

Majority Vote Processes on Trees

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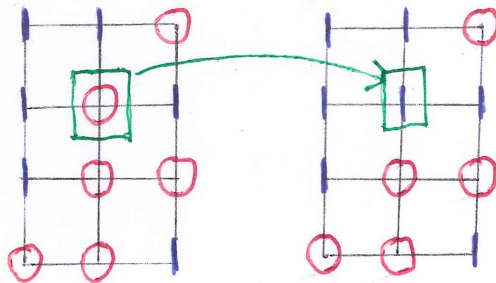
Majority Vote Process (MVP) - one of the classical interacting particle systems

others include voter model, contact process, exclusion process

The MVP is the continuous time Markov process such that:

- On \mathbb{Z}^d , each site has "opinion" either 0 or 1. $d \geq 1$
- Independently, at rate $1-\epsilon$ at each site, sites align w/ the majority opinion of their immediate neighbors. $\epsilon \in [0, 1]$
 \leftarrow MV point
- Independently, at rate ϵ at each site, the opinion "flips".
 \leftarrow noise point

Are primarily interested in case where ϵ is small, but positive.



2

How many equilibria does MVP have?

- Is open except in $d=1$.

Only 1 equilibrium (Gray (1992))

Compare with:

VM when $E=0$

Voter Model with Noise (VMN) -

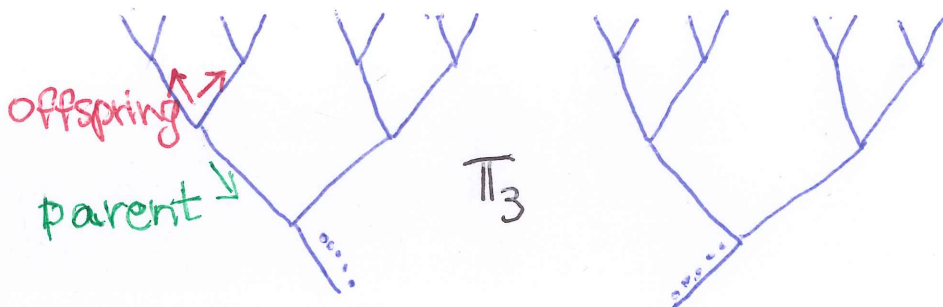
- Same as MVP, except sites randomly choose one of neighbors with which to align.

Not difficult to check:

uses duality

- VMN converges exp'l quickly to its unique equilibrium.

Here, will investigate behavior of MVP on d -regular trees \mathbb{T}_d .



Results will discuss:

↑ (B.-Gray (2018))

Theorem 1. For $d \geq 5$ and $\epsilon > 0$ small, there exist uncountably many mutually singular equilibria.

↖ will spend more time on
↙ less time on

Theorem 2. For any initial state close enough to such an equilibrium, process converges exp'l quickly to this equilibrium.

Comments:

- Method does not depend on precise model. Includes, for example, stochastic Ising model on \mathbb{T}^d .
- Argument simplifies somewhat for oriented MVP, where above results are true for $d \geq 4$.
↖ parents depend only on offspring
- Gandolfo, Ruiz, Shlosman (2012) constructed uncountably many pure Gibbs states for the Ising model on trees.

Ideas behind proof of Theorem 1

related ideas for Theorem 2,
but more involved

Notation

$(\xi_t^J)_{t \geq 0}$ \leftarrow initial state is MVP w/ $\xi_t^J(x) = 0$ or 1 for $x \in \mathbb{T}^d$

Will consider special J defined using a tiling:

- Partition \mathbb{T}^d into even/odd sites.

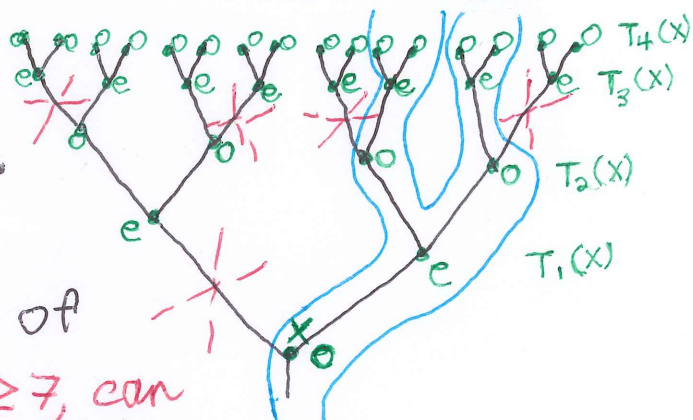
- From each odd site, delete an edge to one of offspring.

\leftarrow For $d \geq 7$, can instead delete an edge from all sites.

- Each connected subset is a tile; corresponding collection is a tiling.

- J is compatible (wrt a

given tiling) if $x, y \in T$ implies $J(x) = J(y)$.



$T_n(x)$ = set of n^{th} generation offspring of x in T

Note: For a given tiling,

(1) there are uncountably many compatible J .

Reasoning for Theorem 1

For $d \geq 5$ and small noise ε , we will show that, for each tile T ,

$$(2) \quad \overline{\lim}_{n \rightarrow \infty} \frac{1}{\#T_n(x)} \sum_{y \in T_n(x)} \left| \sum_t \xi_t^J(y) - J(y) \right| \leq 10\varepsilon \text{ a.s.}$$

for each $x \in T$ and t .

← assumed compatible wrt tiling
says ξ_t is close to J for all t .

Assume (2)

Let

$\mu_t^J =$ Cesaro avg. over $[0, t]$ of measures of ξ_s^J .

One has

$\mu_{t_i}^J \xrightarrow{i \rightarrow \infty} \mu^J$ on some subsequence $t_i \rightarrow \infty$,

where μ^J satisfies analog of (2).

Consequently, for $J \neq J'$ (w/ both compatible wrt tiling),

(3) μ^J and $\mu^{J'}$ are mutually singular.

Theorem 1 follows from (1) and (3).

Demonstrate (2)

To demonstrate (2), construct a process δ_t on $\{0,1\}^{\mathbb{T}^d}$ s.t.

$$(4) \left| \mathbb{E}_t^J(y) - J(y) \right| \leq \delta_t(y) \quad \text{for all } y$$

with

$$(2') \quad \overline{\lim}_{n \rightarrow \infty} \frac{1}{\#T_n(x)} \sum_{y \in T_n(x)} \delta_t(y) \leq 10\epsilon.$$

Construction of δ_t :

- $\delta_0 \equiv 0$.
 - $\delta_t(x) = 1$ at noise points (x,t) .
 - δ_t behaves like \mathbb{E}_t^J at MV points (x,t) , except that neighborhood is now $T_1(x)$, and, for $\delta_t(x) = 1$ to occur, 1 (2) fewer 1's required if x is even (odd).
- only points where change of state occurs

These properties imply (4).

They also imply siblings evolve independently of one another.

← same parent

To demonstrate (2'), set

$$P_t = P(\delta_t(x) = 1) \quad \text{for } x \in \Pi_d.$$

P_t depends on whether x is even or odd.

To simplify computations, instead delete an edge from all sites (not just odd).

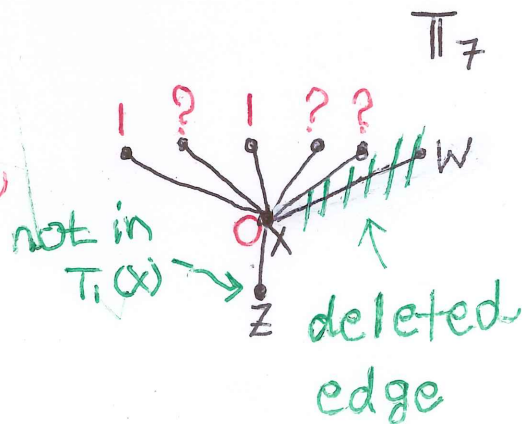
This suffices for $d \geq 7$.

It follows that, for $d \geq 7$,

$$P_t' \leq \underbrace{\varepsilon (1 - P_t)}_{\substack{\text{due to noise} \\ \text{due to rate} \\ \text{back to } 0}} + \underbrace{C_d P_t^2}_{\substack{\text{correct power improves} \\ \text{as } d \text{ increases}}}$$

since siblings evolve independently; correct power improves as d increases

even if $\delta_t(z) = \delta_t(w) = 1$,
still require ≥ 2 sites
 $y \in T_t(x)$ w/ $\delta_t(y) = 1$

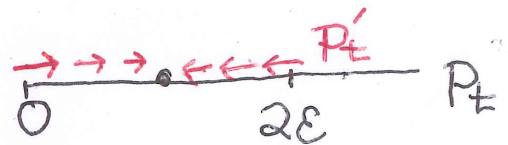


Consequently,

$$(5) \quad P_t' \leq \varepsilon - P_t (1 - 2C_d P_t).$$

Since $P_0 = 0$, for small ε ,

$$(6) \quad P_t \leq 2\varepsilon \quad \text{for all } t.$$



Since all sites in $T_n(x)$ evolve independently,
(2') follows from (6) and SLLN.

Basics behind proof of Theorem 2

Proof of Theorem 2 is more complicated
than that of Theorem 1.

Some ideas:

The limit

$$\mathbb{E}_t^J \xrightarrow{d} \mathbb{E}_\infty^J \text{ as } t \rightarrow \infty$$

is equivalent to

$$(7) \quad \mathbb{E}_0^{J, -s} \xrightarrow{d} \mathbb{E}_\infty^J \text{ as } s \rightarrow \infty.$$

process starts at
time $-s$ instead of 0

will show a.s.
convergence in this
setting

Employ δ_t^{-s} . ← analog of δ_t , but starting
at time $-s$

Since δ_t^{-s} is increasing in s

$\lim_{s \rightarrow \infty} \delta_t^{-s}$ exists. ← want to also show
for \mathbb{E}_0^{-s}

Trace paths of possible influence backwards in time.

← refer to as search histories

Is complicated. ← no duality

But can control these using δ_{\pm}^{-s} .

Corresponding to these paths are supercritical branching processes, at least one of which must eventually die out, but not die out quickly if $E_0^{I_1 - S_1} \neq E_0^{I_1 - S_2}(x)$ for large $S_1 \leq S_2$.

unlikely event

deaths are associated w/ noise pts;
(births w/ MV pts)

