z}^d$ such that for all $1 \leq k \leq d$ the components satisfy $\vec{v}_k \leq \vec{w}_k$. We will denote this with $\vec{v} \preceq \vec{w}$. For every pair $\vec{v} \preceq \vec{w}$ the non-empty subset $B := \mathbb{Z}^d \cap \prod_{k=1}^d [\vec{v}_k, \vec{w}_k] = \{\vec{v} \in \mathbb{Z}^d \mid \vec{v} \preceq \vec{v} \preceq \vec{w}\}$ constitutes a *solid block* for which we use the notation $B = [\vec{v}, \vec{w}]$. With $\vec{1}$ we refer to the element in \mathbb{Z}^d in which all components equal 1, i.e. $\vec{1} := (1, 1, \dots, 1) \in \mathbb{Z}^d$ and by \mathbb{N}_0 to the set of non-negative integers.

Definition 3.1. A \mathbb{Z}^d shift X is called *block gluing at gap $g \in \mathbb{N}_0$* if for any two (disjoint) solid blocks $B_1 = [\vec{v}^{(1)}, \vec{w}^{(1)}], B_2 = [\vec{v}^{(2)}, \vec{w}^{(2)}] \subsetneq \mathbb{Z}^d$ with distance $\delta_\infty(B_1, B_2) > g$ any pair of patterns on B_1 and B_2 occurring in X can be put together to form a valid point of X , i.e. $\forall P_1 \in \mathcal{L}_{B_1}(X), P_2 \in \mathcal{L}_{B_2}(X) \exists x \in X$ such that $x|_{B_1} = P_1$ and $x|_{B_2} = P_2$ (or equivalently $\forall y, z \in X \exists x \in X$ such that $x|_{B_1} = y|_{B_1}$ and $x|_{B_2} = z|_{B_2}$).

A \mathbb{Z}^d shift X is called *block gluing* if it is block gluing at gap g for some $g \in \mathbb{N}_0$.

Obviously \mathbb{Z}^d full shifts are block gluing at gap $g = 0$. Note that a \mathbb{Z} shift is block gluing if and only if it satisfies the specification property [2].

In analogy with the matrix-SFT case in [22, 14] we say that a general \mathbb{Z}^d shift X has a *safe symbol* if its alphabet contains a letter a such that for any $x \in X$, any way of changing some letters in x to a s results in a new point still in X .

It is well known that every \mathbb{Z} SFT of entropy at least $\log N$ factors onto the \mathbb{Z} full shift on N symbols. The same holds for $d > 1$ in the case of \mathbb{Z}^d group shifts [4] (note that those have entropy $\log N$ for some $N \in \mathbb{N}$) but not for general \mathbb{Z}^d SFTs [5, 11]. Desai proved for $d \geq 2$ that a corner gluing \mathbb{Z}^d SFT of entropy strictly greater than $\log N$ factors onto a \mathbb{Z}^d full shift on N symbols [9]. (Corner gluing is a condition strictly stronger than block gluing – see Appendix A.) We prove the following generalization of Desai’s theorem.

Theorem 3.2. *Suppose $d \geq 1$ and let X be a block gluing \mathbb{Z}^d shift. Then the following hold.*

- (1) *If $N \in \mathbb{N}$ and $h(X) > \log N$, then X factors topologically onto the \mathbb{Z}^d full shift on N symbols.*
- (2) *X factors topologically onto a family of strongly irreducible \mathbb{Z}^d SFTs with entropies dense in $[0, h(X)]$.*
- (3) *X factors topologically onto any lower entropy \mathbb{Z}^d SFT having a safe symbol.*

Proof. For all three claims we use ideas of Desai’s proof of part (1) of our theorem from [9] (stated there under the stronger assumption of X being a corner gluing \mathbb{Z}^d SFT), with some adaptations. Due to space considerations, we just sketch our argument here, and refer to Desai’s paper [9] for finer details.

We may assume X to be block gluing with gap $g \in \mathbb{N}_0$ and entropy $h(X) > 0$. Fix any $\varepsilon > 0$ such that $\varepsilon < h(X)$. Theorem B.2(4) guarantees the existence of a pattern $P \in \mathcal{L}(X)$ such that the subshift $X^{\mathcal{R}} := \{x \in X \mid P \text{ does not occur in } x\}$ still has entropy $h(X^{\mathcal{R}}) > h(X) - \varepsilon$.

Because $h(X^{\mathcal{R}}) > 0$, there is a pattern $Q \in \mathcal{L}_B(X^{\mathcal{R}})$ on some solid block $B \subsetneq \mathbb{Z}^d$ such that for $\vec{v} \in \mathbb{Z}^d$ with $\|\vec{v}\|_\infty \leq g + \|P\|_\infty$, the pattern Q can not overlap with a copy of Q translated by \vec{v} , i.e. $\exists \vec{j} \in B \cap B - \vec{v} : Q_{\vec{j}} \neq Q_{\vec{j} + \vec{v}}$ (we can use Desai’s original construction to show this). Since X is block gluing, there is a solid block pattern $M \in \mathcal{L}(X)$ with P at each corner and also containing Q (see the left part of Figure 1), w.l.o.g. we may assume P, Q and thus M have the shape of a hypercube.

We define a “surrounded word” to be the central pattern of a configuration on a solid hypercube as pictured in Figure 1. The surrounded words play the role of Desai’s “follower words” in the original proof.

In Figure 1, each G represents any valid pattern of appropriate shape in which P does not occur, i.e. $G \in \mathcal{L}(X^{\mathcal{R}})$ (different G symbols may not represent the same word). The unspecified shaded regions represent valid configurations in X (over which we have no control) filling gaps of thickness g . By the block gluing property,

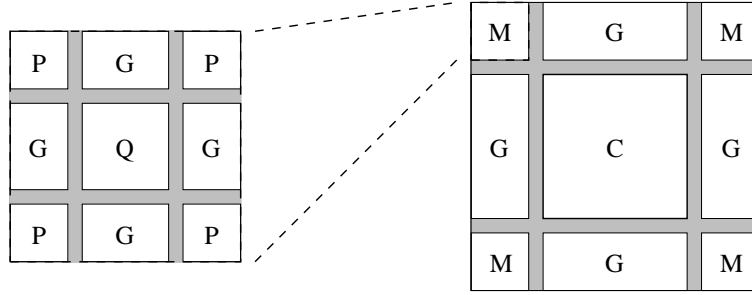


FIGURE 1. The marker pattern M (left), and a surrounded word C (right).

all possible patterns $C \in \mathcal{L}(X^{\mathbb{R}})$ of appropriate size appear as surrounded words in a configuration in X .

To establish claim (1) we choose $\varepsilon := \frac{1}{2}(h(X) - \log N) > 0$ and define a factor map ϕ from X to the full shift on symbols $\mathcal{A} := \{0, \dots, N - 1\}$ which will be defined as suggested by Figure 2. For arbitrarily fixed $n \in \mathbb{N}$ we choose a function $\Phi : \mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathbb{R}}) \rightarrow \mathcal{A}^{[\vec{1}, (n+r)\vec{1}]}$ where $r := 2g + \|M\|_{\infty}$. If in a point in X we see a configuration with a surrounded word $C \in \mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathbb{R}})$ at coordinates $\vec{v} + [\vec{1}, n\vec{1}]$ ($\vec{v} \in \mathbb{Z}^d$), the image configuration on the solid block $\vec{v} - g\vec{1} + [\vec{1}, (n+r)\vec{1}] \subsetneq \mathbb{Z}^d$ indicated as the larger thicklined region in Figure 2 will be $\Phi(C)$. Wherever this rule does not apply (we do not see a complete surrounded-word-pattern in our coding window), we fill in the image coordinates with 0s. The map ϕ induced by this rule is well defined because the marker structure containing Q ensures that two configurations defining surrounded words can overlap only by complete overlap of M corners.

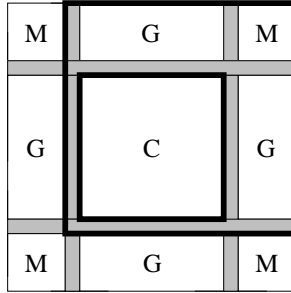


FIGURE 2. A surrounded word C and the (larger) region it uniquely determines in the ϕ -image

As $h(X^{\mathbb{R}}) > h(X) - \varepsilon = \log N + \varepsilon$ we can fix $n \in \mathbb{N}$ large enough such that $|\mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathbb{R}})| > N^{(n+r)^d}$. Thus we may choose Φ surjective. This implies that ϕ is onto: one simply places copies of the pattern M gridding the plane with appropriate distance, and interpolates surrounded words C , to produce a specified image point.

For the second part of the theorem we explicitly construct a family of strongly irreducible \mathbb{Z}^d SFTs. For any fixed $\varepsilon > 0$ we choose $2 \leq m \in \mathbb{N}$ such that $\frac{\log m^d}{m^d} < \frac{\varepsilon}{2}$. For every $j \in \mathbb{N}$ let $\mathcal{A}_{Y_j} := \{0, 1, \dots, (j-1)\}$ and define \mathbb{Z}^d SFTs

$$Y_j := \{y \in \mathcal{A}_{Y_j}^{\mathbb{Z}^d} \mid \forall \vec{v} \in \mathbb{Z}^d : y|_{\vec{v} + [\vec{1}, m\vec{1}]} \text{ contains at most one non-zero symbol}\}.$$

Since there cannot be more than one non-zero symbol in a block of size $[\vec{1}, m\vec{1}]$ we have $|\mathcal{L}_{[\vec{1}, m\vec{1}]}(Y_j)| \leq m^d \cdot j$. Using this and the fact that we can place symbols freely on a grid $(m\mathbb{Z})^d$ while filling in the remaining coordinates with 0s, we have

$$j^{t^d} \leq |\mathcal{L}_{[\vec{1}, tm\vec{1}]}(Y_j)| \leq (m^d \cdot j)^{t^d}$$

for any $t \in \mathbb{N}$. The entropy of Y_j can thus be estimated to be:

$$(EB) \quad \frac{\log j}{m^d} \leq h(Y_j) \leq \frac{\log j}{m^d} + \frac{\log m^d}{m^d} < \frac{\log j}{m^d} + \frac{\varepsilon}{2}.$$

We claim that each Y_j is strongly irreducible with gap m . For this let $P_1 \in \mathcal{L}_{F_1}(Y_j), P_2 \in \mathcal{L}_{F_2}(Y_j)$ be two patterns on finite non-empty sets $F_1, F_2 \subsetneq \mathbb{Z}^d$ such that $\delta_\infty(F_1, F_2) > m$. Clearly the points $y^{(1)}, y^{(2)} \in \mathcal{A}_{Y_j}^{\mathbb{Z}^d}$ defined as $y^{(i)}|_{F_i} := P_i$ and $y^{(i)}|_{\vec{v}} := 0$ for every $\vec{v} \in \mathbb{Z}^d \setminus F_i$ ($i = 1, 2$) are in Y_j (each realizes the pattern P_i in a “minimal” way). Suppose the point $y \in \mathcal{A}_{Y_j}^{\mathbb{Z}^d}$ with $y|_{F_1} := P_1, y|_{F_2} := P_2$ and $y|_{\vec{v}} := 0$ for all $\vec{v} \in \mathbb{Z}^d \setminus (F_1 \cup F_2)$ is not in Y_j . Due to the description of Y_j , there has to be $\vec{v} \in \mathbb{Z}^d$ such that $y|_{\vec{v} + [\vec{1}, m\vec{1}]}$ contains at least two non-zero symbols. But as $y^{(1)}, y^{(2)} \in Y_j$ this can only happen if the block $\vec{v} + [\vec{1}, m\vec{1}]$ contains coordinates of both F_1 and F_2 , which contradicts F_1, F_2 being separated by m . So $y \in Y_j$ and Y_j is strongly irreducible as claimed.

Now choose $j^* \in \mathbb{N}$ such that $h(X) - 2\varepsilon < h(Y_{j^*}) < h(X) - \varepsilon$, which is possible by the choice of m . Again we find $X^{\mathcal{R}} \subsetneq X$ with $h(X^{\mathcal{R}}) > h(X) - \varepsilon$ and we carry out the marker construction as above. This time we fix $n \in \mathbb{N}$ with $n > m$ large enough that

$$\frac{\log |\mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathcal{R}})|}{(n+r)^d} > h(X) - \varepsilon > \frac{\log |\mathcal{L}_{[\vec{1}, (n+r)\vec{1}]}(Y_{j^*})|}{(n+r)^d} > h(X) - 2\varepsilon$$

holds for $r := 2g + \|M\|_\infty$. Hence we have more surrounded words than valid Y_{j^*} -patterns and we can choose a surjection $\Phi : \mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathcal{R}}) \rightarrow \mathcal{L}_{[\vec{1}, (n+r)\vec{1}]}(Y_{j^*})$ to build our factor map ϕ : Whenever ϕ sees a complete marker configuration on $\vec{v} + [(1 - \|M\|_\infty)\vec{1}, (n+r)\vec{1}] \subsetneq \mathbb{Z}^d$ ($\vec{v} \in \mathbb{Z}^d$) with a surrounded word $C \in \mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathcal{R}})$ in the middle, it tries to apply the rule given by Φ to place the pattern $\Phi(C)$ at $\vec{v} + [\vec{1}, (n+r)\vec{1}]$. To ensure that the image point will be valid for Y_{j^*} we have to take care of the local condition defining Y_{j^*} . Therefore ϕ scans a larger neighborhood for surrounded word patterns. If there exists another surrounded word configuration with $C' \in \mathcal{L}_{[\vec{1}, n\vec{1}]}(X^{\mathcal{R}})$ in the middle completely contained within a coding window of size $\vec{v} + [(1 - \|M\|_\infty - n - r - m)\vec{1}, (2n + 2r + m)\vec{1}]$ and the image patterns $\Phi(C), \Phi(C')$ are not compatible with the condition on Y_{j^*} , ϕ just fills in a block of 0s instead. In all other cases ϕ puts 0s too.

As before ϕ is well defined because of the marker structure and it is onto since for every point in Y_{j^*} we can grid \mathbb{Z}^d with markers and fill in appropriate surrounded

words (the rule Φ will always apply, since the image point will satisfy the “non-zero symbols must have distance at least m ”-condition).

Note that for every $j > 1$ there exists a natural factor map $\psi : Y_j \rightarrow Y_{j-1}$ (replacing all occurrences of the symbol $(j-1)$ with 0s), so X factors onto every Y_j with $j \in \{1, 2, \dots, j^*\}$. The choice of m implies that the entropies of the finite family of \mathbb{Z}^d SFTs $(Y_j)_{j \in \{1, 2, \dots, j^*\}}$ is a ε -dense set in $[0, h(X) - \varepsilon]$, i.e. $h(Y_1) = 0$, $h(Y_{j^*}) > h(X) - 2\varepsilon$ and for $2 \leq j \leq j^*$

$$0 \leq h(Y_j) - h(Y_{j-1}) \stackrel{\text{(EB)}}{<} \frac{\log j - \log(j-1)}{m^d} + \frac{\varepsilon}{2} \leq \frac{\log 2}{m^d} + \frac{\varepsilon}{2} < \varepsilon.$$

As $\varepsilon > 0$ was arbitrary we have finished the proof of (2).

For (3) suppose Y is a \mathbb{Z}^d SFT with $h(Y) < h(X)$ and safe symbol 0. The factor map defined for (2) can be used to show that X factors onto Y (we only need that Y is defined by local conditions and replacing symbols with 0s never violates these constraints). As the argument is almost identical, details are left to the reader. \square

We do not know a satisfying characterization of the \mathbb{Z}^d shifts which factor onto full shifts.

4. THE RESTRICTED SUBSYSTEM/QUOTIENT RESULTS

In this section we will state our main results on the existence of sofic/SFT shifts with restricted subsystems and quotients, and indicate the organization of their proofs in subsequent sections. Recall that by a trivial factor we mean the factor which is the trivial dynamical system consisting of one fixed point.

Theorem 4.1. *Given $M > 0$ and $d \geq 2$, there exists a \mathbb{Z}^d sofic shift S with the following properties.*

- (1) $h(S) > M$.
- (2) S (and thus every topological factor of S) contains a fixed point which is its unique minimal subsystem.
- (3) Every topological factor of S (including S) contains a unique \mathbb{Z}^d SFT, which is a fixed point. In particular, S has no non-trivial SFT factor.
- (4) If Y is any non-trivial subshift factor of S , then Y cannot be block gluing. Hence Y also cannot be strongly irreducible, cannot have the uniform filling property (UFP), and cannot be corner gluing. In particular, Y cannot be a full shift.
- (5) In the case $d \geq 3$, S can be chosen such that for every non-trivial topological factor Y of S , there is no invariant Borel probability $\mu \in \mathcal{M}(Y)$ on Y such that (Y, μ) as a measurable system has completely positive entropy.
- (6) S has an equal entropy (sofic) subshift factor of topologically completely positive entropy.
- (7) There is a \mathbb{Z}^d SFT X and a factor map $\pi : X \rightarrow S$ such that $h(X) = h(S)$ and the preimage in X of the unique fixed point in S is a \mathbb{Z}^d SFT $K \subsetneq X$ such that $h(K) = 0$.
- (8) In the case $d = 2$, S can in addition be chosen to be mixing and of topologically completely positive entropy.
- (9) For $d \geq 3$, S can be chosen to be mixing and satisfying all the properties (1)-(7) except (6).

The mixing S we construct to prove Theorem 4.1 in the case $d \geq 3$ has no topological factor which has topologically completely positive entropy.

We have similar results in the case that the subshift must be SFT.

Theorem 4.2. *Given $M > 0$ and $d \geq 2$, there exists a \mathbb{Z}^d SFT X with the following properties.*

- (1) $h(X) > M$.
- (2) X contains a zero entropy SFT K which contains every minimal subsystem of X . In particular, every non-empty subsystem of X has to intersect K .
- (3) If Y is any block gluing shift factor of X , then Y is trivial. In particular, the only full shift factor of X is the trivial shift.
- (4) If Y is any subshift factor of X , and Y_{MIN} is the orbit closure of its minimal subsystems, then $h(Y_{\text{MIN}}) = 0$.
- (5) In the case $d \geq 3$, X can be chosen such that for every non-trivial topological factor Y of X , there is no invariant Borel probability $\mu \in \mathcal{M}(Y)$ on Y such that (Y, μ) as a measurable system has completely positive entropy.
- (6) X has an equal entropy sofic factor with topologically completely positive entropy. In the case $d = 2$, X itself can be chosen mixing with topologically completely positive entropy.

Remark 4.3. Our constructions for some claims in the theorems above require $d = 2$ or $d \geq 3$. However, we know of no obstruction to the existence of examples for all $d \geq 2$ for which all the claims hold. There is some further give and take in the properties we can demand of the examples within the limitations of our current techniques; for this see Remarks 9.1 and 10.1.

Remark 4.4. A \mathbb{Z} sofic shift S contains pairwise disjoint SFTs with entropies dense in $[0, h(S)]$. Desai proved for $d \geq 2$ that a \mathbb{Z}^d sofic shift S contains sofic shifts with entropies dense in $[0, h(S)]$ [8]. Theorem 4.1 shows for $d \geq 2$ that we cannot improve Desai's theorem by replacing sofic with SFT or by requiring disjointness.

Similarly, a \mathbb{Z} SFT X contains pairwise disjoint SFTs with entropies dense in $[0, h(X)]$. Desai (generalizing [25]) proved for $d \geq 2$ that a \mathbb{Z}^d SFT X contains SFTs with entropies dense in $[0, h(X)]$ [8]. Theorem 4.2 shows for $d \geq 2$ we cannot improve Desai's theorem by requiring the SFTs to be disjoint.

Remark 4.5. By way of contrast we recall some facts about the subsystems of a mixing \mathbb{Z} SFT X . By the Jewett Krieger Theorem and one of its proofs, every Borel probability preserving measurable system with entropy strictly less than $h(X)$ has a realization as a uniquely ergodic subshift Y of X [7]. This $Y \subsetneq X$ can be chosen disjoint from any given proper subsystem of X ("there is always room at the Krieger Hotel"). Krieger's Embedding Theorem [15, 19] provides an even more vast array of subsystems. A mixing \mathbb{Z} sofic shift S contains an increasing union of mixing SFTs X_n with $\lim_{n \rightarrow \infty} h(X_n) = h(S)$ [21], so these statements also hold for mixing \mathbb{Z} sofic shifts.

Remark 4.6. The abundance of subsystems in \mathbb{Z} SFTs is related to a number of significant coding constructions and in particular to a considerable control over their factors. For example [3], if S is a \mathbb{Z} sofic shift and Y is a mixing \mathbb{Z} SFT such that $h(S) > h(Y)$, then Y is a factor of S if and only if

$$\forall s \in S, n \in \mathbb{N} \text{ with } \sigma_S^n(s) = s \implies \exists y \in Y \text{ such that } \sigma_Y^n(y) = y .$$

Even for \mathbb{Z} , when Y above is only sofic the statement fails, and the problem of when there is a factor map between mixing \mathbb{Z} sofic shifts of different entropy is only partly solved – see [3, 29]. However, if X is any \mathbb{Z} SFT with $h(X) \geq \log N$, then X factors onto the full \mathbb{Z} shift on N symbols [3, 20]. This led to the following question of Johnson and Madden (see [5] for more motivation).

Question 4.7. [14] If $N \in \mathbb{N}$ is a positive integer and X is a \mathbb{Z}^d SFT with entropy $h(X) \geq \log N$, must there exist a continuous factor map from X onto the \mathbb{Z}^d full shift on N symbols?

Examples proving a negative answer to Question 4.7 for the case $h(X) = \log N$ and $d \geq 2$ were given in [5] and then [11]. Theorem 4.2 proves the negative answer even given a strict entropy inequality.

The proofs of Theorems 4.1 and 4.2 are spread over the next several sections. Sections 5-8 are devoted to the case $d \geq 3$, and are based on a very general subdynamics realization theorem of Hochman [11]. Section 5 establishes properties of a \mathbb{Z} shift which is an ingredient for the construction in Section 6, which proves Theorem 4.1 for the case that $d \geq 3$ and S is not mixing. In Section 7, we elaborate this construction to prove Theorem 4.1 for the case that $d \geq 3$ and S is mixing. As a corollary of this construction, in Section 8 we prove Theorem 4.2 for the case $d \geq 3$.

The constructions for $d = 2$ do not rely on a general theorem but rather on a specific example from another paper of Hochman [12]. In Section 9, we show this example satisfies the conditions of Theorem 4.2 in the case $d = 2$. In Section 10, the case $d = 2$ of Theorem 4.1 follows easily. The powerful constructions of \mathbb{Z}^d sofic and finite type shifts due to Hochman and Meyerovitch [13] and Hochman [11] (for example constructing SFTs of all possible entropies, and \mathbb{Z}^3 sofic shifts with every possible one-dimensional subdynamic) have the feature that the constructed systems lie in the class of subshifts which have an infinite, non-mixing, zero entropy factor. This underlines interest in constructions with stronger mixing properties, and explains our considerable efforts to achieve topologically mixing examples. The property of topologically completely positive entropy is another step away from the class of Hochman-Meyerovitch examples; on the other hand, the lack of any topological factor supporting a c.p.e. measure for our examples with $d \geq 3$ is a pungent residue of the original zero entropy factor.

5. DEFINITION OF A CERTAIN \mathbb{Z} ESS

In this section we define a particular \mathbb{Z} shift $Z \subseteq \{0, 1\}^{\mathbb{Z}}$ which will be used to build our \mathbb{Z}^d examples when $d \geq 3$. For this choose an integer $N \in \mathbb{N}$ so that $\sum_{n=N+1}^{\infty} n^{-1.1} < 1$ and set

$$Z := \{z \in \{0, 1\}^{\mathbb{Z}} \mid \forall n > N : 0^n \text{ appears in every subword of length } \lceil n^{2.1} \rceil \text{ of } z\}.$$

Recall from [11] that a d -dimensional dynamical system is called an *effective symbolic system* (ESS) if it is conjugate to a \mathbb{Z}^d shift X defined over some alphabet \mathcal{A} such that X is given as

$$X := X_{\mathcal{E}} = \{x \in \mathcal{A}^{\mathbb{Z}^d} \mid x \text{ does not contain any pattern from } \mathcal{E}\}$$

where the set of excluded patterns $\mathcal{E} \subseteq \mathcal{L}(\mathcal{A}^{\mathbb{Z}^d})$ is a *recursively enumerable set* (RE-set), i.e. there exists a Turing machine that given input $n \in \mathbb{N}$ outputs a finite d -dimensional pattern $P_n \in \mathcal{L}(\mathcal{A}^{\mathbb{Z}^d})$ over \mathcal{A} such that $\mathcal{E} = \{P_n \mid n \in \mathbb{N}\}$.

Lemma 5.1. *The \mathbb{Z} shift Z defined above has the following properties.*

- (1) Z is topologically mixing.
- (2) For every $n \in \mathbb{N}$ the pattern 0^n appears syndetically in every point in Z . Hence the fixed point 0^∞ is the only minimal subsystem in Z .
- (3) For any pattern $P \in \mathcal{L}(Z)$, P has positive frequency in some point of Z .
- (4) Z is an effective symbolic system.

Proof. To prove (1), take any two valid patterns $P_1, P_2 \in \mathcal{L}_{[1,l]}(Z)$ of length $l \in \mathbb{N}$. For any $m \geq \max\{5, l\}$, define a point $z \in \{0, 1\}^{\mathbb{Z}}$ by $z|_{[1,l]} := P_1$, $z|_{[m+l+1, m+2l]} := P_2$ and $z_i := 0$ for all other $i \in \mathbb{Z}$. We claim that $z \in Z$. To see this, we have to check that for every $n > N$ and every subpattern P of z of length $\lceil n^{2.1} \rceil$, P contains 0^n . We consider two cases: $n \geq m$ and $n < m$.

First suppose $n \geq m$. Using the definition of m this already gives $\lceil n^{2.1} \rceil > 5n \geq 2n + 2l + m$. A subpattern P of z of length $\lceil n^{2.1} \rceil$ is at least $2n + 1$ symbols longer than the concatenation $P_1 0^m P_2$ having length $2l + m$. Therefore P contains more than n consecutive 0s either to the left or to the right of $z|_{[1, m+2l]} = P_1 0^m P_2$.

Now assume $n < m$. Choose any subpattern P of z of length $\lceil n^{2.1} \rceil$. If P contains letters from P_1 as well as from P_2 , then clearly P covers the gap of m 0s between P_1 and P_2 , and so contains 0^n . Otherwise P does not contain letters of both patterns and without loss of generality we may suppose P does not overlap with P_2 . Then P is also a subpattern of the modified point $z' \in \{0, 1\}^{\mathbb{Z}}$, defined by $z'|_{[1,l]} := P_1$ and $z'_i := 0$ for all other $i \in \mathbb{Z}$. Since $P_1 \in \mathcal{L}_{[1,l]}(Z)$ it occurs in some $z'' \in Z$ (with $z''|_{[1,l]} = P_1$). Note that every letter of z' is less than or equal to the corresponding letter of z'' . If $P = z'|_{[j, j+\lceil n^{2.1} \rceil - 1]}$ for some $j \in \mathbb{Z}$, then define $P' := z''|_{[j, j+\lceil n^{2.1} \rceil - 1]}$. As P' is a subpattern of $z'' \in Z$ of length $\lceil n^{2.1} \rceil$, it has to contain 0^n . Finally P occupying the same position in z' as P' in z'' has to be coordinatewise less than or equal to P' and thus P contains 0^n as well.

Therefore any subpattern of z of length $\lceil n^{2.1} \rceil$ contains 0^n and we have shown that $z \in Z$. As P_1, P_2 were chosen arbitrarily Z is (topologically) mixing.

The definition of Z trivially implies (2).

To prove (3), it clearly suffices to consider legal patterns $P_0 \in \mathcal{L}_{[1, l_0]}(Z)$ of big length, e.g. $l_0 > 5^{20}$. The proof of (1) demonstrates that for $m_0 \in \mathbb{N}$ such that the inequality $\lceil n^{2.1} \rceil > 2n + 2l_0 + m_0$ holds for every $n \geq m_0$, the point $z \in \{0, 1\}^{\mathbb{Z}}$ defined by $z|_{[1, l_0]} := z|_{[m_0+l_0+1, m_0+2l_0]} := P_0$ and $z_i := 0$ for all other $i \in \mathbb{Z}$ is in Z . This condition is satisfied by any $m_0 > \sqrt{l_0}$, since for such m_0 and any $n \geq m_0$ we get $\lceil n^{2.1} \rceil > 5n^2 > 2n^2 + 3n > 2n + 2m_0^2 + m_0 > 2n + 2l_0 + m_0$. In particular, the word $P_1 := P_0 0^{m_0} P_0$ is a legal pattern in Z for $m_0 := 2 \lfloor \sqrt{l_0} \rfloor$. Denote by $l_1 \in \mathbb{N}$ the length of P_1 . Then by the same argument, for $m_1 := 2 \lfloor \sqrt{l_1} \rfloor$, the word $P_2 := P_1 0^{m_1} P_1$ is a legal pattern in Z . Continuing in this fashion we inductively define a sequence $(P_n \in \mathcal{L}_{[1, l_n]}(Z))_{n \in \mathbb{N}_0}$, where $P_{n+1} := P_n 0^{2 \lfloor \sqrt{l_n} \rfloor} P_n$, and $l_{n+1} := 2l_n + 2 \lfloor \sqrt{l_n} \rfloor$. Since P_n is a prefix of P_{n+1} for every $n \in \mathbb{N}_0$, there is a point $z \in Z$ so that $z|_{[0, l_n - 1]} = P_n$ for all $n \in \mathbb{N}_0$. We claim that the frequency of P_0 in the point z is positive.

Define the sequence $(r_n := \frac{l_n}{2^n})_{n \in \mathbb{N}_0}$. Then, $r_0 = l_0$ and for $n \in \mathbb{N}_0$ we get $r_{n+1} = \frac{2l_n + 2 \lfloor \sqrt{l_n} \rfloor}{2^{n+1}} = \frac{l_n + \lfloor \sqrt{l_n} \rfloor}{2^n} \leq r_n + 2^{-\frac{n}{2}} \sqrt{r_n}$. For large n we immediately get a loose but sufficient estimate $r_{n+1} < 1.5 r_n$. So we can assume $r_n < 1.62^n$ at least for n big enough. This yields $r_{n+1} \leq r_n + 2^{-\frac{n}{2}} \sqrt{r_n} < r_n + 2^{-\frac{n}{2}} \sqrt{1.62^n} = r_n + 0.9^n$

again for n large. Applying the geometric series, the strictly increasing sequence $(r_n)_{n \in \mathbb{N}_0}$ has a positive limit, call it $r \in \mathbb{R}^+$. Note that by definition, P_0 appears at least 2^n times in $z|_{[0, l_n - 1]} = P_n$ for all $n \in \mathbb{N}_0$. Therefore, the frequency $f_{P_0}(z) \in \mathbb{R}$ of P_0 in z is at least

$$f_{P_0}(z) = \lim_{n \rightarrow \infty} \frac{\# P_0\text{'s in } z|_{[-l_n, l_n]}}{2l_n + 1} \geq \lim_{n \rightarrow \infty} \frac{2^n}{2l_n(1 + \frac{1}{2l_n})} = \lim_{n \rightarrow \infty} \frac{1}{2r_n} = \frac{1}{2r} > 0$$

proving (3).

To get (4), we have to find a RE-set $\mathcal{E} \subseteq \mathcal{L}(\{0, 1\}^{\mathbb{Z}})$ of excluded patterns such that $Z = \{z \in \{0, 1\}^{\mathbb{Z}} \mid z \text{ does not contain any subpattern which is an element of } \mathcal{E}\}$. By simply rewriting the definition of Z , we observe that the following set will do.

$$\mathcal{E} := \bigcup_{n=N+1}^{\infty} \{P \in \{0, 1\}^{[1, \lceil n^{2.1} \rceil]} \mid P \text{ does not contain } 0^n\}$$

It is not difficult to design an algorithm to construct the elements of \mathcal{E} in (length-lexicographic) order. Therefore, \mathcal{E} is a recursively enumerable set and so Z is an effective symbolic system. \square

The \mathbb{Z} shift Z will be used as a main ingredient in the constructions in the following three sections.

6. THE NON-MIXING \mathbb{Z}^d SOFIC EXAMPLES FOR $d \geq 3$

In this section, we prove Theorem 4.1 in the case that $d \geq 3$ and the constructed sofic shift is not required to be mixing.

Proof. Let Z be the ESS obtained in Lemma 5.1. For each integer $1 \leq i \leq d$ define the \mathbb{Z}^d shift T_i as the set of all configurations $t \in \{0, 1\}^{\mathbb{Z}^d}$ satisfying the following:

$$(PC) \quad \begin{aligned} & (t_{n\vec{e}_i})_{n \in \mathbb{Z}} \in Z, \text{ and} \\ & t_{\vec{i}} = t_{\vec{i} + \vec{e}_j} \quad \forall \vec{i} \in \mathbb{Z}^d \text{ and } j \neq i. \end{aligned}$$

So, along the axis $\mathbb{Z}\vec{e}_i$, a point t in T_i is a copy of a point in Z ; in perpendicular directions, t just copies symbols from that axis. Because $d \geq 3$ and Z is a \mathbb{Z} ESS, Hochman's ESS subdynamics realization theorem [11], together with its proof, shows that T_i is a \mathbb{Z}^d sofic shift, such that for each $1 \leq i \leq d$ there is a one-block factor map $\gamma_i : V_i \rightarrow T_i$, where V_i is a zero entropy \mathbb{Z}^d SFT.

Next we map the \mathbb{Z}^d sofic shift $T_1 \times T_2 \times \cdots \times T_d$ onto a sofic shift $S' \subsetneq \{0, 1\}^{\mathbb{Z}^d}$ by the one-block code defined by the rule sending a symbol $(a_1, a_2, \dots, a_d) \in \{0, 1\}^d$ to the product $a_1 a_2 \cdots a_d$. This product is 0 unless every a_i is 1. Note that S' has entropy zero. Finally, we will pick $m \in \mathbb{N}$ and define S by freely splitting the symbol 1 of S' into m symbols $1_1, 1_2, \dots, 1_m$; i.e. a point s is in S if and only if its image under the one-block factor map $\kappa : S \rightarrow S'$ replacing all the 1_j 's ($1 \leq j \leq m$) with 1 is in S' .

By Lemma 5.1, the ESS Z has the following property: every allowed finite configuration occurs with positive frequency in some point of the subshift. This property is inherited by the product $T_1 \times T_2 \times \cdots \times T_d$; and then by the factor S' ; and then by S . As in [13], the entropy of S is $\log m$ multiplied by the maximum frequency of the symbol 1 in a point of S' , which we have noted is positive. We fix $m \in \mathbb{N}$ so that $h(S) > M$, establishing (1).

We have an obvious one-block factor map from the zero entropy \mathbb{Z}^d SFT $X' := V_1 \times V_2 \times \cdots \times V_d$ onto S' , which we denote by $\pi' : X' \rightarrow S'$. We define a \mathbb{Z}^d SFT X by freely splitting each preimage symbol $\tilde{1} \in \pi'^{-1}(1)$ of 1 in X' into m symbols $\tilde{1}_1, \dots, \tilde{1}_m$, and then define a one-block factor map $\pi : X \rightarrow S$ by sending each $\tilde{1}_j$ to 1_j ($1 \leq j \leq m$) and otherwise using the one-block rule for π' . Clearly $h(X) = h(S)$, and the preimage $K := \pi^{-1}(0^{\mathbb{Z}^d})$ of the fixed point $0^{\mathbb{Z}^d} \in S$ has entropy zero. Because $0^{\mathbb{Z}^d}$ and X are SFTs, K is an SFT, and (7) follows.

By construction arbitrarily long blocks of 0s appear syndetically in all points of the one-dimensional ESS Z (5.1(2)). This together with the periodicity constraints (PC) forces the syndetic occurrence of $(d-1)$ -dimensional hyperplanes of 0s of any fixed “thickness” (perpendicular to the \vec{e}_i -axis) in points of T_i . Points from the factor system S' then have a gridlike structure, where the walls of the grid consist of $(d-1)$ -dimensional hyperplanes of 0s separating solid blocks of 1s. In particular, arbitrarily large d dimensional hypercubes of 0s occur in every point, and thus the fixed point $0^{\mathbb{Z}^d}$ is the only minimal subsystem of S . Consequently any factor of S has as its unique minimal subsystem a fixed point, proving (2).

Moreover, for any $n \in \mathbb{N}$ and any finite pattern $P \in \mathcal{L}_F(S)$ ($F \subsetneq \mathbb{Z}^d$ finite) occurring in a point $s \in S$, the pattern P must occur in s inside a pattern $P' := s|_B$ on some larger block $B \supseteq F$ for which all sites within distance n of the boundary are occupied by 0s. If an SFT inside S contains a point s other than $0^{\mathbb{Z}^d}$, then we can find in s a non-zero pattern P' of this form, with n large enough (larger than the order of the SFT), thus a tiling of \mathbb{Z}^d with P' would produce a legal, non-zero periodic point. This is impossible since $0^{\mathbb{Z}^d}$ is the unique minimal subsystem of S . The gridlike structure of 0s is inherited by any \mathbb{Z}^d shift factor of S , proving (3) for subshift factors of S . If there were a factor map ϕ from S onto any compact metrizable system containing a subsystem isomorphic to a non-trivial \mathbb{Z}^d SFT Y , the restriction of ϕ to $\phi^{-1}(Y)$ would be isomorphic to a block code, which we could extend to S to realize an isomorphic copy of Y in some subshift factor of S , which is impossible. This proves (3).

If a non-trivial factor of S were block-gluing, by Theorem B.2 it would contain a subsystem disjoint from the image of the fixed point $0^{\mathbb{Z}^d}$, contradicting (2). This proves (4). The statement (5) follows from Proposition 6.1 as any point in S' contains arbitrarily large hypercubes of 0s and hence condition 6.1(2) is satisfied for $W = \{0^{\mathbb{Z}^d}\}$ being the fixed point of all 0s.

Finally we construct an explicit factor U of S satisfying (6). We define U as the image of S under the one-block map ψ which replaces 1_1 with 0 and changes no other symbol. The (one-block) injective map $\psi \times \kappa : S \rightarrow U \times S'$ has a (one-block) inverse. Thus $h(U) \leq h(S) \leq h(U) + h(S') = h(U)$. This implies $h(S) = h(U)$.

Now suppose there is a factor map $\phi : U \rightarrow Y$ onto a non-trivial subshift Y . Let $\tilde{0}$ and $\tilde{1}$ be distinct symbols in the alphabet of Y such that $\tilde{0}^{\mathbb{Z}^d}$ is the unique fixed point of Y . Let $B := [-n\vec{1}, n\vec{1}]$, where $n \in \mathbb{N}$ is chosen large enough that for all $s \in S$, $s|_B$ determines $(\phi(\psi s))_0$. Choose $s \in S$ such that $(\phi(\psi s))_0 = \tilde{1}$ and let Q be the block $s|_B$. Note that Q contains at least one of the symbols 1_j , $j > 1$.

Let s be a point in S in which Q occurs with positive frequency. Thus there is a subset $\mathcal{C} \subseteq \mathbb{Z}^d$ of positive density in \mathbb{Z}^d such that for all $\vec{v} \in \mathcal{C}$ the restriction of s to $\vec{v} + B$ is Q . Perhaps after reducing the frequency, we may further assume these translates of B by the vectors \vec{v} are pairwise disjoint. Let Q_1 be the pattern which

equals Q , except that every 1_j with $j > 1$ is replaced with 1_1 . If $s|_B = Q_1$, then $(\phi(\psi s))_0 = \tilde{0}$. The block Q_1 can be substituted in s for Q at $\vec{i} + B$ for any subset of vectors \vec{i} from \mathcal{C} to produce a legal point of S . Consequently, in Y all possible patterns of $\tilde{0}$ and $\tilde{1}$ on \mathcal{C} have to appear. This implies that Y has positive entropy, and finishes the proof of the non-mixing and $d \geq 3$ case of Theorem 4.1. \square

Proposition 6.1. *Suppose S is a \mathbb{Z}^d action by homeomorphisms on a compact metric space X such that the following hold:*

- (1) *There exists a topological factor map $\kappa : (X, S) \rightarrow (X', S')$ such that (X', S') is infinite and has zero entropy.*
- (2) *There exists a zero entropy subsystem $(W, S|_W)$ of (X, S) such that for every x' in X' there is a sequence $(\vec{i}_n)_{n \in \mathbb{N}}$ of vectors in \mathbb{Z}^d such that the sequence $S^{\vec{i}_n} \kappa^{-1}(x')$ accumulates on W .*

Then no non-trivial topological \mathbb{Z}^d factor (Y, T) of (X, S) admits an invariant probability $\mu \in \mathcal{M}_T(Y)$ such that (Y, T, μ) as a measurable system has completely positive entropy.

Proof. Let $\phi : (X, S) \rightarrow (Y, T)$ be a continuous factor map, and suppose (Y, T, μ) is c.p.e.. Let $\tilde{\mu} \in \mathcal{M}_S(X)$ be an S -invariant probability such that $\phi \tilde{\mu} = \mu$, and let $\mu' := \kappa \tilde{\mu}$. We have factor maps of \mathbb{Z}^d measure-preserving systems, $\kappa : (X, S, \tilde{\mu}) \rightarrow (X', S', \mu')$ and $\phi : (X, S, \tilde{\mu}) \rightarrow (Y, T, \mu)$. Because one factor is zero entropy and the other is c.p.e., the disjointness theorem of [10] tells us that the map $\kappa \times \phi$ sends $\tilde{\mu}$ onto the product measure $\mu' \times \mu$.

Now consider the decomposition of $\tilde{\mu}$ relative to μ' , which for μ' -almost all x' in X' assigns a measure $\tilde{\mu}_{x'}$ to the fiber $\kappa^{-1}(x')$. The family of measures $(\kappa \times \phi) \tilde{\mu}_{x'}$ on the sets $\{x'\} \times Y$ give a decomposition of μ' relative to the product measure $\mu' \times \mu$. Therefore, on an S -invariant set of full $\tilde{\mu}$ measure, ϕ must send $\tilde{\mu}_{x'}$ to μ . By (2), we may choose a sequence $(\nu_n)_{n \in \mathbb{N}}$ of such measures, $\nu_n := \tilde{\mu}_{x'(\vec{i}_n)}$, which as measures on the compact metric space X converge to a measure ν supported on the subset W . (S -invariance of the measures ν_n is neither available nor needed.) Because $\phi \nu_n = \mu$ for all $n \in \mathbb{N}$ and ϕ is continuous, it follows that $\phi \nu = \mu$. This means that the support of μ is contained in $\phi(W)$. Because the restriction of T to $\phi(W)$ has zero entropy, this contradicts c.p.e., and proves the Proposition. \square

7. THE MIXING \mathbb{Z}^d SOFIC EXAMPLES FOR $d \geq 3$

Elaborating the construction of Section 6, in this section we will construct a family of topologically mixing \mathbb{Z}^d sofic shifts \tilde{S} to establish Theorem 4.1 for $d \geq 3$ in the mixing case.

The construction of \tilde{S} borrows from the *ribbon shift* from [24], the technique of “transmitting information” over long distances in a \mathbb{Z}^d SFT used in [11] and [13], and the results about SFTs mapping onto multidimensional substitution systems from [23]. Begin with the sofic shifts T_i defined in the previous section, i.e. T_i contains the \mathbb{Z} ESS Z as its one-dimensional subdynamics in the \vec{e}_i -direction, and T_i is periodic with respect to \vec{e}_j for all $1 \leq j \neq i \leq d$. Each T_i has an SFT cover V_i . Recall that each V_i is a nearest neighbor shift, i.e. the forbidden patterns in V_i are just pairs of adjacent symbols in any \vec{e}_j -direction ($1 \leq j \leq d$) and that each symbol carries enough information to determine the corresponding symbol in the sofic factor T_i . Denote by γ_i the one-block factor map from V_i to T_i , and denote by \mathcal{A}_i the alphabet of V_i .

For any $1 \leq i \leq d$, define a new SFT W_i as follows. The alphabet for W_i is $\mathcal{A}_i \dot{\cup} \mathcal{A}'_i$, where each symbol of $\mathcal{A}'_i := \{|_a \mid a \in \mathcal{A}_i\}$ is a line in the \vec{e}_i -direction – which from now on, we will assume to be vertical – labeled with a symbol from \mathcal{A}_i . We will call the elements of \mathcal{A}'_i vertical line segments, and the terms “below” and “above” will be used to refer to adjacency in the \vec{e}_i -direction. We will put constraints on the transitions between elements of $\mathcal{A}_i \dot{\cup} \mathcal{A}'_i$ with the effect that points in W_i contain “gently ascending hyperplanes” of elements of \mathcal{A}_i which are separated vertically by one or more vertical line segments from \mathcal{A}'_i .

We begin with the exact definition of the allowed adjacencies. In points of W_i , if a symbol $b \in \mathcal{A}_i$ appears at coordinate $\vec{v} \in \mathbb{Z}^d$, then for any $k \neq i$, a symbol $c \in \mathcal{A}_i$ must appear at $\vec{v} + \vec{e}_k$ or $\vec{v} + \vec{e}_k + \vec{e}_i$, but not in both places. We say that the *rise in direction \vec{e}_k from \vec{v}* is 0 in the former case and 1 in the latter case. Similarly, a symbol $a \in \mathcal{A}_i$ must appear at either $\vec{v} - \vec{e}_k$ or $\vec{v} - \vec{e}_k - \vec{e}_i$, but not in both places. In any of these cases, the pairs of symbols bc and ab must be legal patterns in the \vec{e}_k -direction in V_i . We also impose a homogeneity constraint on the rise: whenever $i \neq j \neq k \neq i$, and a symbol of \mathcal{A}_i appears at coordinate $\vec{v} \in \mathbb{Z}^d$, we require that the rise in direction \vec{e}_k from \vec{v} equal the rise in direction \vec{e}_k from $\vec{v} + \vec{e}_j$ resp. $\vec{v} + \vec{e}_j + \vec{e}_i$. (Note that rises in different directions add up, i.e. if there is a rise of 1 from the \vec{e}_k direction and also from the \vec{e}_j direction then going from coordinate \vec{v} to $\vec{v} + \vec{e}_k + \vec{e}_j$ the \vec{e}_i coordinate of the ascending hyperplane increases by 2.)

We go on to complete the adjacency rules. If $a \in \mathcal{A}_i$ appears at coordinate \vec{v} , then the vertical line segment $|_a \in \mathcal{A}'_i$ must appear at $\vec{v} - \vec{e}_i$, and some symbol $|_b \in \mathcal{A}'_i$ must appear at $\vec{v} + \vec{e}_i$, where the pair ab must be a legal pattern in the \vec{e}_i -direction in V_i . Above a vertical line segment $|_a \in \mathcal{A}'_i$ we can have either another $|_a \in \mathcal{A}'_i$ or the symbol $a \in \mathcal{A}_i$, but nothing else. Hence putting $a \in \mathcal{A}_i$ at coordinate \vec{v} , at position $\vec{v} + \vec{e}_k$ we have to see either $b \in \mathcal{A}_i$ or $|_b \in \mathcal{A}'_i$ with the pattern ab being valid in the \vec{e}_k -direction in V_i for the chosen $k \neq i$. Finally if we see an element $|_a \in \mathcal{A}'_i$ at \vec{v} we can always use a vertical line segment $|_b \in \mathcal{A}'_i$ for which ab is a valid pattern in the \vec{e}_k -direction in V_i to fill coordinate $\vec{v} + \vec{e}_k$. However in addition to this possibility we may put there also a symbol $d \in \mathcal{A}_i$, whenever the coordinate $\vec{v} - \vec{e}_i$ contains an element of \mathcal{A}_i , say $c \in \mathcal{A}_i$ for which cd again has to be a valid pattern in the \vec{e}_k -direction in V_i .

We define a gently ascending hyperplane in a point w in W_i to be a subset H of \mathbb{Z}^d such that $w_{\vec{v}} \in \mathcal{A}_i$ for all $\vec{v} \in H$, and H is a maximal connected set under paths built out of elementary paths (\vec{v}, \vec{j}) such that both \vec{v} and \vec{j} are in H and $\vec{j} - \vec{v} = \pm(\vec{e}_k + \Delta \vec{e}_i)$ for some $k \neq i$ and some $\Delta \in \{0, 1\}$. The homogeneity constraint guarantees that orthogonal projection along \vec{e}_i from a gently ascending hyperplane to $\{\vec{j} \in \mathbb{Z}^d \mid \vec{j}_i = 0\}$ is a bijection. The definition of gently ascending hyperplanes allows them to ascend upwards with a local slope in cardinal directions not exceeding 1.

Note, for every w in W_i and every $\vec{v} \in \mathbb{Z}^d$, the vertical axis $\vec{v} + \mathbb{Z}\vec{e}_i$ intersects every gently ascending hyperplane in w exactly once. The restriction of w to these intersections, read consecutively, gives a legal configuration for Z (which could be empty, finite, half-infinite or bi-infinite). This configuration determines the \mathcal{A}'_i symbols of the remaining sites on the axis, except perhaps in the case that the set $\{n \in \mathbb{Z} \mid \vec{v} + n\vec{e}_i \in H, \text{ for some gently ascending hyperplane } H\}$ is bounded above.

The effect of these adjacency rules is illustrated in Figure 3. The left half is a portion of a potential point from W_i (\vec{e}_i is in the vertical direction); the right half is

the corresponding word in V_i which must be legal in V_i in order for the left half to be legal in W_i . Rows on the right correspond to gently ascending hyperplanes on the left, and the portions between the gently ascending hyperplanes on the left are filled with letters of \mathcal{A}'_i whose values are forced. The reader should note a slight issue which has to be addressed: our construction of the ascending hyperplanes has only been defined for $d \geq 3$, however Figure 3 shows $d = 2$ dimensions. Nevertheless one can consider the picture as a cross-section of a similar picture for $d \geq 3$. The extra issue which appears for $d \geq 3$ but which cannot be expressed in a two-dimensional picture is the homogeneity constraint on the rises as described above. This is largely technical and should not pose a problem to understanding as long as the reader is aware of it.

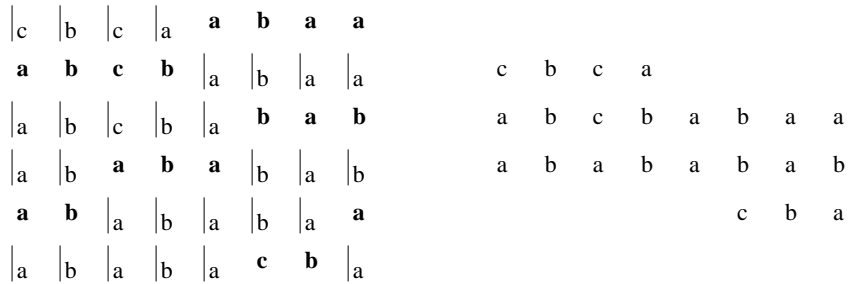


FIGURE 3. A portion of a point from W_i (left) and the portion of a point from V_i (right) which induces it.

The sofic shift $U_i \subseteq \{0, 1\}^{\mathbb{Z}^d}$ is obtained by applying the one-block factor map $\varpi_i : W_i \rightarrow U_i$, where $\varpi_i(a) := \gamma_i(a)$ for $a \in \mathcal{A}_i$, and $\varpi_i(|_a) := 0$ for any $|_a \in \mathcal{A}'_i$ (in order to keep notation simple we will no longer distinguish between a one-block factor map acting on points of a subshift and the corresponding maps acting on the level of valid patterns or even symbols). By a gently ascending hyperplane in U_i we mean the image of a gently ascending hyperplane in W_i . Roughly speaking, flat hyperplanes of 1s in a point of V_i give rise to gently ascending hyperplanes of 1s in U_i , separated vertically by some positive number of 0s.

In a second step we actually have to change W_i and U_i slightly to construct the objects that we will want to work with. The problem with U_i is that there is too much choice in where hyperplanes are allowed to ascend, which will both ruin the grid structure used heavily in Section 6 and cause U_i to already have positive entropy, rendering the proof of part (5) of Theorem 4.1 unusable. We will solve this by superimposing a point from a \mathbb{Z}^{d-1} SFT over each ascending hyperplane which tells the hyperplane when it is possible to ascend one step, and allows ascensions only very rarely. This restrictive SFT is made by using a result of Mozes about substitution dynamical systems in \mathbb{Z}^d for $d \geq 2$, for which we must give some definitions.

Definition 7.1. A \mathbb{Z}^d substitution rule is a map $\tau : \mathcal{A} \rightarrow \mathcal{A}^{[\vec{1}, n\vec{1}]}$ for some integer $n > 1$. For any $m > 1$, τ can also be thought of as a map from $\mathcal{A}^{[\vec{1}, m\vec{1}]}$ to $\mathcal{A}^{[\vec{1}, mn\vec{1}]}$ by identifying $\mathcal{A}^{[\vec{1}, mn\vec{1}]}$ with $(\mathcal{A}^{[\vec{1}, n\vec{1}]})^{[\vec{1}, m\vec{1}]}$. Hence τ can be iterated in a natural way.

Definition 7.2. The \mathbb{Z}^d substitution system associated with a substitution rule τ is the set Ω of all $\omega \in \mathcal{A}^{\mathbb{Z}^d}$ so that every subword of ω is a subword of $\tau^k(a)$ for some $a \in \mathcal{A}$ and some $k \in \mathbb{N}_0$.

In [23], it was shown that for any \mathbb{Z}^d substitution system Ω , there exists an SFT Σ and a one-block factor map ς from Σ onto Ω . If, in addition, the substitution system has a property called unique derivation, then this map can in fact be made into a measure-theoretic isomorphism, though a surjection will suffice for our purposes. Consider the \mathbb{Z}^{d-1} substitution system Ω_1 defined by the following substitution rule $\tau_1 : \{0, 1\} \rightarrow \{0, 1\}^{[\vec{1}, 9\vec{1}]}$ (in the picture, the \vec{e}_1 -direction is horizontal).

$$\begin{array}{ccc}
 & \begin{array}{c} 0 \\ \rightarrow \end{array} & \begin{array}{c} 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \\ 0000000000 \end{array} & & \begin{array}{c} 1 \\ \rightarrow \end{array} & \begin{array}{c} 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \\ 1000000001 \end{array}
 \end{array}$$

The shown substitution rule is for $d - 1 = 2$; for $d - 1 > 2$ it is extended to the remaining directions by concatenating copies of the image-patterns in the picture above. Explicitly: $\tau_1(0) := 0^{[\vec{1}, 9\vec{1}]}$, and $\tau_1(1)$ is defined by $(\tau_1(1))_{\vec{j}} := 0$ if $\vec{j}_1 \in \{2, 3, \dots, 8\}$ and $(\tau_1(1))_{\vec{j}} := 1$ if $\vec{j}_1 \in \{1, 9\}$ (for $\vec{j} \in [\vec{1}, 9\vec{1}]$).

Now by Mozes' theorem, choose a \mathbb{Z}^{d-1} SFT Σ_1 which surjectively factors onto Ω_1 via a one-block factor map ς_1 . Create \mathbb{Z}^{d-1} SFTs Σ_k for $2 \leq k \leq d - 1$ in the same way, i.e. Σ_k factors onto Ω_k via ς_k where Ω_k denotes the substitution system obtained from the rule $\tau_k(0) := 0^{[\vec{1}, 9\vec{1}]}$, and $(\tau_k(1))_{\vec{j}} := 0$ if $\vec{j}_k \in \{2, 3, \dots, 8\}$ and $(\tau_k(1))_{\vec{j}} := 1$ if $\vec{j}_k \in \{1, 9\}$. Next, define the \mathbb{Z}^{d-1} SFT $\Sigma := \Sigma_1 \times \Sigma_2 \times \dots \times \Sigma_{d-1}$ whose alphabet we denote by \mathcal{A}_Σ , as well as the sofic shift $\Lambda \subsetneq \{0, 1\}^{\mathbb{Z}^{d-1}}$ obtained by applying the one-block factor map

$$\eta : \Sigma \rightarrow \Lambda, (s_1, s_2, \dots, s_{d-1}) \mapsto (\max\{(\varsigma_k(s_k))_{\vec{j}} \mid 1 \leq k \leq d - 1\})_{\vec{j} \in \mathbb{Z}^{d-1}}$$

which sends a symbol $(s_1, s_2, \dots, s_{d-1})_{\vec{j}} \in \mathcal{A}_\Sigma$ from the alphabet of Σ to 0 if $(\varsigma_k(s_k))_{\vec{j}} = 0$ for all $1 \leq k \leq d - 1$ and to 1 in any other case. Λ then consists of 0/1-configurations on \mathbb{Z}^{d-1} with a grid structure of sparse \mathbb{Z}^{d-2} -hyperplanes of 1s in each cardinal direction.

Claim 7.3. The number of 1s within a solid block $B := [\vec{1}, n\vec{1}] \subsetneq \mathbb{Z}^{d-1}$ in points of Λ is bounded by $n^{(d-2)+.32}$ for large $n \in \mathbb{N}$. In particular the frequency of 1s in any point of Λ is zero and thus every point in Λ contains arbitrarily large \mathbb{Z}^{d-1} -hypercubes of 0s.

Proof. The sparsely appearing \mathbb{Z}^{d-2} -hyperplanes of 1s along each \vec{e}_k -direction ($1 \leq k \leq d - 1$) in points of Λ stem from the substitution rule τ_k . As $\tau_k(0) = 0^{[\vec{1}, 9\vec{1}]}$ an upper bound to the number of \mathbb{Z}^{d-2} -hyperplanes being perpendicular to \vec{e}_k and carrying symbols 1 in a legal Λ -pattern of shape B is achieved by applying the substitution rule $\lceil \log_9 n \rceil$ times to the symbol 1. In every iteration the number of \mathbb{Z}^{d-2} -hyperplanes of 1s doubles and the total number of 1s is multiplied by a factor $2 \cdot 9^{(d-2)}$. Simply adding the numbers produced in all the $d - 1$ directions,

we slightly overestimate the total number of 1s, getting a bound:

$$\begin{aligned} \# \text{ 1s in any } P \in \mathcal{L}_B(\Lambda) &\leq (d-1) \cdot (2 \cdot 9^{(d-2)})^{\lceil \log_9 n \rceil} \leq \\ &\leq (d-1) \cdot 2 \cdot 2^{\log_9 n} \cdot 9^{(d-2)} \cdot n^{(d-2)} = C \cdot n^{(d-2) + \frac{\log 2}{\log 9}} \end{aligned}$$

with a constant $C := 2 \cdot 9^{(d-2)} \cdot (d-1)$. As $\frac{\log 2}{\log 9} < .3155$ we established the first part of our claim (increasing the exponent to $(d-2) + .32$ we get rid of the constant for large enough values of n).

Moreover this bound shows that the frequency of 1s

$$f_{(1)} \leq \lim_{n \rightarrow \infty} \frac{C \cdot n^{(d-2) + \frac{\log 2}{\log 9}}}{n^{(d-1)}} = C \cdot \lim_{n \rightarrow \infty} n^{\frac{\log 2}{\log 9} - 1} = 0$$

is zero in every point of Λ and thus in particular there can be no bound on the size of hypercubes of 0s seen in each single point in Λ . \square

For each $1 \leq i \leq d$ let $\tilde{\Sigma}_i := \{\tilde{s} \in \mathcal{A}_\Sigma^{\mathbb{Z}^d} \mid \forall n \in \mathbb{Z} : \tilde{s}|_{\{\bar{j} \in \mathbb{Z}^d \mid \bar{j}_i = n\}} \in \Sigma\}$ be the \mathbb{Z}^d SFT obtained as the full \mathbb{Z} extension of Σ along direction \tilde{e}_i . Define the \mathbb{Z}^d SFT \tilde{W}_i constructed just as W_i , with the following changes: the alphabet of \tilde{W}_i is $\tilde{\mathcal{A}}_i \cup \tilde{\mathcal{A}}'_i$, where $\tilde{\mathcal{A}}_i := \mathcal{A}_i \times \mathcal{A}_\Sigma$ and each element of $\tilde{\mathcal{A}}'_i$ is again a vertical line segment this time labeled with a symbol of $\tilde{\mathcal{A}}_i$. Points in \tilde{W}_i arise from elements in the SFT $V_i \times \tilde{\Sigma}_i$ as points in W_i arose from elements in V_i with only one exception: in addition to the restrictions on adjacencies described for W_i above we introduce an extra rule allowing ascensions only at a sparse set of coordinates governed by the superimposed points from $\tilde{\Sigma}_i$. More explicitly, for any $\tilde{v} \in \mathbb{Z}^d$, and for any $1 \leq k \neq i \leq d$, if a symbol $\tilde{a} = (a, s) \in \tilde{\mathcal{A}}_i$ appears at \tilde{v} , then a symbol $\tilde{b} = (b, t) \in \tilde{\mathcal{A}}_i$ is allowed to appear at $\tilde{v} + \tilde{e}_k + \tilde{e}_i$ only when $\eta(t) = 1$; otherwise, \tilde{b} has to appear at coordinate $\tilde{v} + \tilde{e}_k$. Claim 7.3 then shows that on a large scale ascensions of hyperplanes happen very rarely. Finally define the sofic shift \tilde{U}_i (a subsystem of U_i) as the image of \tilde{W}_i under the one-block map given by the rule $\tilde{a} = (a, s) \mapsto \gamma_i(a)$ for all $\tilde{a} = (a, s) \in \tilde{\mathcal{A}}_i$ and $\tilde{a}' \mapsto 0$ for all vertical line segments $\tilde{a}' \in \tilde{\mathcal{A}}'_i$. Denote this one-block map by $\tilde{\omega} : \tilde{W}_i \rightarrow \tilde{U}_i$. A gently ascending hyperplane of a point in \tilde{U}_i is defined with respect to $\tilde{\mathcal{A}}_i$ as a gently ascending hyperplane of a point in U_i was defined with respect to \mathcal{A}_i . If $\tilde{\omega}(w) = u$, then a gently ascending hyperplane in $w \in \tilde{W}_i$ is simply the preimage of a gently ascending hyperplane H in $u \in \tilde{U}_i$, and on a large scale H has rare ascensions.

Next we have to prove some elementary properties of the sofic shifts \tilde{U}_i that will be used later.

Claim 7.4. The topological entropy of \tilde{U}_i ($1 \leq i \leq d$) is zero.

Proof. As all of the shifts \tilde{U}_i are constructed in the same way, it is enough to show that $h(\tilde{U}_1) = 0$. In order to get a sufficient (upper) bound on the entropy of \tilde{U}_1 we do a very coarse estimate on the number of legal patterns in \tilde{U}_1 on some large solid block $B_n := [\bar{\mathbb{1}}, n\bar{\mathbb{1}}] \subsetneq \mathbb{Z}^d$ ($n \in \mathbb{N}$). First note that a gently ascending hyperplane of 1s which intersects B_n must intersect the set $\{\bar{\mathbb{1}} + j\tilde{e}_1 \mid -d(n-1) < j < n\}$ at a single point. For $u \in \tilde{U}_1$, define $J_n(u)$ as the set of $j \in \mathbb{Z}$ such that $-d(n-1) < j < n$ and $\bar{\mathbb{1}} + j\tilde{e}_1$ is part of some gently ascending hyperplane of 1s in u . Trivially, the number $|\{J_n(u) \mid u \in \tilde{U}_1\}|$ is bounded above by $2^{n(d+1)}$. It is not hard to see that $u|_{B_n}$ is

uniquely determined – use the fact that \tilde{U}_1 has alphabet $\{0, 1\}$ – by the set $J_n(u)$, together with the sets of coordinates within each hyperplane at which ascensions actually occur inside B_n . The latter have to be subsets of the coordinates containing the symbol 1 in the corresponding superimposed Λ -point intersected with B_n and their cardinality is bounded by $n^{(d-2)+.32}$ (7.3) for large n . Hence we get an upper bound on the number of valid \tilde{U}_1 -patterns on $B_n = [\vec{1}, n\vec{1}]$ with n large enough

$$|\mathcal{L}_{B_n}(\tilde{U}_1)| \leq 2^{n(d+1)} \cdot \left(\sum_{m=0}^{\lfloor n^{(d-1.68)} \rfloor} \binom{n^{(d-1)}}{m} \right)^{n(d+1)}.$$

For n large enough $n^{(d-1.68)} \ll \frac{1}{2} n^{(d-1)}$, so the binomial coefficients grow with m and the sum can be bounded from above by taking the trivial approximation

$$\begin{aligned} |\mathcal{L}_{B_n}(\tilde{U}_1)| &\leq 2^{n(d+1)} \cdot (n^{(d-1.68)} + 1)^{n(d+1)} \cdot \binom{n^{(d-1)}}{\lfloor n^{(d-1.68)} \rfloor}^{n(d+1)} \\ &\leq (2 \cdot n^{(d-1.68)} + 2)^{n(d+1)} \cdot (n^{(d-1)})^{n^{(d-1.68)} \cdot n(d+1)}. \end{aligned}$$

Therefore the entropy of \tilde{U}_1 turns out to be zero as claimed above by the following estimate:

$$\begin{aligned} 0 \leq h(\tilde{U}_1) &= \lim_{n \rightarrow \infty} \frac{\log |\mathcal{L}_{B_n}(\tilde{U}_1)|}{n^d} \\ &\leq \lim_{n \rightarrow \infty} \frac{\log((2 \cdot n^{(d-1.68)} + 2)^{n(d+1)} \cdot n^{(d^2-1) \cdot n^{(d-0.68)}})}{n^d} = 0. \end{aligned}$$

□

Claim 7.5. Every \mathbb{Z}^d shift \tilde{U}_i ($1 \leq i \leq d$) as defined above is topologically mixing.

Proof. We can show that \tilde{U}_i is topologically mixing in much the same way as was done for the *ribbon shift* in [24]. Consider any two congruent solid blocks $B_1, B_2 \subsetneq \mathbb{Z}^d$ spaced at a large distance $D_{B_1, B_2} \in \mathbb{N}$ apart and legal \tilde{U}_i -patterns $P_1 \in \mathcal{L}_{B_1}(\tilde{U}_i)$ on B_1 and $P_2 \in \mathcal{L}_{B_2}(\tilde{U}_i)$ on B_2 . (The definition of topological mixing uses arbitrary finite subsets, however it is easy to see that checking it for congruent solid blocks suffices.) Without loss of generality we may assume that in our patterns P_j ($j = 1, 2$) there are no symbols 1 along the top or bottom of the solid blocks B_j . If this were the case, just use the fact that P_j is a legal word in \tilde{U}_i , say $P_j = \tilde{u}^{(j)}|_{B_j}$ for some $\tilde{u}^{(j)} \in \tilde{U}_i$, to extend it upwards and downwards to a pattern $P'_j := \tilde{u}^{(j)}|_{B'_j}$ on a solid block $B'_j \supseteq B_j$ in which every hyperplane intersecting the original P_j does not intersect the top or bottom of B'_j . Then, replace all symbols in P'_j above the topmost hyperplane which intersects P_j and below the bottommost hyperplane which intersects P_j with 0s. This results in a pattern P''_j (still on B'_j) which is legal in \tilde{U}_i (since thickening layers of 0s is always possible), which contains P_j as a subpattern, and which has no hyperplanes of 1s intersecting its top and bottom. As the slope of our gently ascending hyperplanes is at most 1 in cardinal directions, the required vertical extension of the pattern P_j to a pattern P''_j is linear in the size of B_j and by adding some extra hyperplanes of 0s on top of the smaller one of the two patterns P''_j if necessary, we reestablish our assumption that both patterns have a congruent shape. For notational simplicity

from now on we will drop the 's, so talking about P_j on B_j we actually refer to the enlarged patterns constructed above.

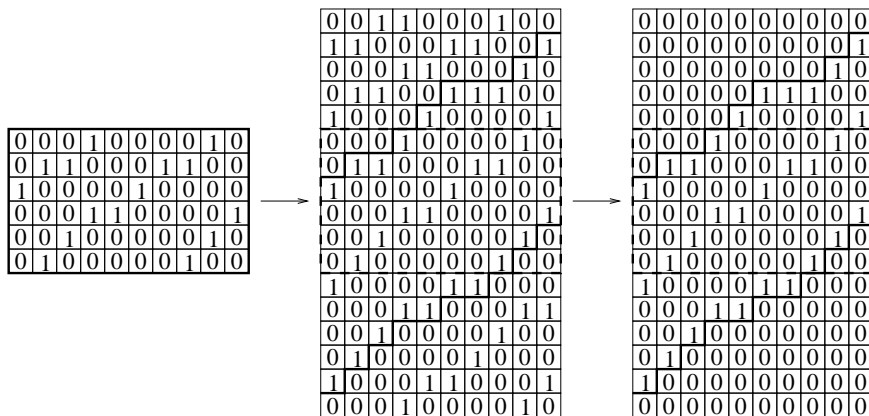


FIGURE 4. Extending a pattern P_j on the solid block B_j (on the left) via a pattern P'_j on B'_j (in the middle) to a pattern P''_j also on B'_j (on the right).

After this preprocessing there are two cases; either B_1 and B_2 are separated by a large distance in the \vec{e}_i -direction, or they have large separation in an \vec{e}_k -direction for some $1 \leq k \neq i \leq d$. If they have large separation in the \vec{e}_i -direction, then the fact that the patterns P_1 and P_2 can be put together to form a valid point in \tilde{U}_i comes from mixing of Z (condition 1 of Lemma 5.1). To create a \tilde{U}_i -configuration on \mathbb{Z}^d containing P_1 on B_1 and P_2 on B_2 , simply first extend all gently ascending (finite) regions of 1s inside P_1 and P_2 to ascending hyperplanes of 1s with no additional ascensions (apart from the ones forced by the homogeneity condition) outside B_1 and B_2 . Fill the gaps between those hyperplanes with 0s and add a layer of 0s just above and below the extended patterns. Then use mixing of Z to interpolate between those infinite configurations, putting in non-ascending hyperplanes of 1s and 0s surrounded vertically by layers of 0s according to a point in Z that contains subwords realizing both patterns P_1 and P_2 . As long as the \vec{e}_i -separation is large enough, this can always be done.

Suppose now that B_1 and B_2 have large separation in some direction \vec{e}_k with $k \neq i$. In this case we exploit the possibility to let hyperplanes of 1s ascend in order to interpolate between P_1 and P_2 . Roughly speaking, we again extend the gently ascending (finite) regions of 1s inside P_1 and P_2 to ascending hyperplanes, but outside B_1 and B_2 we allow only ascensions in the \vec{e}_k -direction, apart from the ascensions forced by the homogeneity constraints. Then, interpolating between these two sets of hyperplanes in \tilde{U}_i may be done as in the first case, since due to the ascending and the sufficiently large initial \vec{e}_k -separation of B_1 and B_2 , the two stacks of hyperplanes observed at the same \vec{e}_k -coordinate will already have large separation in the \vec{e}_i -direction. We have then shown that as long as B_1 and B_2 have large separation in some direction, P_1 and P_2 can be put together to form a valid point of \tilde{U}_i , and so \tilde{U}_i is mixing. \square

We now proceed just as in Section 6. Take the product shift $\tilde{U}_1 \times \tilde{U}_2 \times \dots \times \tilde{U}_d$, and factor it onto the \mathbb{Z}^d shift $\tilde{S}' \subsetneq \{0, 1\}^{\mathbb{Z}^d}$ by the one-block factor map

$$\tilde{\vartheta} : \tilde{U}_1 \times \tilde{U}_2 \times \dots \times \tilde{U}_d \rightarrow \tilde{S}', (\tilde{u}^{(1)}, \tilde{u}^{(2)}, \dots, \tilde{u}^{(d)}) \mapsto (\tilde{u}_i^{(1)} \cdot \tilde{u}_i^{(2)} \cdot \dots \cdot \tilde{u}_i^{(d)})_{i \in \mathbb{Z}^d}.$$

As each \tilde{U}_i was a topologically mixing sofic shift, \tilde{S}' is mixing sofic as well. We also note that $h(\tilde{S}') = 0$; because \tilde{S}' is a factor of $\tilde{U}_1 \times \tilde{U}_2 \times \dots \times \tilde{U}_d$, this follows from our Claim 7.4. Just as every point of T_i contained syndetically appearing thick hyperplanes of 0s, every point of \tilde{U}_i contains a set of arbitrarily thick gently ascending hyperplanes of 0s.

The next result shows that hyperplane ascensions in points of \tilde{U}_i are sufficiently slow that we may recover the key feature of the gridlike structure of 0s from U_i : finite patterns must occur within block patterns which are labeled with 0s in an arbitrarily large neighborhood of the boundary.

Claim 7.6. Suppose $\tilde{s} \in \tilde{S}'$ and $l, m \in \mathbb{N}$. Then there exists a solid block $B := [\vec{v}, \vec{w}] \subsetneq \mathbb{Z}^d$, containing $[-(l+m)\vec{1}, (l+m)\vec{1}]$, such that $\tilde{s}_{\vec{j}} = 0$ whenever $\vec{j} \in [\vec{v}, \vec{w}] \setminus [\vec{v} + l\vec{1}, \vec{w} - l\vec{1}]$.

Proof. We start with a technical observation about the variation in thickness of layers of 0s in points of \tilde{U}_i ($1 \leq i \leq d$). Suppose somewhere in a point $\tilde{u} \in \tilde{U}_i$ there appears a block of $n \in \mathbb{N}$ consecutive 0s along the \vec{e}_i -direction. For convenience, assume this happens on the interval $[0, n-1]\vec{e}_i := \{j\vec{e}_i \in \mathbb{Z}^d \mid 0 \leq j < n\} \subsetneq \mathbb{Z}^d$. We claim that for large n , $\tilde{u}_{\vec{j}} = 0$ as well for any $\vec{j} \in \mathbb{Z}^d$ with $\frac{n}{4} \leq \vec{j}_i < \frac{3n}{4}$ and $\|\vec{j}\|_\infty \leq n^{3.1} + m + l$. In other words, within a distance of $n^{3.1} + m + l$ in all directions, hyperplanes do not ascend by more than $\frac{n}{4}$. This is not hard to see: consider the ascending hyperplanes of 1s which bound our pattern of 0s containing the interval $[0, n-1]\vec{e}_i$ from above and below – if those do not exist, i.e. the region of 0s is unbounded from above/from below, the claim is trivially true. Along each direction \vec{e}_k ($k \neq i$), these bordering hyperplanes may ascend only at coordinates where there is a \mathbb{Z}^{d-2} -hyperplane of 1s perpendicular to \vec{e}_k in the superimposed Λ -point. By earlier computations (see Claim 7.3) the number of those over a distance of $n^{3.1} + m + l$ is bounded by $(n^{3.1} + m + l)^{\frac{\log 2}{\log 9}}$. As the hyperplanes in \tilde{U}_i “climb up” just one coordinate in each ascension, after summing over all \vec{e}_k -directions, we can (over)estimate the total gain in height in direction \vec{e}_i by $(d-1) \cdot (n^{3.1} + m + l)^{.32} < (d-1) \cdot n^{.993} < \frac{n}{4}$ for n large enough.

Now consider any point $\tilde{s} \in \tilde{S}'$ and some $\tilde{\vartheta}$ -preimage $(\tilde{u}^{(1)}, \tilde{u}^{(2)}, \dots, \tilde{u}^{(d)}) \in \tilde{\vartheta}^{-1}(\tilde{s})$. The construction of \tilde{U}_i ensures that the point $\tilde{u}^{(i)}$ contains the word 0^n somewhere in the interval $[m+l, n^{3.1}+m+l-1]\vec{e}_i$ and also somewhere in the interval $[-(n^{3.1}+m+l-1), -(m+l)]\vec{e}_i$. (Here $n^{3.1}$ appears rather than the distance $n^{2.1}$ from the definition of Z to address the fact that successive hyperplanes of 1s are separated vertically with gap at most n .) We define the i th coordinates of the vectors \vec{v}, \vec{w} as follows. Using the technical result of the previous paragraph, we pick some interval $I_i = [\vec{w}_i - \frac{n}{2}, \vec{w}_i] \subseteq [m+l, n^{3.1}+m+l-1] \cap \mathbb{Z}$, such that $\tilde{u}_{\vec{j}}^{(i)} = 0$ for any $\vec{j} \in \mathbb{Z}^d$ with $\vec{j}_i \in I_i$ and $\|\vec{j}\|_\infty \leq n^{3.1} + m + l$. Similarly we pick an interval $[\vec{v}_i, \vec{v}_i + \frac{n}{2}]$ on the negative \vec{e}_i -axis. After applying $\tilde{\vartheta}$, we see that the restriction of \tilde{s} to $[\vec{v}, \vec{w}]$ has 0s at sites within distance $\frac{n}{2}$ of the boundary. We may require $\frac{n}{2} > l$, and this proves the claim. \square

If $z \in Z$, then for each \tilde{U}_i there is a point $\tilde{u}^{(i)} \in \tilde{U}_i$ such that $\tilde{u}_j^{(i)} := z_n$ when $\tilde{j}_i = 2n$ and $\tilde{u}_j^{(i)} := 0$ otherwise. Since the symbol 1 occurs with positive frequency in some points of Z , it follows that 1s appear with positive frequency in some points of \tilde{U}_i and therefore in some point of \tilde{S}' . Hence, just as in Section 6, we can split the symbol 1 into many distinct copies to create the \mathbb{Z}^d sofic shift \tilde{S} of arbitrarily large entropy $h(\tilde{S}) > M$. Denote by $\tilde{\kappa} : \tilde{S} \rightarrow \tilde{S}'$ the 1-block factor map collapsing the different copies of the symbol 1 in \tilde{S} back onto the original. Since \tilde{S}' was topologically mixing, it is clear that \tilde{S} is as well.

The rest of the proof of Theorem 4.1 for $d \geq 3$ now proceeds just as in the non-mixing case, with one exception. We do not claim that our mixing example \tilde{S}' has a topological factor of completely positive entropy. (In fact, it has no such factor, but we will not give a proof.) Presumably this feature of the mixing example is only a limitation of our technique.

8. THE \mathbb{Z}^d SFT EXAMPLES FOR $d \geq 3$

In this section we prove Theorem 4.2 in the case $d \geq 3$. We choose the required SFT X to be the SFT cover of the sofic shift S described in Theorem 4.1(7). We now verify the claims of Theorem 4.2.

Proof. (1) This is instantly true as X is a cover of S and $h(S) > M$.

(2) Let K be the SFT from part (7) of Theorem 4.1. If there is a minimal subsystem W of X which is not contained in K , then it is disjoint from K , and we have $\pi(W)$ a minimal subsystem of S other than its fixed point (4.1(7)), contradicting part (2) of Theorem 4.1.

(3,4) Suppose Y is non-trivial block gluing and $\phi : X \rightarrow Y$ is a factor map. Using Theorem B.2, we have $h(Y) > 0$ (by non-triviality) and then Y contains a subsystem W disjoint from $\phi(K)$. Therefore $\phi^{-1}(W)$ is a subsystem of X disjoint from K , which is impossible. As every minimal subsystem in Y is contained in $\phi(K)$ and $h(\phi(K)) = 0$, we have $h(Y_{\text{MIN}}) = 0$.

(5) Because arbitrarily large hypercubes of 0s occur in every point of S , and K is the preimage of the fixed point $0^{\mathbb{Z}^d}$, it follows that arbitrarily large hypercubes of K -configurations occur in every point of X . Because $h(K) = 0$, this claim (like part (4.1(5))) now follows from Proposition 6.1.

(6) This follows because S is an equal entropy factor of X . (4.1(6)) \square

9. A TOPOLOGICALLY MIXING, TOP.C.P.E. \mathbb{Z}^2 SFT EXAMPLE

Turning our attention to two-dimensional shifts, we present a family of \mathbb{Z}^2 SFTs \tilde{X}_k which prove the case $d = 2$ of Theorem 4.2 (with $X = \tilde{X}_k$ of sufficiently large entropy).

These SFTs are in fact an example of Hochman from [12]. We will begin by giving a brief description of the construction of those \mathbb{Z}^2 SFTs and discuss their properties pertinent to our argument. For a full examination of these systems, see [12].

The construction: For any positive integer $k \in \mathbb{N}$, first a \mathbb{Z}^2 SFT X_k is defined with alphabet \mathcal{A}_k of $32 + k$ symbols, which are most conveniently thought of as square tiles of unit length. Thirty-two of the symbols of \mathcal{A}_k are defined by

assigning a tile one of four colors and one of eight types of arrows. For reasons which will become clear soon, the colors should be thought of as representing the four directions NW (northwest), NE (northeast), SE (southwest), and SE (southeast). The eight types of arrows are four straight arrows in the four cardinal directions (up, down, right, left), and four arrows turning ninety degrees clockwise. The remaining k tiles are called “blanks” and are labeled with an associated integer between 1 and k inclusive. The symbols of \mathcal{A}_k appear in Figure 5.

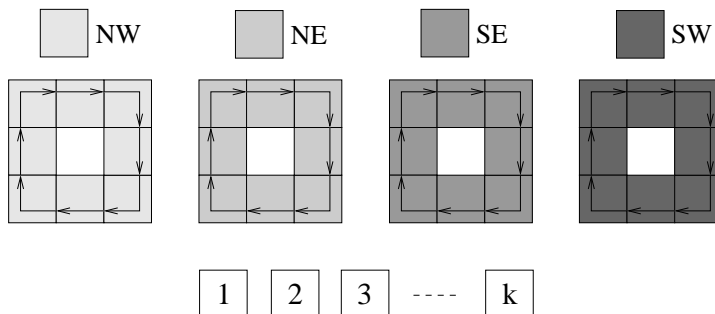


FIGURE 5. The $32 + k$ different symbols of \mathcal{A}_k

The subshift X_k is an SFT of order 2, meaning that there is a set of legal 2×2 patterns $\mathcal{L}_{\{0,1\}^2}(X_k) \subseteq \mathcal{A}_k^{\{0,1\}^2}$, and X_k is precisely the set of elements of $\mathcal{A}_k^{\mathbb{Z}^2}$ in which every 2×2 subpattern is one of the elements in $\mathcal{L}_{\{0,1\}^2}(X_k)$. For the purposes of legal 2×2 patterns, all blanks are considered indistinguishable, and so from now on we refer only to blank tiles in general, with the understanding that the k blank tiles are completely interchangeable in elements of X_k .

We inductively define valid patterns $P_n^\theta \in \mathcal{L}_{[\vec{1}, (2^{n+1} + 2^{n-1} - 2)\vec{1}]}(X_k)$ for any $n \in \mathbb{N}$ and θ one of the four colors NW , NE , SE , and SW as follows (let us denote by $\mathcal{T} := \{NW, NE, SE, SW\}$ the set of colors): for any $\theta \in \mathcal{T}$, P_1^θ is the 3×3 square with a blank tile in the center, surrounded by a clockwise circuit of arrows colored by θ ; for any $n \in \mathbb{N}$ and color $\theta \in \mathcal{T}$, P_{n+1}^θ is constructed by concatenating the four smaller patterns P_n^{NW} , P_n^{NE} , P_n^{SE} , and P_n^{SW} into a square configuration of shape $[\vec{1}, (2(2^{n+1} + 2^{n-1} - 2))\vec{1}]$ with each of the four patterns $P_n^{\theta'}$ in the corner corresponding to $\theta' \in \mathcal{T}$, and then surrounding this square with a clockwise circuit of arrows, colored by θ . For $n \in \mathbb{N}$, $\theta \in \mathcal{T}$, we call P_n^θ the level- n square with θ -colored boundary. As an example, the configuration appearing in Figure 6 is actually P_3^{NW} . The set $\mathcal{L}_{\{0,1\}^2}(X_k)$ of allowed 2×2 patterns in X_k is defined to be the set of all 2×2 configurations which appear in P_3^θ for some $\theta \in \mathcal{T}$. It is verified in [12] that $P_n^\theta \in \mathcal{L}(X_k)$ for all $n \in \mathbb{N}$, $\theta \in \mathcal{T}$, and in fact we will see that in some sense most points in X_k are built up from the patterns P_n^θ . (We are still not concerning ourselves with the labeling of blank tiles; so each P_n^θ actually corresponds to many configurations in $\mathcal{L}(X)$ depending on how the blanks are labeled.)

For any $\omega = (\omega_1, \omega_2, \dots) \in \mathcal{T}^{\mathbb{N}}$, we define $x_\omega \in X_k$ as follows: begin by defining x_ω on a 3×3 square centered at the origin, call this square $B_1 := [-\vec{1}, \vec{1}]$, so that $x_\omega|_{B_1} := P_1^{\omega_1}$. For any $n \geq 1$, assume that x_ω has already been defined on a

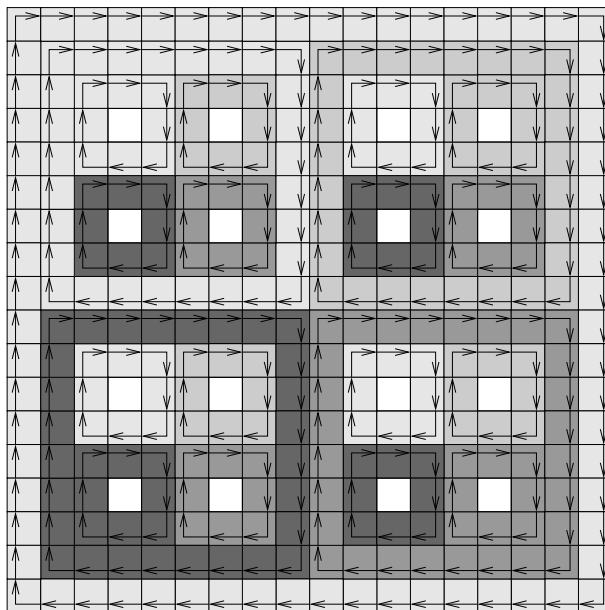
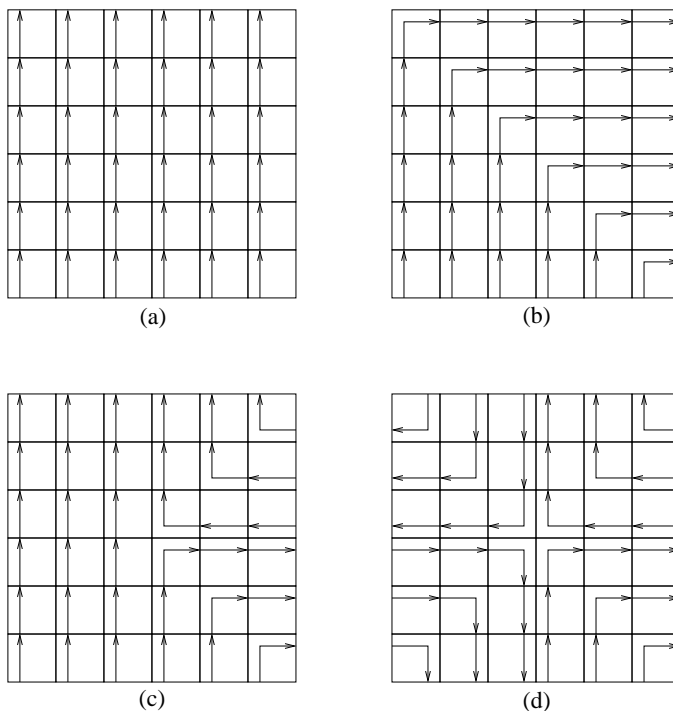


FIGURE 6. Part of a point of X_k

square $B_n \subsetneq \mathbb{Z}^2$ which is a translate of $[\vec{\mathbb{I}}, (2^{n+1} + 2^{n-1} - 2)\vec{\mathbb{I}}]$ so that $x_\omega|_{B_n} = P_n^{\omega_n}$. Since $P_n^{\omega_n}$ appears as a subpattern of $P_{n+1}^{\omega_{n+1}}$ exactly once, there is a unique way to extend the definition of x_ω to a square B_{n+1} which is some translate of $[\vec{\mathbb{I}}, (2^{(n+1)+1} + 2^{(n+1)-1} - 2)\vec{\mathbb{I}}]$ so that $x_\omega|_{B_{n+1}} := P_{n+1}^{\omega_{n+1}}$. In addition, as $P_{n+1}^{\omega_{n+1}}$ has a closed path of arrows on the boundary, it must be the case that $B_n \subseteq B_{n+1} \setminus \partial B_{n+1}$. It is then easily checked by induction that $[-n\vec{\mathbb{I}}, n\vec{\mathbb{I}}] \subseteq B_n$ for all $n \in \mathbb{N}$, and so $\bigcup_{n=1}^{\infty} B_n = \mathbb{Z}^2$, thus x_ω is eventually defined on all of \mathbb{Z}^2 in this way.

In [12], it is proven that every element $x \in X_k$ is of one of two types. If $x = \sigma^{\vec{v}}(x_\omega)$ for some $\vec{v} \in \mathbb{Z}^2$ and $\omega \in \mathcal{T}^{\mathbb{N}}$, then x is transitive. Otherwise x – being a boundary point of the orbit closure of some x_ω – is what Hochman calls an “exceptional point.” Essentially there are four types of exceptional points, shown in Figure 7 without their coloring. Every exceptional point is of one of these four types, possibly rotated by a multiple of ninety degrees, and having the arrows colored in some way allowed in X_k .

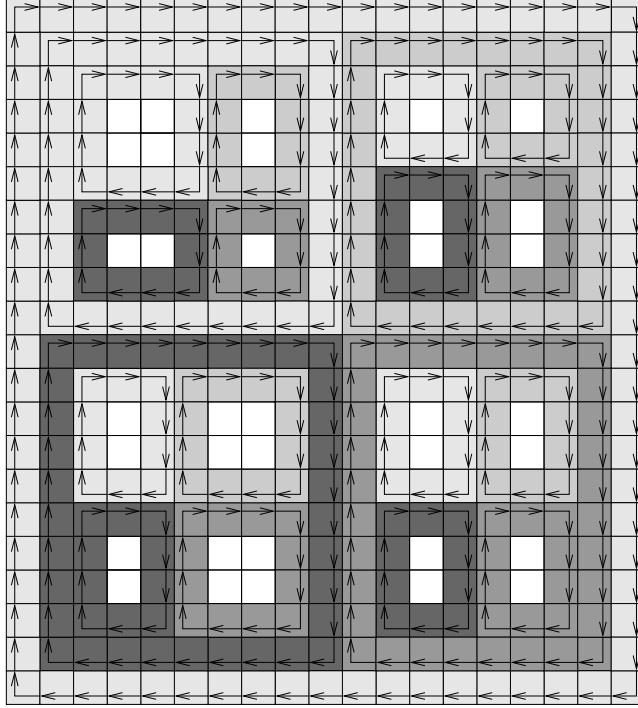
For any subpattern of an exceptional point with shape $B := [\vec{\mathbb{I}}, n\vec{\mathbb{I}}]$, knowing the arrows on the boundary of B suffices to reconstruct the arrows of all interior tiles (there are no blank tiles in an exceptional point and all arrows have to line up). If one knows in addition the colors of the boundary tiles, one can recover the colors in the interior. This is because, due to the finite type restrictions on X_k , any pair of arrows which join up head-to-tail must have the same color. This means that the boundary of any subpattern of an exceptional point forces the configuration on its interior. So if we denote by $E \subsetneq X_k$ the subshift of X_k consisting of all exceptional points, $|\mathcal{L}_{[\vec{\mathbb{I}}, n\vec{\mathbb{I}}]}(E)| \leq |\mathcal{A}_k|^{4n-4} < (32 + k)^{4n}$, and so $h(E) = 0$. (We may use the

FIGURE 7. Exceptional points of X_k (uncolored)

notation E independent of k since elements of E contain no blanks, and therefore the subshifts of exceptional points in X_k for any $k \in \mathbb{N}$ coincide.)

It is easily seen that no X_k is topologically mixing, so the X_k will not satisfy 4.2(6). However without much effort it is possible to adjust the defining rules of the SFTs X_k in order to create similar \mathbb{Z}^2 shifts of finite type which are indeed topologically mixing. For this, we use another example of Hochman, which is very similar to X_k . In X_k , blank tiles were not permitted to be adjacent to each other. Define a new system \tilde{X}_k by allowing free combinations of blank tiles to appear in rectangular configurations of dimensions 1×1 , 1×2 , 2×1 , and 2×2 only. All other adjacency conditions remain unchanged. Figure 8 shows a possible configuration in \tilde{X}_k .

The level- n squares from X_k correspond to a slightly more general class of level- n rectangles in \tilde{X}_k . A level-1 rectangle of \tilde{X}_k is defined to be any rectangular configuration of blanks surrounded by a closed path of arrows (all with the same color). For $n > 1$, a pattern with rectangular shape is a level- n rectangle in \tilde{X}_k if the configuration on its boundary is a closed path of arrows with the same color and its interior is the union of four rectangular patterns, each of which is a level- m rectangle for some $m < n$, and at least one of which is of level $n - 1$. In fact, it is shown in Lemma 6.14 of [12] that the four component rectangular configurations must all be level- $(n - 1)$ rectangles. The description of the transitive (non-exceptional) points of \tilde{X}_k is similar to that of X_k ; Lemma 6.16 of [12] states that for any $\tilde{x} \in \tilde{X}_k$, either


 FIGURE 8. Part of a point of \tilde{X}_k

there exists a sequence of nested rectangles $R_n \subsetneq \mathbb{Z}^2$ ($n \in \mathbb{N}$) with $\bigcup_{n=1}^{\infty} R_n = \mathbb{Z}^2$ where $\tilde{x}|_{R_n}$ is a level- n rectangle for all $n \in \mathbb{N}$, or \tilde{x} is an exceptional point. (The set of exceptional points in \tilde{X}_k is actually the same as the set of exceptional points for X_k , again because exceptional points contain no blanks.) We are now ready to prove that for sufficiently large $k \in \mathbb{N}$, $X = \tilde{X}_k$ satisfies the claims of Theorem 4.2.

Proof. (1) In X_k , any level- n square has shape $[\vec{1}, (2^{n+1} + 2^{n-1} - 2)\vec{1}]$ and contains 4^{n-1} blank tiles, so the density of blank (\square) tiles in any $x_\omega \in X_k$ ($\omega \in \mathcal{T}^{\mathbb{N}}$) is

$$f_{\square} := \lim_{n \rightarrow \infty} \frac{4^{n-1}}{(2^{n+1} + 2^{n-1} - 2)^2} = \frac{1}{25}.$$

Therefore, since any blank tile may have k labels, it is not hard to check that $h(X_k) = \lim_{n \rightarrow \infty} \frac{\log |\mathcal{L}_{[\vec{1}, (2^{n+1} + 2^{n-1} - 2)\vec{1}]}(X_k)|}{(2^{n+1} + 2^{n-1} - 2)^2} \geq \frac{\log k}{25}$ for any positive integer $k \in \mathbb{N}$. Since $X_k \subsetneq \tilde{X}_k$, certainly $h(\tilde{X}_k) \geq h(X_k) \geq \frac{\log k}{25}$. So we may set $X = \tilde{X}_k$ with k chosen large enough that $h(X) > M$.

(2,4) These claims may be proved just as they were proved in the case $d \geq 3$. (Recall that for any $x \in X$ and $n \in \mathbb{N}$, there will exist $l \in \mathbb{N}$ and some rectangular block $B := [\vec{v}, \vec{w}]$ such that $[\vec{v} + l\vec{1}, \vec{w} - l\vec{1}]$ contains $[-n\vec{1}, n\vec{1}]$ and for all \vec{i} in the thick border $[\vec{v}, \vec{w}] \setminus [\vec{v} + l\vec{1}, \vec{w} - l\vec{1}]$, the symbol $x_{\vec{i}}$ is an arrow.)

(6) It is shown in Theorem 6.23 of [12] that \tilde{X}_1 is topologically mixing, and the proof extends trivially to any \tilde{X}_k .

To prove X has top.c.p.e., assume that $\phi : X \rightarrow Y$ is a sliding block code from X onto Y , and denote by $2m + 1$ ($m \in \mathbb{N}_0$) the size of its coding window. If Y is non-trivial, then there exist two patterns $P, P' \in \mathcal{L}_{[-m\vec{1}, m\vec{1}]}(X)$ so that $\phi(P) \neq \phi(P')$. By definition of transitivity, any transitive point of X contains both P and P' as subpatterns, and by the description of the transitive points in Lemma 6.16 of [12], for some $n \in \mathbb{N}$ there exists a level- n rectangle Q_n whose dimensions are $a_n \times b_n$ ($a_n, b_n \in \mathbb{N}$) and which contains both P and P' as subpatterns. Say that $Q_n|_{\vec{i}+[-m\vec{1}, m\vec{1}]} = P$ and $Q_n|_{\vec{j}+[-m\vec{1}, m\vec{1}]} = P'$ for some $\vec{i} \neq \vec{j} \in \mathbb{Z}^2$.

We claim that for $N \in \mathbb{N}$ large enough there exist two level- N rectangles Q_N, Q'_N both having a border colored SE , both of dimensions $a_N \times b_N$ ($a_N, b_N \in \mathbb{N}$) and such that $Q_N|_B = P$ and $Q'_N|_B = P'$ for some fixed $B \subseteq [\vec{1}, (a_N, b_N)] \subsetneq \mathbb{Z}^2$ a translated copy of $[-m\vec{1}, m\vec{1}]$. To prove this it suffices to show that we can decrease the horizontal $|\vec{j}_1 - \vec{i}_1|$ and vertical $|\vec{j}_2 - \vec{i}_2|$ distance between occurrences of P and P' in Q_n by increasing the level of the rectangle.

Without loss of generality, assume that $\vec{j}_1 > \vec{i}_1$ and that the border of Q_n is colored SE . It is shown in [12] that for any $a, b \in I_n := [2^{n+1} + 2^{n-1} - 2, 2^{n+1} + 2^n - 2] \cap \mathbb{Z}$, there exists a level- n rectangle with dimensions $a \times b$. We now construct two level- $(n+1)$ rectangles Q_{n+1} and Q'_{n+1} . Q_{n+1} has a border colored NW , and the four level- n rectangles in it have the following properties: the northwest rectangle has dimensions $a \times a$ for some allowed integers $a, a+1 \in I_n$, the southeast rectangle is Q_n , and the northeast and southwest rectangles are any with the correct dimensions, i.e. the northeast is $a_n \times a$ and the southwest is $a \times b_n$. Q'_{n+1} is described in exactly the same way as Q_{n+1} , except that the northwest subrectangle of Q'_{n+1} has dimensions $a+1 \times a$. The dimensions of Q_{n+1} are then $a+a_n+2 \times a+b_n+2$, and the dimensions of Q'_{n+1} are $a+a_n+3 \times a+b_n+2$. We finally create two level- $(n+2)$ rectangles Q_{n+2} and Q'_{n+2} in a similar fashion. Q_{n+2} has a border colored SE , and the four level- $(n+1)$ rectangles in it have the following properties: the northwest rectangle is Q_{n+1} , the southeast rectangle has dimensions $b+1 \times b+1$ for some allowed integers $b, b+1 \in I_{n+1}$, and the northeast and southwest rectangles are any with the correct dimensions. Q'_{n+2} is described in exactly the same way as Q_{n+2} , except that the southeast subrectangle of Q'_{n+2} has dimensions $b \times b+1$. Then, Q_{n+2} and Q'_{n+2} both have dimensions $a_n + a + b + 5 \times b_n + a + b + 5$. Also, it is not hard to see that there exists $B' \subseteq [\vec{1}, (a_n + a + b + 5, b_n + a + b + 5)] \subsetneq \mathbb{Z}^2$ a translated copy of $[-m\vec{1}, m\vec{1}]$ so that $Q_{n+2}|_{\vec{i}+B'} = P$ and $Q'_{n+2}|_{\vec{j}+B'-\vec{e}_1} = P'$, so the horizontal distance between subpatterns P and P' has decreased by one. By repeating this procedure, we will eventually arrive at a pair $Q_{n'}, Q'_{n'}$ of level- n' rectangles ($n' > n$), both with dimensions $a_{n'} \times b_{n'}$ ($a_{n'}, b_{n'} \in \mathbb{N}$) such that P appears in $Q_{n'}$ at the same \vec{e}_1 -coordinate as P' in $Q'_{n'}$. Similarly we may decrease the (vertical) \vec{e}_2 -coordinate, changing the role of the integers a and $a+1$ resp. b and $b+1$. After finitely many steps we arrive at level- N rectangles Q_N, Q'_N as claimed.

Now for any $N' > N$, it is possible to construct a level- N' rectangle such that all $4^{N'-N}$ of its level- N subrectangles have dimensions $a_N \times b_N$. The dimensions of this level- N' rectangle are then $2^{N'-N}a_N + 2^{N'-N+1} - 2 \times 2^{N'-N}b_N + 2^{N'-N+1} - 2$. Obviously each of the $4^{N'-N-1}$ level- N subrectangles having a border colored SE can be independently filled either by Q_N or by Q'_N . As $Q_N|_B = P$ and $Q'_N|_B = P'$

we know that $\phi(Q_N) \neq \phi(Q'_N)$ and hence all of the $2^{4^{N'-N-1}}$ level- N' rectangles createable in this manner have distinct ϕ -images. This means that the number of distinct $(2^{N'-N}a_N + 2^{N'-N+1} - 2 - 2m) \times (2^{N'-N}b_N + 2^{N'-N+1} - 2 - 2m)$ patterns in the language of Y is at least $2^{4^{N'-N-1}}$, and so the entropy of Y will be

$$\begin{aligned} h(Y) &\geq \lim_{N' \rightarrow \infty} \frac{\log(2^{4^{N'-N-1}})}{(2^{N'-N}a_N + 2^{N'-N+1} - 2 - 2m)(2^{N'-N}b_N + 2^{N'-N+1} - 2 - 2m)} \\ &= \frac{\log 2}{4a_N b_N + 8(a_N + b_N) + 16} > 0. \end{aligned}$$

It follows that every non-trivial subshift factor of X has positive entropy, and therefore no non-trivial topological factor of X has zero entropy. \square

Remark 9.1. Our techniques allow some give and take in the properties enjoyed by an example. Above, if we set X equal to X_k rather than \tilde{X}_k , then X would not be mixing and X would not be top.c.p.e.. However, $X = X_k$ would still satisfy conditions (1)-(4) of Theorem 4.2. Also (in contrast to \tilde{X}_k), X_k would satisfy (5) (i.e. X_k has no topological factor which admits a c.p.e. measure); this follows from Proposition 6.1, because X_k has the zero entropy factor X_1 and arbitrarily large blocks from the zero entropy exceptional shift E occur in every point of X_k . X_k would also have a top.c.p.e. sofic factor (constructed from X_k by sending all non-blank symbols and one of the blank symbols to 0, and leaving the other $k - 1$ blank symbols unchanged).

Alternately, using the technique which built the gently ascending hyperplanes in Section 7, we could build an SFT \tilde{X}_k^* by superimposing on \tilde{X}_k an SFT which allows the blocks of blanks with sizes 2×2 , 1×2 , 2×1 to occur only very rarely, so that $h(\tilde{X}_1^*) = 0$ even though the blank symbol will occur with positive frequency in some points. Now \tilde{X}_k^* is mixing, and has no topological factor admitting a c.p.e. invariant measure, and satisfies the other statements of Theorem 4.2, with the following exception: \tilde{X}_k^* has no factor which is top.c.p.e..

10. A TOPOLOGICALLY MIXING, TOP.C.P.E. \mathbb{Z}^2 SOFIC EXAMPLE

In this section we complete the proof of Theorem 4.1 by establishing the case $d = 2$.

Proof. Let $X = \tilde{X}_k$ be the \mathbb{Z}^2 SFT constructed in the previous section, satisfying the statement of Theorem 4.2 in the case $d = 2$. We define S as the image of X under the one-block factor map π which sends all non-blank symbols in the alphabet of \tilde{X}_k to 0 and leaves each of the k blank symbols unchanged. Hence S is a subshift with an alphabet $\tilde{\mathcal{A}}_k := \{0\} \cup \{\square_j \mid 1 \leq j \leq k\}$ and π collapses the subsystem of exceptional points $E \subsetneq \tilde{X}_k$ to a single (fixed) point. We go on now to verify the claims of Theorem 4.1.

(1, 7) We claim that $h(X) = h(S)$. As S is a factor of X we only need to show $h(X) \leq h(S)$. For this we estimate the number of patterns in $\mathcal{L}_{B_n}(X)$ in terms of image patterns in $\mathcal{L}_{B_n}(S)$ for $B_n := [\vec{1}, n\vec{1}]$ ($n \in \mathbb{N}$). Preimages of the pattern $0^{B_n} \in \mathcal{L}_{B_n}(S)$ are subpatterns of exceptional points and their number is thus bounded by $|\mathcal{A}_k|^{4n-4}$ as before (the symbols along the boundary determine the interior). Every other pattern $P \in \mathcal{L}_{B_n}(S)$ contains at least one blank and every preimage can be recreated from an n -tuple of colors $t \in \mathcal{T}^n$ as follows: Let

$s \in S$ such that $s|_{B_n} = P$. Starting at the lexicographically smallest coordinate $\vec{i} \in B_n$ containing a blank, surround it (and the adjacent blanks if any) by a level-1 border of color t_1 . Knowing t_1 we can also put down the three neighboring level-1 rectangles. In doing so, we completed the interior of the level-2 rectangle containing the blank at \vec{i} ; now surround it by a border of arrows of color t_2 . Continue in this way: for any $i > 2$, take the next color t_i ($1 < i \leq n$) from the list and surround the already placed level- $(i-1)$ rectangles with a level- i border of color t_i (we can easily determine its size from the knowledge of the positions of the blanks) and we can fill in all remaining symbols in its interior (all 4 subrectangles have to be level- $(i-1)$). In every step the region already colored extends at least 1 coordinate further into each direction than in the step before, so after at most n steps we have completely filled B_n . As every preimage of P can be obtained by this procedure, $|\pi^{-1}(P)| \leq 4^n \leq |\mathcal{A}_k|^{4n}$ and we have $|\mathcal{L}_{B_n}(X)| \leq |\mathcal{A}_k|^{4n} \cdot |\mathcal{L}_{B_n}(S)|$. Hence $h(X) \leq h(S)$, proving our claim and thus (7). (1) follows because $h(X) > M$.

(2) Because E contains every minimal subsystem of X , the image of E (which is the fixed point) contains every minimal subsystem of S .

(3,4) These items are proved as in the case $d \geq 3$. (Note that finite configurations must occur inside blocks with arbitrarily large neighborhoods of the boundary colored 0.)

(6,8) S is mixing and top.c.p.e. because it is a non-trivial factor of X which has those properties. Trivially (6) also follows. \square

Remark 10.1. Let \tilde{S}_k^* be the sofic shift which is the image under the map π above of the SFT \tilde{X}_k^* described in Remark 9.1. This shift is mixing, and satisfies the conditions satisfied by S above, except that it has no top.c.p.e. factor. However, no topological factor of \tilde{S}_k^* admits a c.p.e. invariant measure.

APPENDIX A. A HIERARCHY OF MIXING NOTIONS FOR \mathbb{Z}^d SHIFTS

In this appendix, we collect and put into a more general framework some mixing properties for \mathbb{Z}^d shifts dispersed throughout the literature [9, 14, 16, 17, 24, 25, 27], as well as the block gluing notion (3.1) of the current paper. We record some known and some new facts about these mixing properties. We show that all of them can be put into a linear chain, in which – with one possible exception – none of the logical implications can be reversed (A.6).

Recall that δ_∞ represents the maximum-metric on \mathbb{Z}^d , i.e. $\delta_\infty(\vec{v}, \vec{w}) = \|\vec{v} - \vec{w}\|_\infty := \max_{1 \leq k \leq d} |v_k - w_k|$ and also its natural extension to a separation function on non-empty subsets $V, W \subseteq \mathbb{Z}^d$: $\delta_\infty(V, W) := \min_{\vec{v} \in V, \vec{w} \in W} \delta_\infty(\vec{v}, \vec{w})$.

We start with two common definitions from the \mathbb{Z}^d shift literature.

Definition A.1. A \mathbb{Z}^d shift X is *topologically mixing* if for any two non-empty finite subsets $V, W \subsetneq \mathbb{Z}^d$ there exists a constant $D_{V,W} \in \mathbb{N}_0$ so that for any $\vec{i} \in \mathbb{Z}^d$ such that $\delta_\infty(V, W + \vec{i}) > D_{V,W}$, any pair of patterns on V and $W + \vec{i}$ occurring in X can be put together to form a valid point of X , i.e. $\forall x, y \in X, \exists z \in X$ such that $z|_V = x|_V$ and $z|_{W+\vec{i}} = y|_{W+\vec{i}}$.

Furthermore remember that $\vec{v} \preceq \vec{w}$ is a shorthand for the fact that the components of $\vec{v}, \vec{w} \in \mathbb{Z}^d$ satisfy the inequality $v_k \leq w_k$ for all $1 \leq k \leq d$. As in Section 3, the notation $B = [\vec{v}, \vec{w}] := \{\vec{i} \in \mathbb{Z}^d \mid \vec{v} \preceq \vec{i} \preceq \vec{w}\}$ will be used for solid blocks in \mathbb{Z}^d and $\vec{1}$ refers to the element in \mathbb{Z}^d in which all components equal 1.

Definition A.2. [14] A \mathbb{Z}^d shift X is *corner gluing at gap $g \in \mathbb{N}_0$* if for any $\vec{u}, \vec{v}, \vec{w} \in \mathbb{Z}^d$ with $\vec{u} + \vec{1} \preceq \vec{v} \preceq \vec{w} - g\vec{1}$, any patterns on $L := [\vec{u}, \vec{w}] \setminus [\vec{v}, \vec{w}]$ – which we will call the corner-shape – and a solid block $R := [\vec{v} + g\vec{1}, \vec{w}]$ which are legal in X can be put together in the same point of X , i.e. $\forall x, y \in X, \exists z \in X$ such that $z|_L = x|_L$ and $z|_R = y|_R$.

Next we introduce the general concepts of \mathcal{F} -gluing and \mathcal{F} -filling and exhibit some special cases of these new notions which are already present in the literature.

Definition A.3. Let $\mathcal{F} \subseteq \{F \subsetneq \mathbb{Z}^d \mid F \neq \emptyset \wedge F \text{ finite}\}$ be a family of non-empty finite subsets in \mathbb{Z}^d .

A \mathbb{Z}^d shift X is called *\mathcal{F} -gluing at gap $g \in \mathbb{N}_0$* if for any two elements $V, W \in \mathcal{F}$ with separation $\delta_\infty(V, W) > g$ and any two patterns on V and W which occur in X there is a valid point in X realizing these two patterns, i.e. $\forall x, y \in X, \exists z \in X$ such that $z|_V = x|_V$ and $z|_W = y|_W$.

Similarly a \mathbb{Z}^d shift X is called *\mathcal{F} -filling at gap $g \in \mathbb{N}_0$* if for any element $V \in \mathcal{F}$ and any (possibly infinite) subset $\emptyset \neq W \subseteq \mathbb{Z}^d$ with separation $\delta_\infty(V, W) > g$ every pair of patterns on V and W which occur in X can be put together to form a valid point in X , i.e. $\forall x, y \in X, \exists z \in X$ such that $z|_V = x|_V$ and $z|_W = y|_W$.

A \mathbb{Z}^d shift X is called *\mathcal{F} -gluing resp. \mathcal{F} -filling* if it is \mathcal{F} -gluing at gap g resp. \mathcal{F} -filling at gap g for some $g \in \mathbb{N}_0$.

Observation A.4. Note that by compactness the notions of \mathcal{F} -gluing at gap g and \mathcal{F} -filling at gap g coincide for $\mathcal{F} = \{F \subsetneq \mathbb{Z}^d \mid F \neq \emptyset \wedge F \text{ finite}\}$ the collection of all finite subsets. \mathbb{Z}^d shifts satisfying this complete gluing/filling condition are usually (following [6]) called strongly irreducible (SI) and we will use this terminology for the two equivalent properties.

However, for general sets \mathcal{F} the notions of \mathcal{F} -gluing and \mathcal{F} -filling need not coincide, as we can check with the following:

Observation A.5. For $\mathcal{B} := \{B = [\vec{v}, \vec{w}] \mid \vec{v} \preceq \vec{w} \in \mathbb{Z}^d\}$ the collection of all solid blocks, the notion \mathcal{B} -gluing at gap g becomes block gluing at gap g as defined in Section 3.

The property of being \mathcal{B} -filling already appears in the literature as the uniform filling property (UFP). For simplicity we will also use this common expression. As we will see in Theorem A.6, block gluing and having the UFP (i.e. being block filling) are not equivalent for \mathbb{Z}^d shifts.

Theorem A.6. For $g \in \mathbb{N}_0$ and a \mathbb{Z}^d shift X we have the following implications.

X is SI at gap $g \xrightarrow{(1)} X$ has the UFP at gap $g \xrightarrow{(2)} X$ is corner gluing at gap $g \xrightarrow{(3)} X$ is block gluing at gap $g \xrightarrow{(4)} X$ is topologically mixing.

Implications (2), (3), and (4) cannot be reversed even if we allow an increase in g .

Proof. Implications (1), (2) and (4) are fairly clear from the definitions.

We prove Implication (3). Let the \mathbb{Z}^d shift X be corner gluing at gap $g \in \mathbb{N}_0$ and let $B_1 = [\vec{v}^{(1)}, \vec{w}^{(1)}], B_2 = [\vec{v}^{(2)}, \vec{w}^{(2)}] \subsetneq \mathbb{Z}^d$ be two disjoint solid blocks with separation $\delta_\infty(B_1, B_2) > g$. There must be some $1 \leq j \leq d$ such that the intervals $[(\vec{v}^{(1)})_j, (\vec{w}^{(1)})_j], [(\vec{v}^{(2)})_j, (\vec{w}^{(2)})_j] \subsetneq \mathbb{Z}$ have separation larger than g ; otherwise, since each B_i ($i = 1, 2$) is a product $\prod_{1 \leq k \leq d} [(\vec{v}^{(i)})_k, (\vec{w}^{(i)})_k]$, we could find points in B_1

and B_2 at a distance at most g . Without loss of generality, suppose $(\vec{w}^{(1)})_j + g < (\vec{v}^{(2)})_j$. Define $L := \mathbb{Z}^d \cap (\prod_{k=1}^d [\min\{\vec{w}_k^{(1)}, \vec{v}_k^{(2)}\} - g - 1, \max\{\vec{w}_k^{(1)}, \vec{w}_k^{(2)}\}]) \setminus (\prod_{k=1}^d [\vec{v}_k^{(2)} - g, \max\{\vec{w}_k^{(1)}, \vec{w}_k^{(2)}\}])$ and $R := \mathbb{Z}^d \cap \prod_{k=1}^d [\vec{v}_k^{(2)}, \max\{\vec{w}_k^{(1)}, \vec{w}_k^{(2)}\}]$. By corner gluing, any patterns on L and R which are legal in X can occur in the same point. Because $B_1 \subseteq L$ and $B_2 \subseteq R$, the same holds for B_1 and B_2 .

It will be shown in Example C.1 that corner gluing (at gap 2) does not imply the UFP, and so Implication (2) cannot be reversed.

Block gluing at gap 1 does not imply corner gluing, as will be shown in Example A.7 below. Therefore, Implication (3) cannot be reversed.

Finally, we will see in Appendix B that non-trivial block gluing \mathbb{Z}^d shifts have positive entropy (B.2), and so since there exist topologically mixing shifts with zero entropy (for instance, those in [24]), topological mixing does not imply block gluing at any gap either. Thus Implication (4) cannot be reversed. \square

We point out that whether or not SI and UFP define the same notion seems to be still an open question.

The following example shows that the condition of being block gluing is in fact strictly weaker than corner gluing.

Example A.7. Consider the SFT $X \subseteq \{0, 1\}^{\mathbb{Z}^2}$ with forbidden pattern $Q = \begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix}$. In other words, $x \in \{0, 1\}^{\mathbb{Z}^2}$ is in X if and only if x does not contain Q as a subword. We claim that X is block gluing at gap 1, but not corner gluing at any gap.

Take any pair of solid blocks $B_1, B_2 \subsetneq \mathbb{Z}^2$ with $\delta_\infty(B_1, B_2) > 1$, and any pair of patterns $P_1 \in \mathcal{L}_{B_1}(X)$ on B_1 and $P_2 \in \mathcal{L}_{B_2}(X)$ on B_2 . We claim that the point $x \in \{0, 1\}^{\mathbb{Z}^2}$ defined by $x|_{B_1} := P_1$, $x|_{B_2} := P_2$, and $x_{\vec{i}} := 1$ for all $\vec{i} \notin B_1 \cup B_2$ is in X . All that must be checked is that x does not contain Q as a subword. Consider any 2×2 solid block $B \subsetneq \mathbb{Z}^2$. There are two cases. If $B \subseteq B_1$ or $B \subseteq B_2$, then $x|_B$ is a subword of P_1 or P_2 and therefore does not equal Q . If $B \not\subseteq B_1$ and $B \not\subseteq B_2$, then $|B \setminus (B_1 \cup B_2)| \geq 2$. Therefore, $x|_B$ contains at least two 1s, and is again not equal to Q . So, x does not contain Q as a subword, and $x \in X$. This means that X is block gluing at gap 1.

Now consider, for any $g \in \mathbb{N}_0$, take the solid blocks $B_1 = [\vec{0}, (g+1)\vec{1}]$, $B_2 = [\vec{1}, (g+1)\vec{1}]$, and define the set $L := B_1 \setminus B_2 \subsetneq \mathbb{Z}^2$. The pattern $P := 0^L$ on L consisting of all 0s is certainly valid for X – for instance, it is a subword of the point $0^{\mathbb{Z}^2}$, which is in X . However, we claim that if P appears in some point $x \in X$, i.e. there exists $\vec{i} \in \mathbb{Z}^2$ so that $x|_{\vec{i}+L} = P$, this actually forces the entire solid block on whose lower-left boundary the pattern P occurs to be filled with 0s, i.e. $x|_{\vec{i}+B_1} = 0^{B_1}$. Suppose this is not the case. Then there would exist $x \in X$ and $\vec{i} \in \mathbb{Z}^2$ so that $x|_{\vec{i}+L} = P$ and $x|_{\vec{i}+B_1}$ contains a 1. Consider the lowest row of $x|_{\vec{i}+B_1}$ containing a 1 and then take the leftmost 1 in this row. The 2×2 solid block with this 1 in its upper-right corner will have 0s in its other locations, and so equals Q . This cannot happen since $x \in X$, so we have a contradiction and in fact $x|_{\vec{i}+B_1}$ has all 0s. Now define the singleton set $R := \{(g+1)\vec{1}\} \subsetneq \mathbb{Z}^2$, and let the pattern on R be just a single 1. Certainly the pattern 1 appears in X , e.g. in the point $1^{\mathbb{Z}^2}$. However, there does not exist $x \in X$ so that $x|_L = P$ and $x|_R = 1$. As $\delta_\infty(L, R) > g$ this means that X is not corner gluing at gap g , and since g was arbitrary, X is not corner gluing (at any gap-size).

APPENDIX B. BLOCK GLUING \mathbb{Z}^d SHIFTS

In this section we consider some consequences of the very natural block gluing condition, some of which are used in the proof of Theorems 3.2, 4.1 and 4.2.

Lemma B.1. *Let $d \geq 1$, $g \in \mathbb{N}_0$ and let X be a \mathbb{Z}^d shift. Then the following statements are equivalent.*

- (1) X is block gluing at gap g .
- (2) For $n \in \mathbb{N}_0$ and any countable family $(P_{\vec{v}} \in \mathcal{L}_{[\vec{1}, n\vec{1}]}(X))_{\vec{v} \in (g+n)\mathbb{Z}^d}$ of individually valid patterns there exists a point $x \in X$ such that $x|_{\vec{v} + [\vec{1}, n\vec{1}]} = P_{\vec{v}}$ for all $\vec{v} \in (g+n)\mathbb{Z}^d$.

Proof. (2) \implies (1) is trivial, and the converse is a straightforward compactness argument. \square

Theorem B.2. *Suppose $d \geq 1$ and let X be a non-trivial block gluing \mathbb{Z}^d shift. Then the following hold.*

- (1) $h(X) > 0$.
- (2) Every subshift factor of X is again block gluing.
- (3) X has topologically completely positive entropy.
- (4) For every $\varepsilon > 0$, X contains a proper subshift $W \subsetneq X$ such that $h(W) > h(X) - \varepsilon$.
- (5) For every non-empty proper subsystem $V \subsetneq X$ and every $\varepsilon > 0$ there is a non-empty subshift $W \subsetneq X$ disjoint from V with $h(W) > h(X) - \varepsilon$.

Proof. Assume that X is block gluing at gap $g \in \mathbb{N}_0$.

(1) This follows easily from condition (2) of Lemma B.1.

(2,3) Suppose $\phi : X \rightarrow Y$ is a factor map between subshifts, such that $x|_{[-n\vec{1}, n\vec{1}]}$ ($n \in \mathbb{N}_0$) determines $(\phi(x))_0$. Patterns on Y -blocks with separation $g + 2n$ are the images of patterns on X -blocks with separation g , and therefore Y is block gluing at gap $g + 2n$. (3) now follows from (1) and (2).

(5) Let $P \in \mathcal{L}(X) \setminus \mathcal{L}(V)$ be a pattern appearing in some point in X but not in any point of the proper subsystem V ; w.l.o.g. we assume $P \in \mathcal{L}_{[\vec{1}, n\vec{1}]}(X)$ for some fixed $n \in \mathbb{N}$. Let $\mathcal{P}_N := \{Q \in \mathcal{L}_{[\vec{1}, N\vec{1}]}(X) \mid Q|_{[\vec{1}, n\vec{1}]} = P\}$ with $N \geq n$. By block gluing, every element of $\mathcal{L}_{[\vec{1}, (N-n-g)\vec{1}]}(X)$ occurs as the upper corner subblock of some element of \mathcal{P}_N , so $|\mathcal{P}_N| \geq |\mathcal{L}_{[\vec{1}, (N-n-g)\vec{1}]}(X)|$. Given $N > g$, let $B := [\vec{1}, (2N+g)\vec{1}]$ and define a subshift

$$W_N := \{x \in X \mid \forall \vec{v} \in \mathbb{Z}^d : \text{the block } x|_{\vec{v}+B} \text{ contains the subpattern } P\}.$$

Clearly W_N is disjoint from V . By condition (2) of Lemma B.1, patterns from \mathcal{P}_N may appear independently on the lattice $(g+N)\mathbb{Z}^d$ in points of X . Such points are in W_N and it follows that $\lim_{N \rightarrow \infty} h(W_N) = h(X)$.

(4) We will show that X contains a proper subshift, and then (5) will imply (4). Given $N \in \mathbb{N}$, pick a pattern $P \in \mathcal{L}_{[\vec{1}, N\vec{1}]}(X)$, and let U_N be the set of points x in X such that $x|_{\vec{v} + [\vec{1}, N\vec{1}]} = P$ for all $\vec{v} \in (g+N)\mathbb{Z}^d$. By condition (2) of Lemma B.1, U_N is non-empty. Also U_N is closed, and invariant under shifts from $N\mathbb{Z}^d$. Thus the \mathbb{Z}^d shift orbit closure V_N of U_N is the union of images of U_N under finitely many (i.e. N^d) shifts. Clearly $\lim_{N \rightarrow \infty} h(V_N) = 0$, and for large N it follows that V_N is a proper subshift of X . \square

We are interested in the implications of block gluing (or other mixing conditions) on richness and separation properties of subsystems. Within a mixing \mathbb{Z} SFT X of positive entropy, any two disjoint subsystems are contained in disjoint mixing subSFTs of X , with entropy arbitrarily close to that of X . This technical fact is tremendously useful for constructions with \mathbb{Z}^d SFTs when $d = 1$ (e.g. see [7]), but it fails for $d > 1$ [24]. As part of the effort to understand the boundaries of how properties do and do not extend to higher dimensional SFTs, we ask the following.

Question B.3. Given $\varepsilon > 0$ and $d \geq 2$, is any pair of disjoint subsystems V_1, V_2 inside a block gluing \mathbb{Z}^d SFT X contained in a pair $W_1 \supseteq V_1, W_2 \supseteq V_2$ of disjoint block gluing subSFTs of entropy at least $h(X) - \varepsilon$?

For $d = 1$, the block gluing SFTs are the mixing SFTs, and the answer to Question B.3 is well known to be yes [7, Section 26].

For $d = 1, 2$ (but not for $d \geq 3$), it is known that finite orbits of a \mathbb{Z}^d strongly irreducible SFT are dense [30]. The same argument works for block gluing SFTs, and we have the same question.

Question B.4. Are there block gluing \mathbb{Z}^d SFTs (for $d \geq 3$) that do not have dense periodic points? Are there examples with no periodic points?

From [24] we know that every \mathbb{Z}^d SFT having the UFP is entropy minimal (i.e., every proper subsystem has strictly smaller entropy). We note that Example C.1 provides the existence of non-trivial block gluing \mathbb{Z}^d SFTs ($d \geq 2$) which are not entropy minimal.

APPENDIX C. A NON-ENTROPY-MINIMAL 4-CORNER-GLUING \mathbb{Z}^2 SFT

This appendix contains a \mathbb{Z}^2 SFT which is corner gluing in all 4 corners, i.e. in the NW-, NE-, SW- and SE-corners, but which still does not have the UFP. With a small adjustment we can even violate entropy minimality. In particular this example proves that implication (2) in Theorem A.6 cannot be reversed.

The \mathbb{Z}^2 SFT we construct models a system of meandering streams running in a \mathbb{Z}^2 plane that has a large slope in north-south direction, but no slope in east-west direction. Therefore a stream starting somewhere north has to constantly run south (it will never run uphill), with some possible meandering in east-west direction. It can split into different “branches” and those branches can unify again. However there are neither sources nor sinks, thus a stream has to run forever without starting or ending at a certain coordinate in \mathbb{Z}^2 .

Example C.1 (The meandering streams SFT). The formal definition of the meandering streams shift X_{MS} uses an alphabet of 16 symbols as displayed in Figure 9. We will refer to symbol 1 as the blank symbol and to symbols 2 up to 16 as the non-blanks.

Now the 16 symbols can be placed next to each other only in a way that conserves the current indicated by the direction of the arrows. For example a symbol with an arrow pointing east resp. west at its right edge can only sit to the left of a symbol also having an arrow pointing east resp. west at its left edge and a symbol with an arrow ending on its lower edge can sit above a symbol if and only if this has an arrow starting from its upper edge.

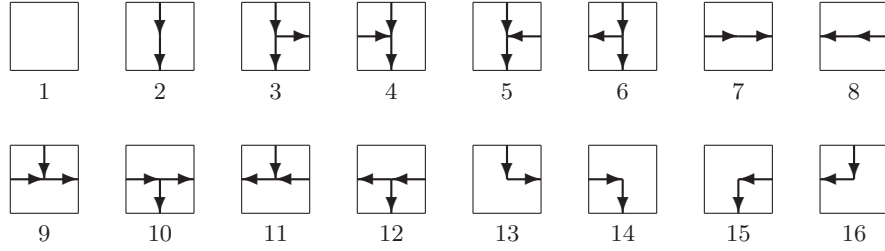


FIGURE 9. The alphabet of the meandering streams shift X_{MS}

Posing these nearest-neighbor restrictions as local rules we obviously define a non-trivial \mathbb{Z}^2 SFT. Note that all allowed transitions are realized in some point, i.e. the corresponding transition matrices are already strictly essential. Next we prove the desired properties of our example.

Claim C.2. The meandering streams SFT X_{MS} is corner gluing at gap 2 in each of the 4 corners.

Proof. Since X_{MS} is a nearest-neighbor SFT and the horizontal resp. vertical transitions only depend on the arrows that are (are not) present at the vertical resp. horizontal borders of each symbol we may prove the corner gluing property by filling the gap of width 2 between arbitrary configurations along the borders of the rectangle and the corner-shape region that appear in the definition of corner gluing.

In fact we will give a procedure to fill the gap between the following two infinite regions forming a NW-corner in \mathbb{Z}^2 with parts of either two or three streams. Let $a, b \in \mathbb{Z}$, define an infinite corner-shape region $L' := \{\vec{u} \in \mathbb{Z}^2 \mid u_1 < a - 2 \vee u_2 > b + 2\}$ and an infinite rectangle $R' := \{\vec{v} \in \mathbb{Z}^2 \mid v_1 \geq a \wedge v_2 \leq b\}$ and let $x, y \in X_{MS}$ be two arbitrary points. Now we construct a new point $z \in X_{MS}$ such that $z|_{L'} = x|_{L'}$ and $z|_{R'} = y|_{R'}$. The coordinates in $\mathbb{Z}^2 \setminus (L' \cup R')$ are filled with symbols as follows.

First we put $z|_{(a-1, b+1)} = 15$ and for every $i \in \mathbb{N}$ we put $z|_{(a-1, b+1-i)} \in \{2, 3, 5\}$ and $z|_{(a-1+i, b+1)} \in \{8, 12\}$. In this manner we have already nurtured all currents on the top-border of R' and nurtured/drained all currents on the left border of R' . Second we assign a pattern to the remaining coordinates depending on the two symbols sitting to the left and above the coordinate $(a - 2, b + 2)$.

If $z|_{(a-3, b+2)} \in \{1, 2, 4, 6, 14, 16\}$, we use the first drawing of Figure 10, i.e. we put $z|_{(a-2, b+2)} \in \{5, 15\}$ and for all $i \in \mathbb{N}$ we put $z|_{(a-2, b+2-i)} \in \{2, 4, 6\}$ and $z|_{(a-2+i, b+2)} \in \{8, 11\}$.

If $z|_{(a-3, b+2)} \in \{3, 7, 9, 10, 13\}$, we use the second drawing of Figure 10, i.e. we put $z|_{(a-2, b+2)} \in \{4, 14\}$ and for all $i \in \mathbb{N}$ we put $z|_{(a-2, b+2-i)} \in \{2, 4, 6\}$. Let $i' := \min\{i \in \mathbb{N} \mid z|_{(a-2+i, b+3)} \in \{2, 3, 4, 5, 6, 10, 12, 14, 15\}\}$, then for every $i \in \mathbb{N}$ with $i < i'$ put $z|_{(a-2+i, b+2)} = 1$, put $z|_{(a-2+i', b+2)} = 13$ and for $i > i'$ put $z|_{(a-2+i, b+2)} \in \{7, 9\}$.

Similarly if $z|_{(a-3, b+2)} \in \{5, 8, 11, 12, 15\}$ and $z|_{(a-2, b+3)} \in \{1, 7, 8, 9, 11, 13, 16\}$, we use the third drawing of Figure 10 and finally if $z|_{(a-3, b+2)} \in \{5, 8, 11, 12, 15\}$ and $z|_{(a-2, b+3)} \in \{2, 3, 4, 5, 6, 10, 12, 14, 15\}$, we use the fourth drawing of Figure 10.

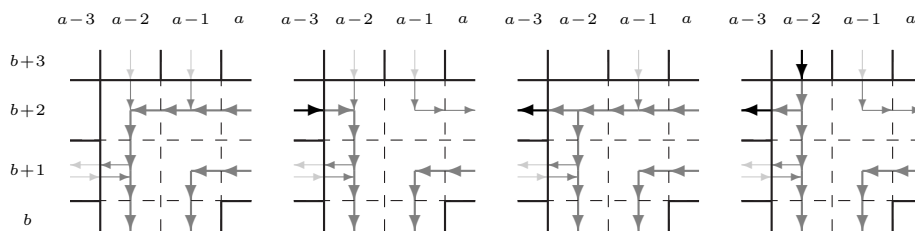


FIGURE 10. Corner gluing in the NW-corner

In all four cases we have constructed a new point $z \in X_{\text{MS}}$ which combines the infinite patterns $x|_{L'}$ and $y|_{R'}$. Obviously this implies corner-gluing at gap size 2 for the NW-corner, as every finite NW-corner-gluing configuration (on finite regions L and $R = [\vec{v}, \vec{w}]$) of gap size at least two can be reduced to the case treated above by choosing $a = v_1, b = w_2$ and extending L to $L' \supseteq L$ and R to $R' \supseteq R$ (keeping larger parts of the points x, y).

As the situation is completely symmetric in the east-west direction one can flip the drawings in Figure 10 about their north-south axis to get appropriate fillings for the NE-corner. Similarly, because the situation is almost symmetric (except for the vertical current of a stream) in the north-south direction one can produce corresponding fillings for the SW- resp. SE-corner by flipping the drawings in Figure 10 resp. those obtained for the NE-corner about their east-west axis and reversing the direction of the current in all vertical parts afterwards. \square

Claim C.3. The meandering streams SFT X_{MS} does not have the UFP.

Proof. Note that X_{MS} contains the fixed point of all blanks. However as a stream cannot stop at any coordinate in \mathbb{Z}^2 there are no points seeing non-blank symbols only at a finite subset of \mathbb{Z}^2 . To show this suppose there exists a point $x \in X_{\text{MS}}$ and a non-empty finite subset $F \subsetneq \mathbb{Z}^2$ of coordinates such that $x|_{\vec{v}} \neq 1$ for all $\vec{v} \in F$ but $x|_{\vec{v}} = 1$ for all $\vec{v} \in \mathbb{Z}^2 \setminus F$. Let $b := \max\{r \in \mathbb{Z} \mid \mathbb{Z} \times \{r\} \cap F \neq \emptyset\}$ be the top-most row in \mathbb{Z}^2 containing elements of F and let $a := \min\{c \in \mathbb{Z} \mid (c, b) \in F\}$ be the left-most coordinate in this row. Now the symbol at coordinate (a, b) has to be a non-blank sitting beneath a blank and also sitting to the right of a blank symbol. Checking the alphabet of X_{MS} we find that $x|_{(a,b)} = 15$ – the only non-blank symbol with no arrow on its top nor its left border. Since by definition of b all symbols in row $\mathbb{Z} \times \{b+1\}$ have to be blanks, the only possibilities for $x|_{(a+1,b)}$ are the symbols 8 or 12. Continuing in this fashion we see that $x|_{(a+i,b)} \in \{8, 12\}$ for all $i \in \mathbb{N}$. Therefore $(a + \mathbb{N}_0) \times \{b\} \subseteq F$ which contradicts the finiteness of F .

This proves that in X_{MS} there is a unique way to fill in a rectangular frame of blanks, namely with all blanks. Thus it is not possible for any $g \in \mathbb{N}_0$ to glue together a non-blank symbol at the origin $U := [\vec{0}, \vec{0}] = \{\vec{0}\}$ to a pattern of only blank symbols on $V := \mathbb{Z}^2 \setminus [-g\vec{1}, g\vec{1}]$ and thus X_{MS} does not have the UFP. \square

Now we modify the meandering streams shift by splitting the blank symbol into $N \in \mathbb{N}$ distinct but completely interchangeable copies. This modified meandering

streams SFT $X_{\text{MS}}^{(N)}$ is still corner gluing at gap 2 in each of the 4 corners and still does not have the UFP. Moreover for N large enough it is not entropy minimal.

Claim C.4. For $N \geq 120$ the modified meandering streams SFT $X_{\text{MS}}^{(N)}$ contains a proper subsystem of full entropy.

Proof. For $n \in \mathbb{N}$ we estimate the number of patterns in $\mathcal{L}_{[\vec{1}, n\vec{1}]}(X_{\text{MS}}^{(N)})$ containing a fixed number $0 \leq m \leq n^2$ of non-blank symbols. Those patterns are uniquely determined by specifying the finite parts of streams intersecting the solid block $B := [\vec{1}, n\vec{1}]$ and putting in arbitrary blanks in the remaining positions. The key point in our argument is the fact that every non-blank in a point of $X_{\text{MS}}^{(N)}$ is part of some infinite stream and those have to hit the boundary of B at least twice (once entering, once leaving).

This enables us to design an algorithm which allows a Turing machine to recreate every single pattern in $\mathcal{L}_{[\vec{1}, n\vec{1}]}(X_{\text{MS}}^{(N)})$ given a specific piece of input consisting of a finite list $L \subseteq B$ of coordinates, an m -tuple \mathcal{I} of instructions taken from a set of 120 different commands and an $(n^2 - m)$ -tuple $\mathcal{N} \in \{1, 2, \dots, N\}^{n^2 - m}$ of numbers specifying blank symbols. The list L is a subset of the boundary $\partial B := B \setminus [2\vec{1}, (n-1)\vec{1}] \subsetneq \mathbb{Z}^2$ of B that contains a starting coordinate for every connected finite part of a stream seen in B . Let L be given an ordering by the lexicographic ordering on \mathbb{Z}^2 . Every instruction in \mathcal{I} is itself a 3-tuple (a_i, f_i, d_i) ($1 \leq i \leq m$) where $a_i \in \{2, 3, \dots, 16\}$ represents one of the 15 non-blanks in the alphabet of $X_{\text{MS}}^{(N)}$, $f_i \in \{0, 1\}$ is a flag that signals either “revert” or “continue” and $d_i \in \{\text{up, down, right, left}\}$ is one of the four standard directions in \mathbb{Z}^2 . Now the Turing machine processes its input and puts down symbols on the solid block B as follows: The machine starts by moving its writing-head to the coordinate given by the first entry in the list L (if L is empty, the machine skips the first part of the algorithm and continues with placing the blank symbols, see below). There it puts down the symbol a_1 from the first instruction in \mathcal{I} starting a finite part of some stream. If f_1 is “continue” it moves its writing-head one step in the direction given by d_1 where it executes the next instruction in the same manner. If some f_i is “revert” the machine moves back along the non-blank symbols written so far until it comes to the first junction – one of the symbols $\{3, 4, 5, 6, 9, 10, 11, 12\}$ as in Figure 9 – where one of the three branches is a dead-end (i.e. the branch points to a place still inside B but where the machine has not already placed another non-blank symbol). From there the machine moves one step in the direction specified by d_i and continues with the $(i + 1)$ th instruction. If there is no dead-end, the machine moves its writing-head to the next coordinate from the list L , where it starts another stream using the next instruction from \mathcal{I} . After executing all commands in \mathcal{I} , the machine has placed exactly m non-blanks. It then jumps to coordinate $\vec{1}$, from which it starts filling in the remaining $n^2 - m$ coordinates of B with blank symbols as specified by the numbers in \mathcal{N} .

As every pattern in $\mathcal{L}_{[\vec{1}, n\vec{1}]}(X_{\text{MS}}^{(N)})$ can be created by our Turing machine using a particular input, the number of different “programs” gives an upper bound on the number of elements in $\mathcal{L}_{[\vec{1}, n\vec{1}]}(X_{\text{MS}}^{(N)})$ and thus on the entropy of $X_{\text{MS}}^{(N)}$. Overestimating the number of different lists L with the number of subsets of the boundary ∂B , summing over programs with exactly m instructions ($0 \leq m \leq n^2$), and assuming

$N \geq 120$, we get

$$|\mathcal{L}_{[\bar{1}, n\bar{1}]}(X_{\text{MS}}^{(N)})| \leq 2^{|\partial B|} \cdot \sum_{m=0}^{n^2} (120^m \cdot N^{(n^2-m)}) \leq 2^{4n-4} \cdot (n^2 + 1) \cdot N^{n^2}.$$

As a lower bound we take the number of different patterns containing only blanks, so $|\mathcal{L}_{[\bar{1}, n\bar{1}]}(X_{\text{MS}}^{(N)})| \geq N^{n^2}$. Putting both bounds together we see that $h(X_{\text{MS}}^{(N)}) = \log N$ for all $N \geq 120$. As $X_{\text{MS}}^{(N)}$ contains the full \mathbb{Z}^2 shift on N (blank) symbols as a proper subshift, we are done. \square

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