

1 **A MATHEMATICAL MODEL OF ASYNCHRONOUS DATA FLOW**
2 **IN PARALLEL COMPUTERS** *

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4 **Abstract.** We present a simplified model of data flow on processors in a high performance
5 computing framework involving computations necessitating inter-processor communications. From
6 this ordinary differential model, we take its asymptotic limit, resulting in a model which treats the
7 computer as a continuum of processors and data flow as an Eulerian fluid governed by a conservation
8 law. We derive a Hamilton-Jacobi equation associated with this conservation law for which the
9 existence and uniqueness of solutions can be proven. We then present the results of numerical
10 experiments for both discrete and continuum models; these show a qualitative agreement between
11 the two and the effect of variations in the computing environment's processing capabilities on the
12 progress of the modeled computation.

13 **Key words.** data flow, high-performance computing, asymptotic approximation, conservation
14 laws, Hamilton-Jacobi equation

15 **AMS subject classifications.** 35L65, 93A30, 70H20, 41A60

16 **1. Introduction.** It has been well-established that current and future genera-
17 tions of extreme scale computers have achieved and, for the foreseeable future, are
18 expected to achieve increases in performance via greater levels of parallelism at multi-
19 ple levels — e.g., within the processors as well as increasing the number of processors
20 and nodes — as opposed to increases in clock speeds, which are expected to remain
21 relatively flat. Additionally, extremely concurrent codes, involving dynamic parallel-
22 ism and greater degrees of asynchronous parallel executions, are increasingly needed
23 to leverage this large scale parallelism [15, 29].

24 As machine improvements depend on increasingly complex architectures and as
25 additional constraints on system development and planning (such as power consump-
26 tion [15]) arise, a need for predictive, quantitative models of computational perfor-
27 mance will grow greater. Previously developed modeling tools such as LogP [12, 13]
28 result in easily evaluated models which can prove difficult to extend and modify. Al-
29 ternatively, PRAM models have been used as abstractions of codes; these however
30 have scalability issues due to the complexity of simulating them [25]. Other modern
31 tools [23, 24] are similarly still limited to fine-grain simulations of at most a few dozen
32 nodes, again due to their computational complexity during simulations.

33 Core counts are now in the hundreds of thousands and millions on machines in
34 the TOP500 list of supercomputers; node counts consistently are in the thousands [1].
35 Such numbers mean that fine-grained simulation tools (such as those listed above) are

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Funding: This research is sponsored by the Office of Advanced Scientific Computing Research;
U.S. Department of Energy. The work was performed at the Oak Ridge National Laboratory, which
is managed by UT-Battelle, LLC under Contract No. De-AC05-00OR22725.

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36 incapable of describing large-scale phenomena; essentially the simulation tools begin
 37 to require computational resources beyond those of the systems they are simulating.
 38 Alternative approaches have been proposed to address these issues: miniapp codes
 39 can mimic key features of the performance of exascale codes with a much smaller
 40 codebase [16]. Aspen, a framework for performance modeling [28, 31], uses a domain
 41 specific language which encodes both abstracted features of machines hardware and
 42 specific software applications to provide coarse-grained simulations. However, these
 43 suffer from the need to develop specialized simulation codes which can be problem
 44 dependent, resulting in possibly labor-intensive tools. A workflow modeling approach,
 45 Pegasus, has been developed to model workflows using a graph-theoretic perspective
 46 to detect and manage anomalies in the computing environment [14].

47 We propose developing a macroscopic model of extreme scale computers which
 48 views such computing environments in a continuum framework. Such a model has
 49 several potential benefits: in addition to being computationally tractable, it will open
 50 up the possibility of using the theoretical tools of partial differential equations to
 51 understand and control the performance of high-performance computing systems.
 52 Specifically, our goal is to derive a fluid-limit model of data flow — which can be
 53 described by a partial differential equation — from a simplified deterministic model
 54 of data processing and flow in an extreme scale computer with interprocessor com-
 55 munications and asynchronous executions. Fluid models, beyond their obvious utility
 56 in physical systems, have been used to model flows in networks, such as vehicular
 57 traffic flows [3], supply chains [2], and gas networks [4, 8]. In particular, as discussed
 58 in [22] and [2], such fluid models lie at the end of a hierarchy of models which begin
 59 with microscopic or discrete models. That is, similar to the derivation of physical
 60 fluid laws from many body physics, one may derive continuum-level flow equations
 61 from discrete-level models of the dynamics of agent interactions. With such a model,
 62 standard numerical simulation tools and analytical methods may be brought to bear
 63 for studying large-scale phenomena in extreme-scale computing.

64 We begin in Section 2 with a microscopic model of a network of processors per-
 65 forming a multi-stage computational task which necessitates inter-processor commu-
 66 nications. In Section 3, we derive a formal asymptotic limit of this agent-based model
 67 as the scale of the system increases, resulting in an Eulerian fluid flow model. Along
 68 with the resulting nonlinear conservation law, we present a related Hamilton-Jacobi
 69 equation and establish the existence of solutions in Section 3. In Section 4, we pre-
 70 sent the results of numerical experiments to show agreement between the microscopic
 71 and fluid models and then illustrate the behavior of solutions under heterogeneous
 72 computing layouts.

73 **2. The discrete model.** In this section, we introduce the microscopic model,
 74 which is based on a highly simplified, deterministic, semi-discrete ordinary differen-
 75 tial equation (ODE). We imagine the computer as a network of processors $\{\mathcal{P}_i\}_{i=1}^{i^{\max}}$
 76 that are arranged in a one-dimensional, periodic lattice. The computer is assigned a
 77 computational job involving a sequence of k^{\max} tasks which are identical in the sense
 78 that each one takes the same effort to complete. This computational job is divided
 79 by distributing data amongst processors. We denote by $q_{i,k}(t)$ the amount of data in
 80 \mathcal{P}_i that sits in stage k at time t .

81 **2.1. Conservation law.** The dynamics of $q_{i,k}$ are given by a conservation law
 82 of the form

83 (2.1)
$$\dot{q}_{i,k}(t) = F_{i,k-1}(t) - F_{i,k}(t), \quad k = 1, \dots, k^{\max}, \quad i = 1, \dots, i^{\max},$$

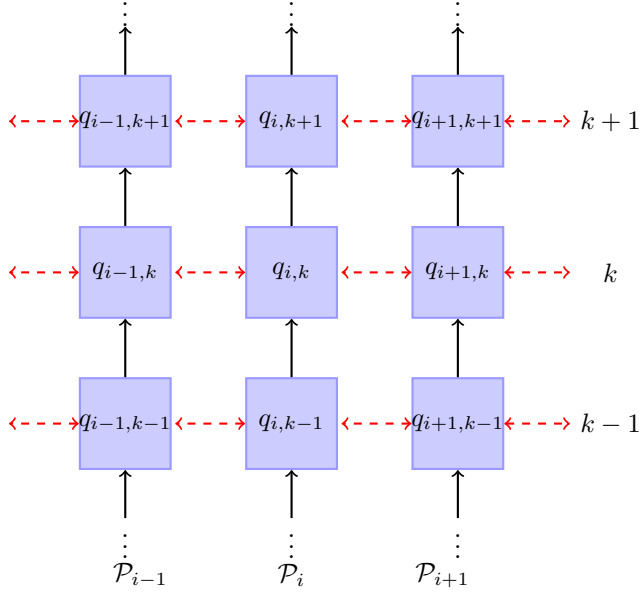


Fig. 1: Schematic of network of processors. Dashed lines denote inter-processor communications

84 where $F_{i,k}$ ($i = 1, \dots, i^{\max}$, $k = 1, \dots, k^{\max} - 1$) is the rate of data moving in processor
 85 i from stage k to $k + 1$, referred to as the *throughput*. At the first stage $k = 1$, $F_{i,0}$
 86 ($i = 1, \dots, i^{\max}$) is the rate of data being loaded into processor i to be processed,
 87 referred to as the *inflow*, and at the final stage, $F_{i,k^{\max}}$ ($i = 1, \dots, i^{\max}$) is the rate of
 88 data completing the final stage of the job, referred to as the *outflow*. Equation (2.1)
 89 implies that the data in each processor is neither created or destroyed, only moved in
 90 and out of the processor or in between stages; that is,

$$91 \quad (2.2) \quad \frac{d}{dt} \left(\sum_{k=1}^{k^{\max}} q_{i,k} \right) = F_{i,0} - F_{i,k^{\max}}.$$

92 A key aspect of the model is that it does not separately track data that moves between
 93 processors; instead the effects of communication delays will be incorporated directly
 94 into the definition of the throughputs.

95 A fundamental quantity of interest in the discrete model is $Q_{i,k}(t)$, which is
 96 defined as the amount of data at time t that has gone through the first $k - 1$ stages
 97 of \mathcal{P}_i . For each $t \geq 0$,

$$98 \quad (2.3) \quad Q_{i,k}(t) = \left(\sum_{j=k}^{k^{\max}} q_{i,j}(t) \right) + \int_0^t F_{i,k^{\max}}(s) ds.$$

99 **2.2. Processor throttling.** We now turn to specifying the form of $F_{i,k}$. In the
 100 absence of throttling, each processor \mathcal{P}_i moves data between stages at a rate $a_i \geq 0$,

101 which we refer to as the *maximum throughput*. For the purposes of the current paper,
 102 we assume that a_i is given. In practice, it must be determined from experiments,
 103 fine-scale models, or a combination of both. It may also depend on k , although for
 104 simplicity, we assume here that it does not. Throttling is said to occur whenever
 105 $F_{i,k}(t) < a_i$; this happens for one of two reasons.

- 106 1. **Self-throttling:** Given an amount of data $q_{i,k}$ to be processed at stage k in
 107 processor \mathcal{P}_i , we define the self-throttling function

$$108 \quad (2.4) \quad v_1(q_{i,k}; q_*) = \max \left\{ 0, \min \left\{ 1, \frac{q_{i,k}}{q_*} \right\} \right\}.$$

109 Clearly if no data is available to be processed, then $F_{i,k} = 0$. Furthermore,
 110 we assume that if the amount of data to be processed drops below a certain
 111 threshold $q_* > 0$, then \mathcal{P}_i cannot maintain the throughput a_i and the
 112 throughput is reduced.

- 113 2. **Neighbor throttling:** As the computational task is not entirely parallel
 114 across processors, \mathcal{P}_i requires sufficient information from its neighbors to
 115 perform task k at full throughput. The neighbor throttling function v_2 models
 116 this dependence. It gives the amount of available data on \mathcal{P}_i at stage k

$$117 \quad (2.5) \quad v_2(q_{i,k}, \Delta_{i+1,k}, \Delta_{i-1,k}; \beta) = \min \left\{ q_{i,k}, \frac{1}{\beta} \max\{\Delta_{i+1,k}, 0\}, \frac{1}{\beta} \max\{\Delta_{i-1,k}, 0\} \right\}.$$

118 Here $\Delta_{i\pm 1,k}$ denotes the data on the right/left neighbor which is available
 119 to be used by \mathcal{P}_i to process $q_{i,k}$. The parameter $\beta \in (0, 1]$ allows for the
 120 possibility that computations do not rely in a one-to-one fashion upon the
 121 availability of data from neighbors. If $\Delta_{i\pm 1,k} = 0$ the processing of data stops
 122 due to the absence of a necessary component of the computational task and
 123 so $F_{i,k} = 0$. Alternatively, if both $\Delta_{i+1,k}$ and $\Delta_{i-1,k}$ exceed $\beta q_{i,k}$, then \mathcal{P}_i
 124 has sufficient data from its neighbors to process $q_{i,k}$ and no throttling occurs.
 125 The data from the left/right neighbor which is available for processing at
 126 stage k is given by

$$127 \quad (2.6) \quad \Delta_{i\pm 1,k} = Q_{i\pm 1,k} - Q_{i,k+1} = Q_{i\pm 1,k} - (Q_{i,k} - q_{i,k}).$$

128 The data on each neighbor must have completed the same stage for it to be
 129 available; additionally, this data is not reused on \mathcal{P}_i for the same stage. This
 130 means that the data available to be used from the neighbors can be written
 131 as above and so the amount of data available to be processed on \mathcal{P}_i at stage
 132 k is given by

$$133 \quad (2.7) \quad v_2(q_{i,k}, Q_{i+1,k} - Q_{i,k} + q_{i,k}, Q_{i-1,k} - Q_{i,k} + q_{i,k}; \beta).$$

134 The throughput $F_{i,k}$ is a composition of the throttling functions v_1 and v_2 :

$$135 \quad (2.8) \quad F_{i,k} = a_i v_1 \left(v_2(q_{i,k}, Q_{i+1,k} - Q_{i,k} + q_{i,k}, Q_{i-1,k} - Q_{i,k} + q_{i,k}; \beta); q_* \right).$$

136 At first glance, this definition of $F_{i,k}$ appears circular since it depends on $Q_{i,k}$, which
 137 in turn depends on $F_{i,k}^{\max}$. However, as a consequence of the conservation law (2.2),

$$138 \quad (2.9) \quad \int_0^t F_{i,k}^{\max}(s) ds = \int_0^t F_{i,0}(s) ds + \sum_{j=1}^{k^{\max}} q_{i,j}(0) - \sum_{j=1}^{k^{\max}} q_{i,j}(t)$$

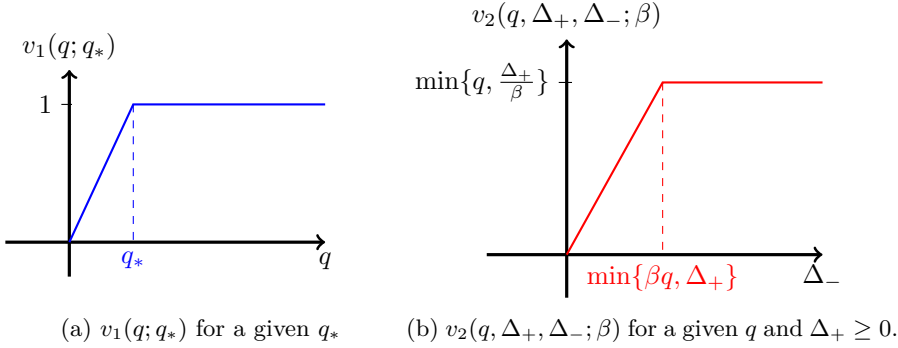


Fig. 2: Throttling functions v_1 and v_2 , defined in (2.4) and (2.5), respectively

139 Thus to complete the model, we need only prescribe initial data $q_{i,k}(0)$ and the inflow
 140 $F_{i,0}$. To prescribe the inflow, we specify $q_{i,0}$ and then let $F_{i,0}$ be evaluated according
 141 to (2.8).

142 **PROPOSITION 2.1.** *The system (2.1) with (i) throughput $F_{i,k}$ defined in (2.8)*
 143 *for $i = 1, \dots, i^{\max}$ and $k = 1, \dots, k^{\max}$; (ii) prescribed initial data $q_{i,k}(0)$ for $i =$*
 144 *$1, \dots, i^{\max}$ and $k = 1, \dots, k^{\max}$; and (iii) prescribed inflow data $F_{i,0}$ for $i = 1, \dots, i^{\max}$*
 145 *and $t \geq 0$ has a unique solution for all $t \geq 0$. Moreover, if $q_{i,k}(0) \geq 0$ for all*
 146 *$i = 1, \dots, i^{\max}$ and $k = 1, \dots, k^{\max}$, then $q_{i,k}(t) \geq 0$ for all $t \geq 0$ and $i = 1, \dots, i^{\max}$*
 147 *and $k = 1, \dots, k^{\max}$.*

148 *Proof.* Since $F_{i,k}$ is globally Lipschitz in its arguments for every i, k , standard
 149 ODE theory (see, for example Theorem III.VI of [32]) implies the existence of a
 150 unique solution. Moreover, it is clear from (2.8) that $0 \leq F_{i,k} \leq a_i v_1(q_{i,k}; q_*)$. Hence
 151 according to (2.1),

$$152 \quad (2.10) \quad \dot{q}_{i,k}(t) \geq -a_i v_1(q_{i,k}; q_*).$$

153 Standard comparison results for ordinary differential equations imply then that $q(t) \geq$
 154 0 . (See, for example, Lemma 1.2 of [30] and for comparison, use the zero function,
 155 which satisfies (2.10) as an equality.) \square

156 **3. The continuum model.** In this section, we derive a continuous model that
 157 is formally accurate in the limit as i^{\max} and k^{\max} tend to infinity. We assume, in
 158 taking this limit, that the job performed by the computer is fixed – that is, the total
 159 amount of work does not change. For given i^{\max}, k^{\max} , we define the quantities:

$$160 \quad (3.1) \quad \delta := (k^{\max})^{-1}, \quad \varepsilon := (i^{\max})^{-1}, \quad \eta := \frac{\varepsilon}{\delta}.$$

161 Here, δ is the fraction of the work done in each stage and ε is the average amount of
 162 data in a processor. Finally, η is simply the ratio of k^{\max} and i^{\max} which will be of
 163 use in the following analysis.

164 **3.1. Formal Derivation.** To derive a continuum model, we first express the
 165 ODE (2.1) in terms of the following $O(1)$ quantities:

(3.2)

$$166 \quad r_* := \frac{q_*}{\varepsilon\delta}, \quad r_{i,k} := \frac{q_{i,k}}{\varepsilon\delta}, \quad R_{i,k} := \frac{1}{\varepsilon}Q_{i,k}, \quad D_{i,k}^\pm := \pm \frac{R_{i\pm 1,k} - R_{i,k}}{\varepsilon}, \quad \alpha_i := \frac{a_i}{\varepsilon}.$$

168 In terms of these rescaled quantities, the neighbor throttling function as can be written
 169 as

(3.3)

$$170 \quad v_2(q_{i,k}, Q_{i+1,k} - Q_{i,k} + q_{i,k}, Q_{i-1,k} - Q_{i,k} + q_{i,k}; \beta)$$

$$171 \quad = \min \left\{ \varepsilon\delta r_{i,k}, \frac{1}{\beta} \max\{\varepsilon R_{i+1,k} - \varepsilon R_{i,k} + \varepsilon\delta r_{i,k}, 0\}, \frac{1}{\beta} \max\{\varepsilon R_{i-1,k} - \varepsilon R_{i,k} + \varepsilon\delta r_{i,k}, 0\} \right\}$$

$$172 \quad = \varepsilon \min \left\{ \delta r_{i,k}, \frac{1}{\beta} \max\{R_{i+1,k} - R_{i,k} + \delta r_{i,k}, 0\}, \frac{1}{\beta} \max\{R_{i-1,k} - R_{i,k} + \delta r_{i,k}, 0\} \right\}$$

$$173 \quad = \varepsilon\delta \min \left\{ r_{i,k}, \frac{1}{\beta} \max\{\eta D_{i,k}^+ + r_{i,k}, 0\}, \frac{1}{\beta} \max\{-\eta D_{i,k}^- + r_{i,k}, 0\} \right\}. \quad \blacksquare$$

175 Therefore,

176

$$177 \quad (3.4) \quad v_1 \left(v_2(q_{i,k}, v_2(q, Q_{i+1,k} - Q_{i,k} + q_{i,k}, Q_{i-1,k} - Q_{i,k} + q_{i,k})); q_* \right)$$

$$178 \quad = \min \left\{ 1, \frac{\min \left\{ r_{i,k}, \frac{1}{\beta} \max\{\eta D_{i,k}^+ + r_{i,k}, 0\}, \frac{1}{\beta} \max\{-\eta D_{i,k}^- + r_{i,k}, 0\} \right\}}{r_*} \right\}.$$

180 With (3.4) in mind, we define the rescaled throttling functions

$$181 \quad (3.5a) \quad w_1(r, r_*) = \max \left\{ 0, \min \left\{ 1, \frac{r}{r_*} \right\} \right\}.$$

$$182 \quad (3.5b) \quad w_2(r, D^-, D^+; \eta, \beta) = \min \left\{ r, \frac{1}{\beta} \max\{\eta D^+ + r, 0\}, \frac{1}{\beta} \max\{\eta D^- + r, 0\} \right\}$$

184 and the composite function

$$185 \quad (3.6) \quad w(r, D^-, D^+; r_*, \alpha, \eta, \beta) := \alpha w_1(w_2(r, D^-, D^+; \eta, \beta); r_*).$$

186 The dynamics in (2.8) can now be re-expressed in terms of the $O(1)$ quantities in
 187 (3.2), thereby obtaining a evolution formula for $r_{i,k}$:

$$188 \quad (3.7) \quad \dot{r}_{i,k}(t) = \frac{f_{i,k-1}(t) - f_{i,k}(t)}{\delta}$$

189 for $i = 1, \dots, i^{\max}$ and $k = 1, \dots, k^{\max}$, where

$$190 \quad (3.8) \quad f_{i,k}(t) = w \left(r_{i,k}(t), -D_{i,k}^-(t), D_{i,k}^+(t); r_*, \alpha_i, \eta, \beta \right)$$

191 for $i = 1, \dots, i^{\max}$, $k = 0, \dots, k^{\max}$, and $r_{i,0}$ is prescribed for $i = 1, \dots, i^{\max}$

192 The next step is to interpret (3.7) as a conservative finite-difference formula for a
 193 sufficiently smooth function $\rho = \rho(x, y, t)$, defined on $[0, 1] \times [0, 1] \times [0, \infty)$, such that

$$194 \quad (3.9) \quad \rho(x_i, z_k, t) = r_{i,k}(t),$$

195 on grid points

$$196 \quad (3.10) \quad x_i = (i - 0.5)\varepsilon \quad \text{and} \quad z_k = k\delta,$$

197 for $i = 1, \dots, i^{\max}$ and $k = 0, \dots, k^{\max}$. We also let $\alpha = \alpha(x)$ be a continuous function
 198 such that $\alpha(x_i) = \alpha_i$.

199 Let $\psi = \psi(x, z, t)$ be a smooth test function with compact support on $[0, 1) \times$
 200 $[0, 1] \times [0, \infty)$ and set $\psi_{i,k}(t) = \psi(x_i, z_k, t)$. From (3.7),

$$201 \quad (3.11) \quad \delta \sum_{k=1}^{k^{\max}} \psi_{i,k}(t) \dot{r}_{i,k}(t) = \sum_{k=0}^{k^{\max}-1} [\psi_{i,k+1}(t) - \psi_{i,k}(t)] f_{i,k}(t)$$

$$202 \quad (3.12) \quad + \psi_{i,0}(t) f_{i,0}(t) - \psi_{i,k^{\max}}(t) f_{i,k^{\max}}(t).$$

204 Let the function $\phi = \phi(x, z, t)$ interpolate the fluxes on the grid:

$$205 \quad (3.13) \quad \phi(x_i, z_k, t) = f_{i,k}(t),$$

206 for $i = 1, \dots, i^{\max}$, $k = 1, \dots, k^{\max}$, $t \geq 0$. Then (3.11) can be interpreted formally as
 207 the weak formulation (with respect to z) of a conservation law for ρ with flux ϕ :

$$208 \quad \int_0^1 \psi(x, \xi, t) \partial_t \rho(x, \xi, t) d\xi = \int_0^1 \partial_z \psi(x, \xi, t) \phi(x, \xi, t) d\xi$$

$$209 \quad (3.14) \quad + \psi(x, 0, t) \phi(x, 0, t) - \psi(x, 1, t) \phi(x, 1, t) + O(\delta).$$

211 To derive a closed model from (3.14), we approximate ϕ in terms of ρ . Such an
 212 approximation depends on $D_{i,k}^{\pm}$ via the formula for $f_{i,k}$ in (3.8). From the definition
 213 of $Q_{i,k}$ in (2.3) and the scalings in (3.2), it follows that

$$214 \quad (3.15) \quad R_{i,k}(t) = \delta \sum_{j=k}^{k^{\max}} r_{i,j}(t) + \int_0^t f_{i,k^{\max}}(s) ds$$

215 and, moreover, that for any finite $t > 0$,

$$216 \quad (3.16) \quad \pm D_{i,k}^{\pm}(t) = \frac{\delta}{\varepsilon} \sum_{j=k}^{k^{\max}} [r_{i\pm 1,j} - r_{i,j}(t)] + \frac{1}{\varepsilon} \int_0^t [f_{i\pm 1,k^{\max}}(s) - f_{i,k^{\max}}(s)] ds$$

$$217 \quad (3.17) \quad = \delta \sum_{j=k}^{k^{\max}} \left[\pm \partial_x \rho(x_i, z_j, t) + \frac{\varepsilon}{2} \partial_x^2 \rho(x_i, z_j, t) + O(\varepsilon^2) \right]$$

$$218 \quad + \int_0^t \left[\pm \partial_x \phi(x_i, 1, s) + \frac{\varepsilon}{2} \partial_x^2 \phi(x_i, 1, s) + O(\varepsilon^2) \right] ds$$

$$219 \quad (3.18) \quad = \pm \partial_x P(x_i, z_k, t) + \frac{\varepsilon}{2} \partial_x^2 P(x_i, z_k, t) + O(\varepsilon^2) + O(\delta)$$

221 where P is given by

$$222 \quad (3.19) \quad P(x, z, t) = \int_z^1 \rho(x, \xi, t) d\xi + \int_0^t \phi(x, 1, s) ds$$

224 Motivated by the above calculation, we approximate ϕ by one of two flux functions:

225 (3.20a) $\Phi^{(0)}(\rho, \partial_x P; r_*, \alpha, \eta, \beta) = w(\rho, -\partial_x P, \partial_x P; r_*, \alpha, \eta, \beta)$

(3.20b)

226 $\Phi^{(1)}(\rho, \partial_x P, \partial_x^2 P; r_*, \alpha, \eta, \beta) = w\left(\rho, -\partial_x P + \frac{\varepsilon}{2}\partial_x^2 P, \partial_x P + \frac{\varepsilon}{2}\partial_x^2 P; r_*, \alpha, \eta, \beta\right).$

228 Using (3.18) and the Lipschitz continuity of w with respect to D^\pm , we conclude that

229 (3.21) $w(r_{i,k}(t), D_{i,k}^-(t), D_{i,k}^+(t); r_*, \alpha_i, \eta, \beta)$

230 $= \Phi^{(1)}(\rho(x_i, z_k, t), \partial_x P(x_i, z_k, t), \partial_x^2 P(x_i, z_k, t); r_*, a, \eta, \beta) + O(\varepsilon^2) + O(\delta).$

231 $= \Phi^{(0)}(\rho(x_i, z_k, t), \partial_x P(x_i, z_k, t); r_*, a, \eta, \beta) + O(\varepsilon) + O(\delta).$

233 Thus for $0 \ll \varepsilon, \delta \ll 1$, with $\eta \in (0, \infty)$ fixed, (3.14) is formally consistent with the
234 continuum model

235 (3.22a) $\partial_t \rho + \partial_z \Phi^{(\ell)}(\rho, \partial_x P, \partial_x^2 P; r_*, a, \eta, \beta) = 0, \quad (x, z, t) \in \mathbb{T}^1 \times (0, 1) \times (0, \infty),$

236 (3.22b) $\rho(x, 0, t) = \rho_{bc}(x, t), \quad (x, t) \in \mathbb{T}^1 \times (0, \infty),$

237 (3.22c) $\rho(x, z, 0) = \rho_0(x, z), \quad (x, z) \in \mathbb{T}^1 \times (0, 1)$

239 where

240 (3.23a) $P(x, z, t) = \int_z^1 \rho(x, \xi, t) d\xi + \int_0^t \phi^{(\ell)}(x, 1, s) ds,$

241 (3.23b) $\phi^{(\ell)}(x, z, t) = \Phi^{(\ell)}(\rho(x, z, t), \partial_x P(x, z, t), \partial_x^2 P(x, z, t); r_*, \alpha, \eta, \beta),$

243 and $\Phi^{(\ell)}$, $\ell \in \{0, 1\}$, is given in (3.20). For the sake of compactness, we have slightly
244 abused notation in (3.22a), as the definition of $\Phi^{(0)}$ is independent of $\partial_x^2 P$. Addition-
245 ally, we have identified $[0, 1)$ with the one-dimensional torus \mathbb{T}^1 in order to reflect the
246 periodic layout of the processors.

247 As in the discrete case, it may appear that the model in (3.22) is circular due to
248 the definition of P in (3.23a). However, as with F in (2.9), $\Phi^{(\ell)}$ can be unwrapped,
249 this time using the conservation law (3.22a); that is

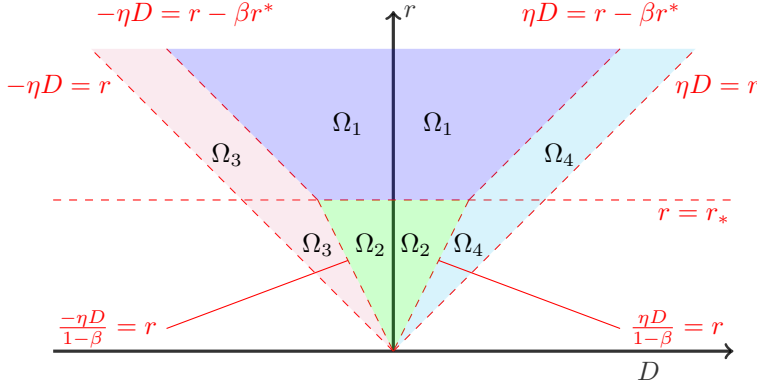
250 (3.24) $\int_0^t \phi^{(\ell)}(x, 1, s) ds = \int_0^t \phi(x, 0, s) ds + \int_0^1 \rho_0(x, \xi) d\xi - \int_0^1 \rho(x, \xi, t) d\xi$

251 Thus the continuum model is complete once initial condition ρ_0 and inflow condition
252 $\phi_{bc} := \phi(\cdot, 0, \cdot)$ are specified. In practice, ρ_{bc} is prescribed and then ϕ_{bc} is evaluated
253 using (3.23b) and (3.20).

254 We use the flux function $\Phi^{(0)}$ for all of the numerical simulations in Section 4.
255 This function is a piecewise constant that can be expressed in the following form:

256 (3.25) $\Phi^{(0)}(r, D; r_*, \alpha, \beta) = \begin{cases} \alpha & (r, D) \in \Omega_1, \\ \frac{\alpha r}{r_*} & (r, D) \in \Omega_2, \\ \frac{\alpha(r + \eta D)}{\beta r_*} & (r, D) \in \Omega_3, \\ \frac{\alpha(r - \eta D)}{\beta r_*} & (r, D) \in \Omega_4, \end{cases}$

257 where the subdomains Ω_i are depicted in Figure 3.


 Fig. 3: Flux $\Phi^{(0)}$ defined in (3.25)

258 **3.2. A Hamilton-Jacobi formulation.** To our knowledge, there are no immediate
 259 conclusions available from the literature regarding the existence or uniqueness
 260 of solutions to (3.22). However, we may consider instead a related Hamilton-Jacobi
 261 equation which in turn opens up the extensive theory of viscosity solutions. For back-
 262 ground, we refer to [6, 11, 20]. We are interested primarily in obtaining, for regular
 263 inputs α , ϕ_{bc} , and ρ_0 , the existence and uniqueness of P . It is possible that for more
 264 general inputs such results are available in the extensive viscosity solution literature
 265 (e.g. [5, 10, 21]).

266 Integrating (3.22a) with respect to z gives

$$267 \quad (3.26) \quad \partial_t \int_z^1 \rho(x, \xi, t) d\xi + \Phi^{(\ell)}(x, 1, t) - \Phi^{(\ell)}(x, z, t) = 0$$

268 Meanwhile, differentiating (3.23a) gives

$$269 \quad (3.27) \quad \partial_t P(x, z, t) = \partial_t \int_z^1 \rho(x, \xi, t) d\xi + \Phi^{(\ell)}(x, 1, t)$$

270 Combining (3.26) and (3.27) and using the fact that $\rho = -\partial_z P$ gives a closed
 271 Hamilton-Jacobi equation for P with initial and boundary conditions that are derived
 272 by applying the definition of P in (3.23a) to (3.22c) and (3.22b), respectively.
 273 The complete model is, for some $T > 0$,

$$274 \quad (3.28a) \quad \partial_t P - \Phi^{(\ell)}(-\partial_z P, \partial_x P, \partial_x^2 P; r_*, \alpha, \eta, \beta) = 0, \quad (x, z, t) \in \mathbb{T}^1 \times (0, 1) \times (0, T),$$

$$275 \quad (3.28b) \quad P(x, 0, t) - \int_0^1 \rho_0(x, \xi) d\xi - \int_0^t \phi_{bc}(x, s) ds = 0, \quad (x, t) \in \mathbb{T}^1 \times (0, T),$$

$$276 \quad (3.28c) \quad P(x, z, 0) - \int_z^1 \rho_0(x, \xi) d\xi = 0, \quad (x, z) \in \mathbb{T}^1 \times (0, 1),$$

277 where (3.28b) is derived by integrating (3.22b) over $z \in (0, 1)$ and applying (3.24).

278 **THEOREM 3.1.** Assume that α and ρ_0 are (i) non-negative, (ii) uniformly Lip-
 279 schitz in their arguments, and (iii) periodic in x (that is, $\alpha(0) = \alpha(1)$ and $\rho_{bc}(0, t) =$
 280

281 $\rho_{bc}(1, t)$ for all $t \in [0, T]$. Further, assume that there is an M where $\int_0^T \phi_{bc}(x, s) ds \leq$
 282 M for all $x \in \mathbb{T}^1$. Then there exists a unique, continuous, viscosity solution (in the
 283 sense of [11]) to (3.28).

284 *Proof.* We show that Lemma A.6 applies by first modifying the domain in (3.28).
 285 We extend α , ρ_0 , ρ_{bc} , and ϕ_{bc} as functions of x from \mathbb{T}^1 to all of \mathbb{R} by tiling; for
 286 simplicity, in the remainder of the proof we still refer to these extensions by the same
 287 name. The assumption that $\alpha(0) = \alpha(1)$ means that the extended version of α is
 288 uniformly Lipschitz on \mathbb{R} . We then consider (3.28) defined on $\Omega := \mathbb{R} \times (0, 1) \times (0, T)$,
 289 and to more closely align with the results in the appendix, let

$$290 \quad (3.29) \quad H^{(\ell)}(x, z, t, P, \nabla P, \nabla^2 P) = -\Phi^{(\ell)}(-\partial_z P, \partial_x P, \partial_x^2 P; r_*, \alpha(x), \eta, \beta)$$

291 for $\ell \in \{0, 1\}$, where $\nabla = (\partial_x, \partial_z)$. By the hypothesis on α , both $H^{(0)}$ and $H^{(1)}$ are
 292 uniformly Lipschitz on all of $\Omega \times \mathbb{R} \times \mathbb{R}^2 \times \mathcal{S}^2$, where \mathcal{S}^n is the space of all $n \times n$
 293 symmetric matrices. Moreover, $H^{(\ell)}$ is nonnegative, bounded by α , and independent
 294 of the argument P . This means that it immediately satisfies Hypotheses 1 and 3-8 of
 295 Lemma A.6. Thus the only condition of Lemma A.6 left to be verified is Hypothesis
 296 2, which is the degenerate ellipticity condition on $H^{(\ell)}$. Verifying this condition can
 297 be done in a sequence of simple steps, starting with the definitions of w_1 and w_2 .

298 w_1 is non-decreasing WRT r and w_2 is non-decreasing WRT D^+, D^- (see (3.5))
 299 $\implies w$ is non-decreasing WRT D^+, D^- (see (3.6))
 300 $\implies \Phi^{(\ell)}$ is non-decreasing WRT $\partial_x^2 P$ (see (3.20))
 301 $\implies H^{(\ell)}$ is degenerate elliptic (see (3.29))

303 Finally, to invoke Theorem A.8, we must establish the existence of subsolutions
 304 and supersolutions as defined in (A.2) and (A.3), respectively. This is done by the
 305 usual construction found in, for instance, [20, Section 2.3.2.1]. Let

$$306 \quad (3.30) \quad P^\pm(x, z, t) := \int_z^1 \rho_0(x, \xi) d\xi \pm M \pm t \cdot \sup_x \alpha(x).$$

307 Clearly $\partial_t P^\pm = \pm \sup \alpha(x)$ and since $|\Phi_\ell| \leq |\alpha(x)|$, it follows that

$$308 \quad (3.31) \quad \pm [\partial_t P^\pm - \Phi^{(\ell)}(-\partial_z P^\pm, \partial_x P^\pm, \partial_x^2 P^\pm; r_*, \alpha, \eta, \beta)] \geq 0.$$

309 Thus P^\pm satisfy the interior conditions in (A.3a) and (A.2a), respectively. Next
 310 write (3.28b) and (3.28c) in the form $h(t, x, P, \nabla P) = 0$. Then it is straight-forward
 311 to verify that $\pm h(t, x, P^\pm, \nabla P^\pm) \geq 0$. Hence P^\pm satisfies the parabolic boundary
 312 conditions in (A.3a) and (A.2a), respectively. Therefore P^- is a subsolution and P^+
 313 is a supersolution for (3.28). This completes the proof. \square

314 *Remark 3.2.* In general, results regarding the regularity of solutions to (3.28)
 315 using $\Phi^{(0)}$ (no such results are known by the authors for $\Phi^{(1)}$) require additional
 316 smoothness of (and possibly convexification of) $\Phi^{(0)}$ as well as other technical condi-
 317 tions (see [26], [7], [9]). Therefore obtaining the existence of a L^1 function ρ solving
 318 (3.22) (in some generalized sense) via the existence of P solving (3.28) is still an open
 319 problem.

320 **3.3. Higher Dimensional Models.** Both the discrete and continuum models
 321 above can be readily extended to systems of processors arranged in an n -dimensional

322 periodic lattice. Assuming that processors only communicate with their nearest neigh-
 323 bors (i.e., no diagonal communication), the n -dimensional analog of the system formed
 324 by (2.1), (2.3), and (2.8) is:

$$325 \quad (3.32) \quad \frac{dq_{\mathbf{i},k}}{dt} = F_{\mathbf{i},k-1} - F_{\mathbf{i},k}, \quad Q_{\mathbf{i},k}(t) = \left(\sum_{j=k}^{k^{\max}} q_{\mathbf{i},j}(t) \right) + \int_0^t F_{\mathbf{i},k^{\max}}(s) ds.$$

(3.33)

$$326 \quad F_{\mathbf{i},k} = a_{\mathbf{i}} v_1 \left(\min_{1 \leq d \leq n} \left\{ v_2 (q_{\mathbf{i},k}, Q_{\mathbf{i}-\mathbf{e}_d,k} - Q_{\mathbf{i},k} + q_{\mathbf{i},k}, Q_{\mathbf{i}+\mathbf{e}_d,k} - Q_{\mathbf{i},k} + q_{\mathbf{i},k}; \beta) \right\}; q_* \right)$$

327

328 where $\mathbf{i} = (i_1, \dots, i_n)$ is a multi-index and $(\mathbf{e}_d)_i = \delta_{d,i}$. As in the one-dimensional case,
 329 v_2 provides the amount of available data to process, after accounting for the throttling
 330 from neighbors over a given axis. The multidimensional discrete model then takes the
 331 minimum over all possible axes in order to determine what is available to be processed
 332 to the next stage. As in the one-dimensional case, self-throttling is computed using v_1
 333 based on the amount of data available for processing.

334 If i_d^{\max} denotes the number of processors along the d direction, we let $\varepsilon_d =$
 335 $(i_d^{\max})^{-1}$; the definition of δ is unchanged. For notational convenience, we set $V :=$
 336 $\prod_{d=1}^n \varepsilon_d$. We define $w_{2,d}$ by replacing η in the definition of w_2 with $\eta_d := \varepsilon_d/\delta$. Then
 337 we define the quantities analogous to those in (3.2)

(3.34)

$$338 \quad r_* := \frac{q_*}{\delta V}, \quad \hat{\alpha}_{\mathbf{i}} := \frac{\alpha_{\mathbf{i}}}{V}, \quad r_{\mathbf{i},k} := \frac{q_{\mathbf{i},k}}{\delta V}, \quad R_{\mathbf{i},k} := \frac{1}{V} Q_{\mathbf{i},k}, \quad D_{\mathbf{i},k}^{d,\pm} := \pm \frac{R_{\mathbf{i} \pm \mathbf{e}_d,k} - R_{\mathbf{i},k}}{\varepsilon_d}.$$

339

340 Continuing as in Section 3.1, we define

$$341 \quad (3.35) \quad x_{\mathbf{i}} = \left(\varepsilon_1(i_1 + 0.5), \dots, \varepsilon_n(i_n + 0.5) \right), \quad z_k = k\delta$$

342 and the smooth density function $\rho(x, z, t)$ defined on $[0, 1]^n \times (0, 1) \times [0, \infty)$ for which
 343 $\rho(x_{\mathbf{i}}, z_k, t) = r_{\mathbf{i},k}(t)$. Arguments analogous to those used to obtain (3.16) and (3.18)
 344 give us that

$$345 \quad (3.36) \quad \pm D_{\mathbf{i},k}^{d,\pm} \approx \pm \partial_{x_d} P(x_{\mathbf{i}}, z_k, t) + \frac{\varepsilon_d}{2} \partial_{x_d}^2 P(x_{\mathbf{i}}, z_k, t).$$

346 We assume that all of the i_d^{\max} are of the same order, so that the order of accuracy
 347 of the above approximation is consistent across all dimensions.

348 As advection in the z -direction is unchanged, we have the continuum model

$$349 \quad (3.37a) \quad \partial_t \rho + \partial_z \Phi^{(\ell)}(\rho, \nabla_x P, \nabla_x^2 P; r_*, \alpha, \vec{\eta}, \beta) = 0, \quad (x, z, t) \in \mathbb{T}^n \times (0, 1) \times (0, \infty),$$

$$350 \quad (3.37b) \quad \rho(x, 0, t) = \rho_{bc}(x, t), \quad (x, t) \in \mathbb{T}^n \times (0, \infty),$$

$$351 \quad (3.37c) \quad \rho(x, z, 0) = \rho_0(x, z), \quad (x, z) \in \mathbb{T}^n \times (0, 1),$$

352

353 where \mathbb{T}^n denotes the n -dimensional torus parameterized by $[0, 1]^n$, P and $\phi^{(\ell)}$ are
 354 defined as in (3.23a) and (3.23b), respectively, and the form of $\Phi^{(\ell)}$, $\ell \in \{0, 1\}$ is a

355 slightly generalized version of (3.20).

(3.38a)

$$356 \quad \Phi^{(0)}(\rho, \nabla_x P, \nabla_x^2 P; r_*, \alpha, \vec{\eta}, \beta) = \alpha w_1 \left(\min_{1 \leq d \leq n} \left([w_2(\rho, \partial_{x_d} P, -\partial_{x_d} P; \eta_d, \beta)] \right); r_* \right)$$

(3.38b)

$$357 \quad \Phi^{(1)}(\rho, \nabla_x P, \nabla_x^2 P; r_*, \alpha, \vec{\eta}, \beta) =$$

(3.38c)

$$358 \quad \alpha w_1 \left(\min_{1 \leq d \leq n} \left([w_2(\rho, \partial_{x_d} P + \frac{\varepsilon_d}{2} \partial_{x_d}^2 P, -\partial_{x_d} P + \frac{\varepsilon_d}{2} \partial_{x_d}^2 P; \eta_d, \beta)] \right); r_* \right)$$

359

360 For notational convenience we have defined $\Phi^{(\ell)}$ using the full tensor $\nabla_x^2 P$; how-
 361 ever, we note that the flux functions do not depend on mixed second derivatives. The
 362 procedure used in Section 3.2 to obtain a Hamilton-Jacobi equation for P can be
 363 repeated here; the only changes are (i) the multi-dimensional version of $\Phi^{(\ell)}$ in (3.20)
 364 and (ii) the domain of the x variable. Verifying that these newly-defined flux functions
 365 satisfy the conditions of Theorem A.6 is essentially the same as before. Existence and
 366 uniqueness of viscosity solutions P then follow.

367 **4. Numerical Simulations.** In this section, we perform numerical simulations
 368 of the one dimensional processor system in order to (i) test the ability of the macro-
 369 scopic model to approximate the discrete model when ε and δ are small and (ii)
 370 explore how model parameters affect the model output. All simulations are based on
 371 the flux $\Phi^{(0)}$, although results with $\Phi^{(1)}$ demonstrate similar. Problem data is speci-
 372 fied in terms continuum model of continuum models quantities. These quantities are
 373 translated back to discrete model quantities in order to implement ODE simulations.

374 **4.1. ODE Implementation.** The explicit two-step Adams-Bashforth (Section
 375 III of [18]) is used to simulate the discrete model formed by (2.1), (2.3), and (2.8).
 376 Given η , values i^{\max} and k^{\max} are chosen so that $k^{\max}/i^{\max} = \eta$ (cf. (3.1)). We then
 377 compute a solution to the discrete model as follows. Using (3.2) and (3.9), we convert
 378 $r_*, a, \rho_0, \rho_{bc}$ to their discrete counterparts:

$$379 \quad (4.1) \quad q_* = \varepsilon \delta r_*, \quad q_{i,k}(0) = \varepsilon \delta \rho_0(x_i, z_k), \quad q_{i,0}(t) = \rho_{bc}(x, t), \quad a_i = \varepsilon \alpha(x_i).$$

380 This discrete model data is used to set the time step:

$$381 \quad (4.2) \quad \Delta t = \frac{q_*}{2(\max_i a_i) \sqrt{i^{\max} k^{\max}}}.$$

382 The outflow at $F_{i,k^{\max}}$ is tracked and accumulated over time in order to compute $Q_{i,k}$
 383 from (2.3). At the final time T , the result of the explicit time stepping is converted
 384 back, via the formula in (3.2), i.e., $r_{i,k}(T) = (\varepsilon \delta)^{-1} q_{i,k}(T)$. In order to compare this
 385 against solutions to the continuum model (see below), we use these point-wise values
 386 to generate a piecewise constant function r over the cells $C_{i,k} = (x_i - .5\varepsilon, x_i + .5\varepsilon) \times$
 387 $(z_k, z_k + \delta)$:

$$388 \quad (4.3) \quad r(x, z) = \sum_{i,k} r_{i,k} \chi_{C_{i,k}}(x, z).$$

389 **4.2. Hamilton-Jacobi Implementation.** The Hamilton Jacobi equation (3.28) ■
 390 is solved numerically using a fifth-order WENO interpolation in x and z and the op-
 391 timal third-order SSP Runge-Kutta method for time integration. Details of these

392 algorithms can be found in Sections 3.2 and 6, respectively, of [27]. Once a numerical
 393 solution for P is computed, we again use WENO interpolation to approximate ρ via
 394 the relation $\rho(x, z, t) = -\partial_z P(x, z, t)$.

395 To condense the notation, let $\sigma = \partial_x P$, $\tau = \partial_z P$ and $v = \partial_{xx} P$. Then for fixed
 396 r_* , α , η , β , and ℓ , let $H(\sigma, \tau, v) = -\Phi^{(\ell)}(-\tau, \sigma, v; r_*, a, \eta, \beta)$. The numerical solution
 397 for P is computed on a grid $\{x_n, z_m\}$ where

$$398 \quad (4.4) \quad x_n = n\Delta x, \quad n = 1, 2, \dots, N, \quad \Delta x = N^{-1},$$

$$399 \quad (4.5) \quad z_m = m\Delta z, \quad m = 1, 2, \dots, M, \quad \Delta z = M^{-1}.$$

401 The semi-discrete method for the grid function $P_{n,m}(t) \approx P(x_n, z_m, t)$ is

$$402 \quad (4.6) \quad \frac{d}{dt} P_{n,m}(t) = -\hat{H}(\sigma_{n,m}^-, \sigma_{n,m}^+, \tau_{n,m}^-, \tau_{n,m}^+; v_{n,m}),$$

403 where the numerical approximations $\sigma_{n,m}^\pm \approx \sigma(x_n^\pm, z_m)$ and $\tau_{n,m}^\pm \approx \tau(x_n, z_m^\pm)$ are
 404 obtained via WENO interpolation and $v_{n,m} \approx v(x_n, z_m)$ is computed by central
 405 difference. The numerical flux function \hat{H} , based on the global Lax-Friedrichs flux:

$$406 \quad (4.7) \quad \hat{H}(\sigma^-, \sigma^+, \tau^-, \tau^+; v) = H\left(\frac{\sigma^- + \sigma^+}{2}, \frac{\tau^- + \tau^+}{2}, v\right) - \frac{1}{2}\lambda^x(\sigma^+ - \sigma^-) - \frac{1}{2}\lambda^z(\tau^+ - \tau^-),$$

407 where

$$408 \quad (4.8) \quad \lambda^x = \max_{\sigma, \tau} |H_\sigma| = \frac{\alpha\eta}{\beta r_*}, \quad \lambda^z = \max_{\sigma, \tau} |H_\tau| = \frac{\alpha}{\beta r_*}.$$

409 The time step for the SSP integrator is given by

$$410 \quad (4.9) \quad \Delta t \left(\frac{\lambda^x}{\Delta x} + \frac{\lambda^z}{\Delta z} \right) \leq 0.6.$$

411 **4.3. Experiments.** We perform a sequence of exploratory experiments below,
 412 modifying the parameters η and β , as well as the throughput function α . In all cases,
 413 α , ρ_0 , and ρ_{bc} are periodic with respect to x and the parameter $r_* = 1$. Results are
 414 presented as two-dimensional color maps or line-outs in the z direction. In all figures,
 415 the horizontal axis corresponds to the z -axis. Profiles of α for each experiment are
 416 depicted in Figure 4.

417 **EXAMPLE 1** (Agreement between models). *The purpose of this example is to*
 418 *demonstrate that the macroscopic model approximates the microscopic model when ε*
 419 *and δ are sufficiently small. We set $\beta = 1$ and consider $\eta \in \{0.2, 1, 5\}$. The initial*
 420 *condition, boundary condition, and processor speed are given by*

$$421 \quad (4.10) \quad \rho_0(x, z) = 1.5(\sin(2\pi z))^6 \chi_{[0, 0.5]}(z), \quad \rho_{bc}(x, t) = 0, \quad \alpha(x) = 1 - 0.4(\sin(\pi x))^2,$$

422 *respectively. Both models are simulated up to a final time $t = 0.5$.*

423 *For this example, the Hamilton-Jacobi simulation is performed with a 1000×1000*
 424 *mesh and a time step chosen according to (4.9) in order to generate a highly resolved*
 425 *numerical solution of the macroscopic model. For the microscopic model, we use*
 426 *$i^{\max} = 1000$ and $k^{\max} = 200$ when $\eta = 0.2$, $i^{\max} = k^{\max} = 500$ when $\eta = 1$, and*
 427 *$i^{\max} = 200$ and $k^{\max} = 1000$ when $\eta = 5$. These solutions to the microscopic model*
 428 *are then used to obtain the piecewise-constant function r on the 1000×1000 mesh*
 429 *from the Hamilton-Jacobi simulation.*

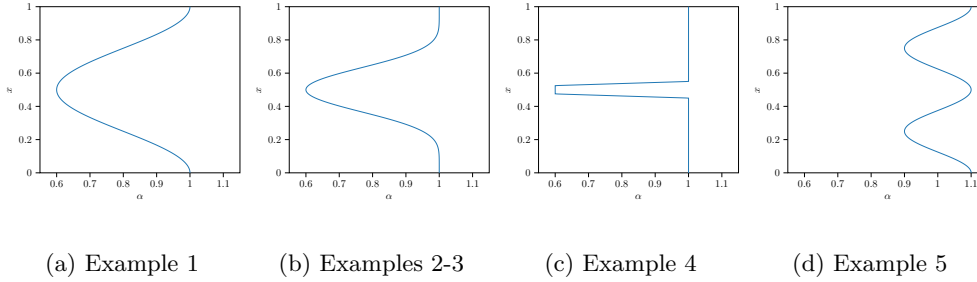


Fig. 4: Profiles of the processor speed α used in the numerical experiments. The non-standard orientation of the graphs is set to match the axes in the numerical results that follow.

430 Numerical results for $\eta = 0.2$, $\eta = 1.0$, and $\eta = 5.0$ are shown in [Figure 5](#),
 431 [Figure 6](#), and [Figure 7](#), respectively. While the results demonstrate general qualitative
 432 agreement between the models, discrepancies develop over time, especially for smaller
 433 values of η ; see [Figures 5i](#) and [5l](#). For the worst case scenario ($\eta = 0.2$), we increase
 434 the size of the discrete model by a factor of 2.5 (giving $i^{\max} = 2500$ and $k^{\max} = 500$),
 435 at which point the discrepancy between models decreases noticeably; see [Figure 8](#).

436 For the remaining examples, the Hamilton-Jacobi simulations are performed on
 437 a coarser mesh of 100×100 .

438 **EXAMPLE 2** (Variations in η). *In this example, we examine the effect of η on*
 439 *solutions to the macroscopic model while $\beta = 1.0$ is fixed. The initial condition,*
 440 *boundary condition, and processor speed are given by*

$$441 \quad (4.11) \quad \rho_0(x, z) = 1.5\chi_{z \leq 0.2}(x, z) \quad \rho_{bc}(x, t) = 0, \quad \alpha(x) = 1 - 0.4(\sin(\pi x))^6,$$

442 *respectively. It is expected that the slower processor speed around $x = 0.5$ will slow*
 443 *down neighboring processors due to neighbor-based throttling, encoded in the definition*
 444 *of w_2 in (3.5b). Moreover, the effect should become more global in x as η increases,*
 445 *since larger values of η correspond to a larger number of stages per processor. Indeed*
 446 *as the stages increase, interactions between neighbors begin to have a cumulative global*
 447 *effect. This trend can be observed by comparing results across the first three rows of*
 448 [Figure 9](#) *and in the line-outs in the final row.*

449 **EXAMPLE 3** (Variations in β). *In this example, we examine the effect of β on*
 450 *solutions to the macroscopic model, while holding $\eta = 1.0$ fixed. The initial condition,*
 451 *boundary condition, and processor speed are again given by (4.11).*

452 *Based on the definition of the function w_2 in (3.5b), the expectation is that smaller*
 453 *values of β will lead to reduced throttling effects. Such behavior is confirmed by the*
 454 *numerical results in [Figure 10](#).*

455 **EXAMPLE 4** (Highly localized slowdown). *In this example, we explore the effects*
 456 *of a highly localized slowdown in processor speed when $\eta = \beta = 1$. The initial*
 457 *and boundary conditions are given in (4.11), while the processor speed is given by*

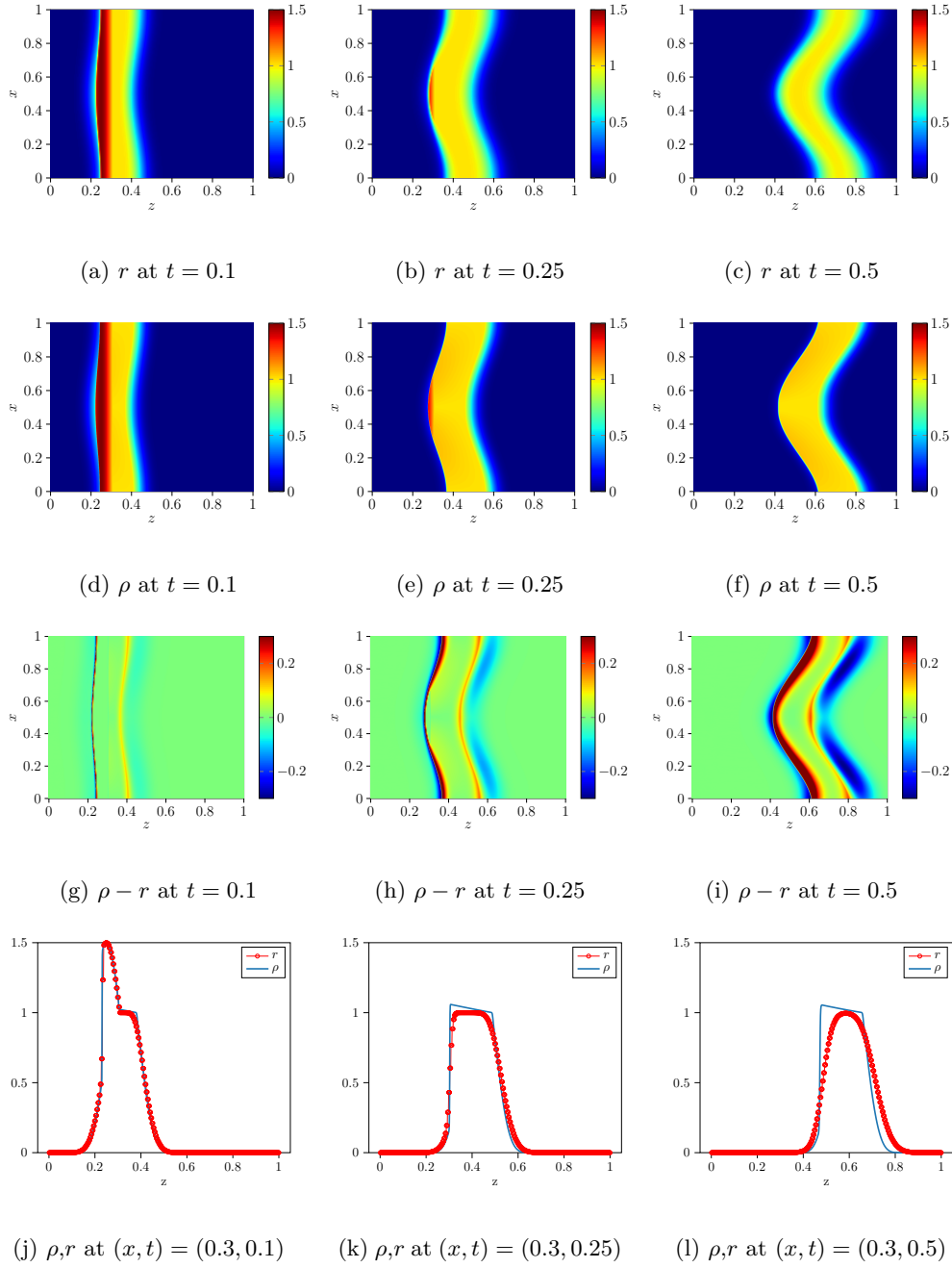


Fig. 5: Discrete solution r and continuum solution ρ when $\eta = 0.2$. From left to right, columns correspond to solutions at $t = 0.1, t = 0.25$, and $t = 0.5$. Discrete solution is computed with $(i^{\max}, k^{\max}) = (1000, 200)$. Continuum solution is computed on a $10^3 \times 10^3$ mesh.

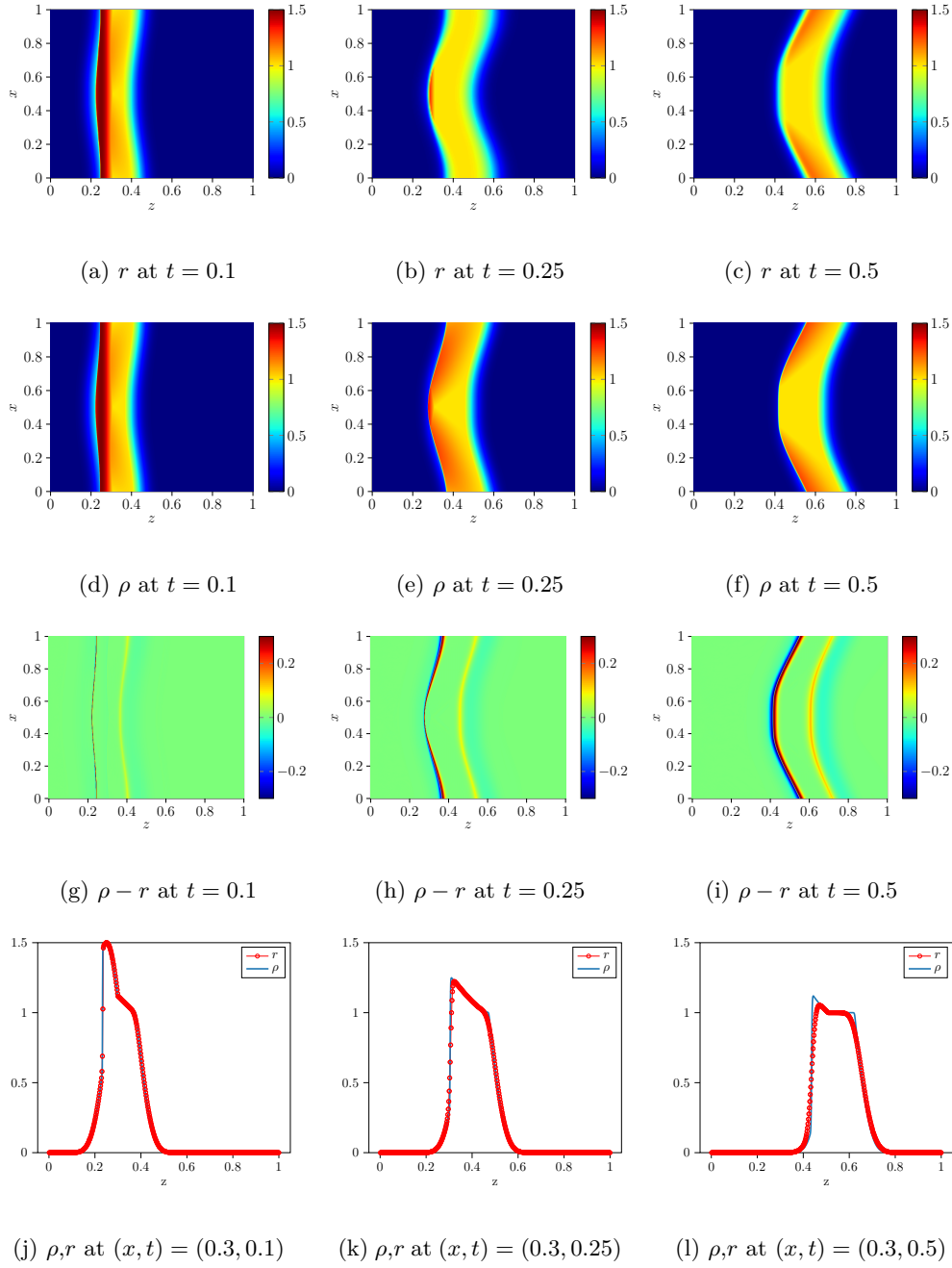


Fig. 6: Discrete r and continuum ρ solutions when $\eta = 1$ case. From left to right, column correspond to solutions at $t = 0.1$, $t = 0.25$, and $t = 0.5$. Discrete solution is computed with $(i^{\max}, k^{\max}) = (500, 500)$. Continuum solution is computed on a $10^3 \times 10^3$ mesh.

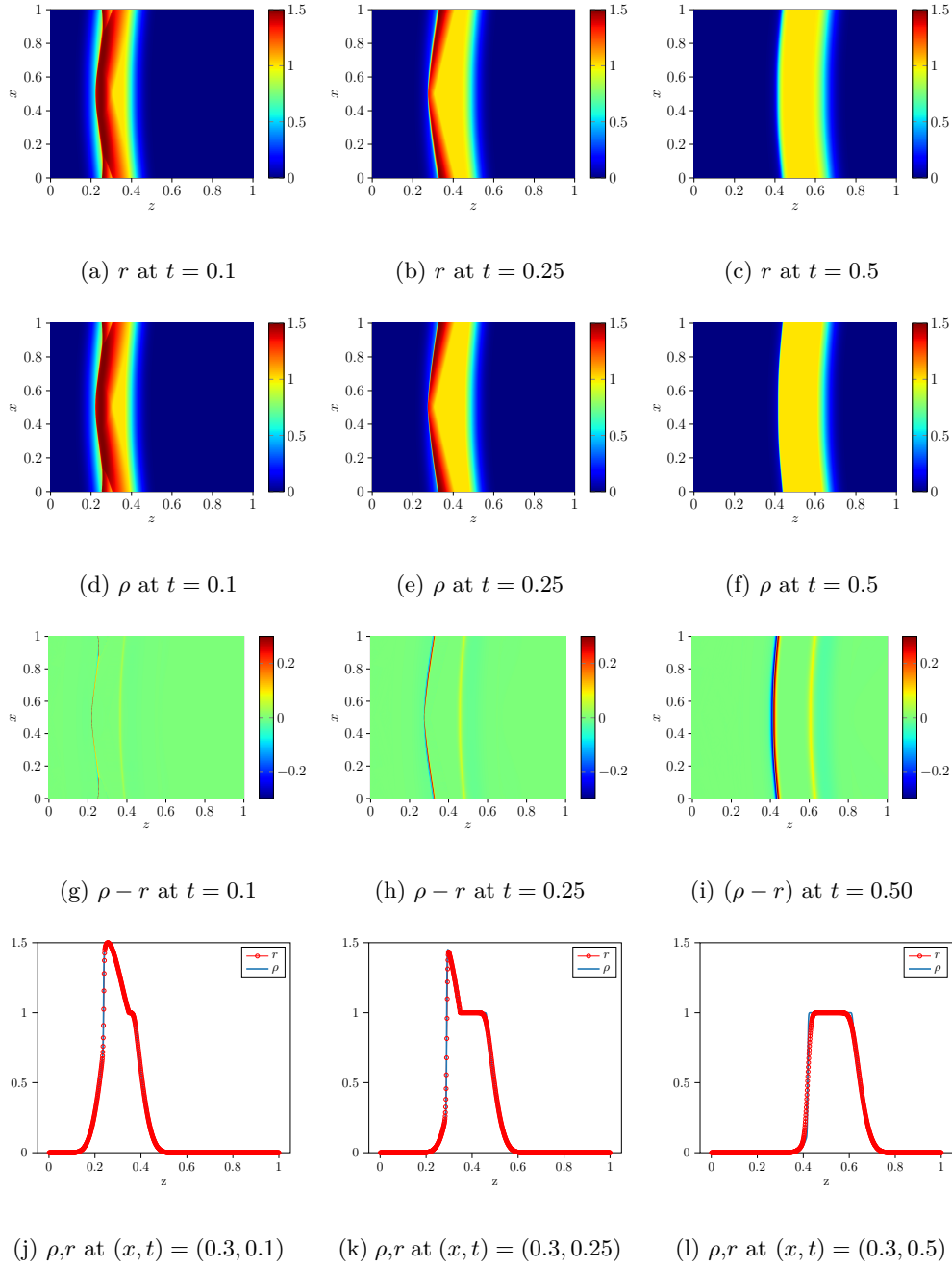


Fig. 7: Discrete r and continuum ρ solutions when $\eta = 5$ case. From left to right, column correspond to solutions at $t = 0.1$, $t = 0.25$, and $t = 0.5$. Discrete solution is computed with $(i^{\max}, k^{\max}) = (200, 1000)$. Continuum solution is computed on a $10^3 \times 10^3$ mesh.

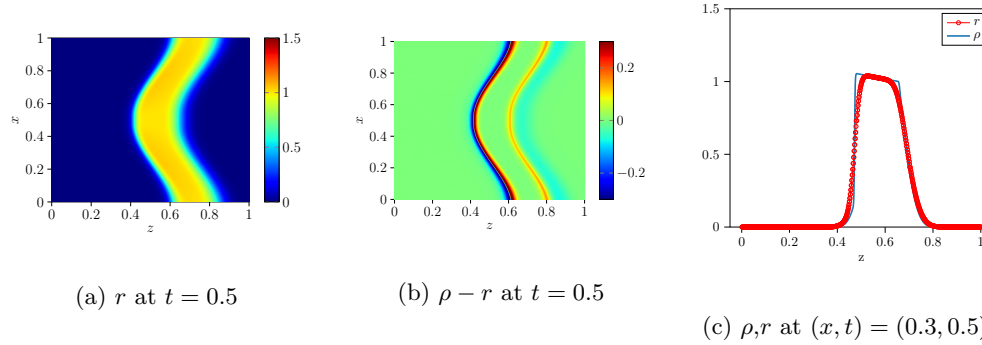


Fig. 8: Comparison of the discrete model with $(i^{\max}, k^{\max}) = (2500, 500)$ and the continuum model for $\eta = 0.2$ at time $t = 0.5$. As expected, the discrete model shows better agreement with the continuum model than the previous version with only $(i^{\max}, k^{\max}) = (1000, 200)$ processors and stages; cf. Figure 5

458 $\alpha(x) = 1 - 0.4c(x)$, where

$$459 \quad (4.12) \quad c(x) = \begin{cases} 0 & |x - .5| > .05 \\ 40x - 18 & x \in [.45, .475] \\ -40x + 22 & x \in [.525, .55] \\ 1 & |x - .5| < .025 \end{cases}.$$

460 In particular, $\alpha \neq 1$ only on the interval $(0.45, 0.55)$. Simulation results from this
 461 example are shown in Figure 11. At early times, slower processors in the center
 462 of the x domain prohibit neighboring processors from moving data to later stages of
 463 the calculation (i.e. along the z -direction). The result is a buildup of data in the
 464 neighboring processors. As time progresses, the build-up of data spreads as throttled
 465 processors near the initial slowdown around $x = 0.5$ begin to effect neighbors further
 466 away. Eventually these buildups dissipate as the slower processors begin catch up with
 467 their throttled neighbors.

468 EXAMPLE 5 (Long-term behavior). In previous examples, we have observed that
 469 under some conditions, solutions eventually resemble a traveling profile of the form

$$470 \quad (4.13) \quad \rho_*(x, z, t) = \chi_{[\zeta_0(x), \zeta_1(x)]}(z - st),$$

471 where s is a positive constant and the profiles ξ_0 and ξ_1 are constant in time and
 472 satisfy $\zeta_1(x) < \zeta_1(x)$ for all $x \in [0, 1]$. Our intuition is that for a wide range of
 473 conditions, traveling profiles of this type will arise after sufficiently long times, if
 474 the z domain is extended to $(0, \infty)$. Moreover the shape of ζ_1 and ζ_2 is closely related
 475 to the initial data and the shape of α .¹ Rather than make a precise conjecture at this
 476 point, we instead provide an example which further demonstrates our intuition. Initial
 477 and boundary conditions are given in (4.11). Because the domain in z is limited, we
 478 introduce relatively small variations in α , which allow the system to settle faster:

$$479 \quad (4.14) \quad \alpha(x) = 1 + 0.1 \cos(4\pi x).$$

¹A more systematic study of such profiles in special case can be found in [19].

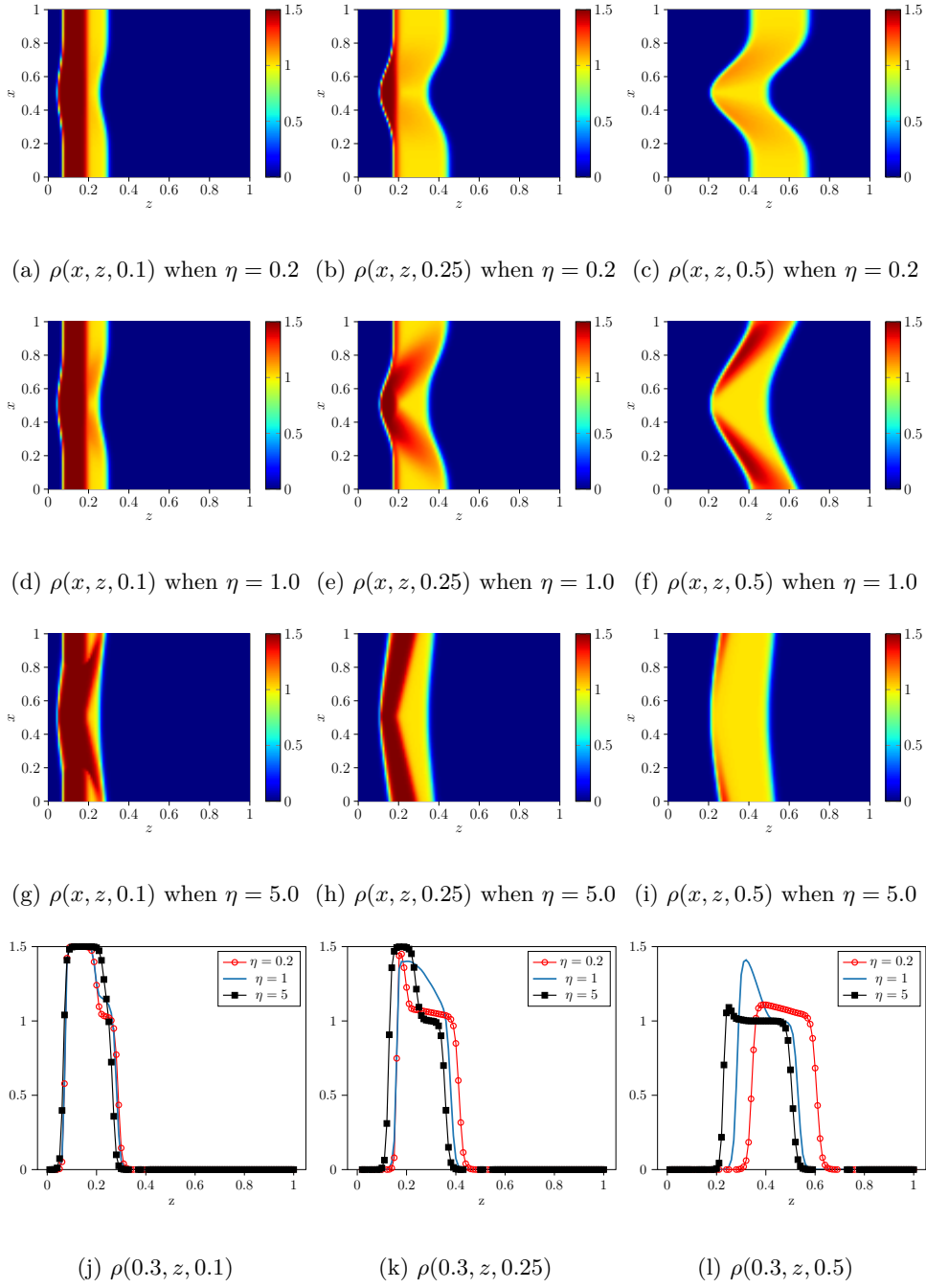


Fig. 9: The effects on ρ due to variations in η . As η increases the throttling effect of a local slowdown spreads more quickly, and data is not processed as quickly.

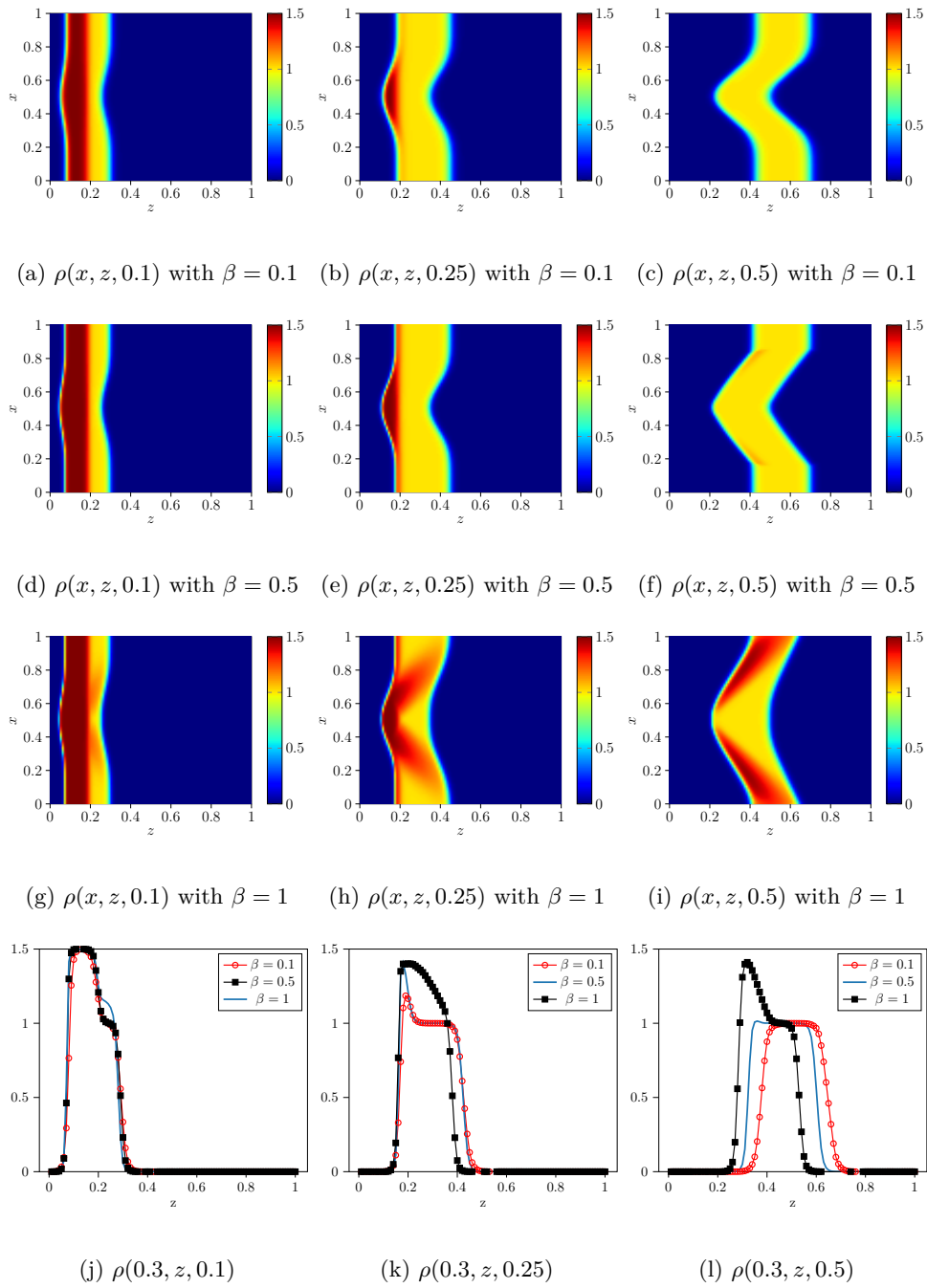


Fig. 10: Plots of the solution ρ from Example 3 for different values of β . Larger values of β lead to more throttling.

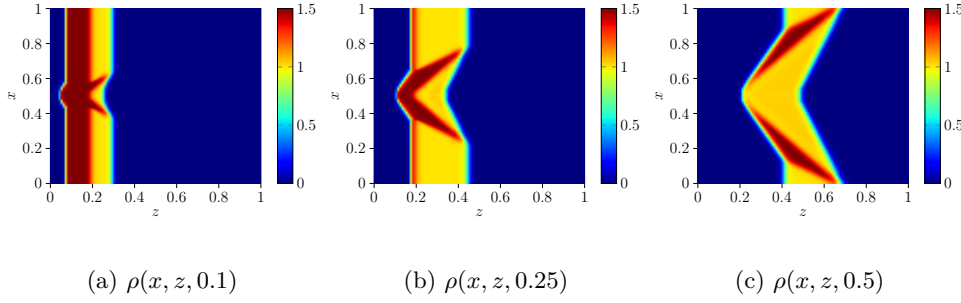


Fig. 11: The effect of a highly localized slowdown on ρ

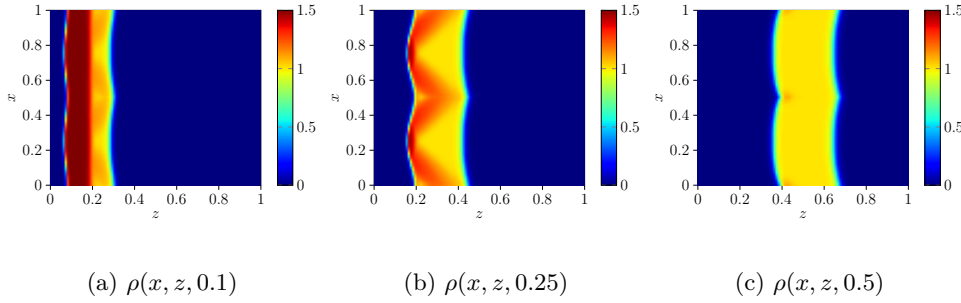


Fig. 12: The effect of small variation in processor speed on ρ . After sufficiently time, a profile emerges with the periodicity of α .

480 *Simulation results for this example are presented in Figure 12. When $t = 0.5$, the*
 481 *solution has nearly settled to a profile of the form (4.13), with cusps that appear where*
 482 *the waves caused by throttling meet, at $x = 0.5$ and at the periodic boundary. In*
 483 *particular the solution has the periodicity of α .*

484 **5. Conclusion.** We have presented a simple discrete model of a network of
 485 processors in a high performance computing environment where the computational
 486 throughput depends on the on the availability of data from neighboring processors.
 487 This discrete, microscopic-level model has been then used to derive a continuum-level
 488 model which treats computational progress as an Eulerian fluid flow. Currently, the
 489 existence and uniqueness of solutions to the partial differential equation in this fluid
 490 model is open. However, a Hamilton-Jacobi model is available for which we can es-
 491 tablish the existence and uniqueness of continuous viscosity solutions; the solution for
 492 the governing equation corresponds to the total amount of data that has been pro-
 493 cessed through a particular stage in the computation. Numerical experiments have
 494 shown that this continuum model can capture the asymptotic behavior of the discrete
 495 model. Additionally, we have used these experiments to give an initial understanding
 496 of solutions' dependence on parameters associated with the parallelism of the modeled
 497 computation as well as the effects heterogeneities in processing capacity.

498 In future work, we intend to explore control strategies for α that can alleviate
 499 bottlenecks caused by local slowdowns in the processor speed. We will also extend

500 the model to allow for more complicated interactions, including stochastic effects, and
 501 explore strategies for optimal communication. Finally, we hope to tune the parameters
 502 of the model with data taken from processor components of a real supercomputer and
 503 then compare predictions of the macroscopic model with the real global behavior of
 504 the supercomputer.

505 **Appendix A. Hamilton-Jacobi Theory.**

506 We recall a few standard definitions from the theory of nonlinear second-order
 507 Hamilton-Jacobi equations as in, for instance, [6, 11]:

508 **DEFINITION A.1** (Degenerately elliptic function). *Let $F : \mathbb{R}^n \times \mathbb{R}^n \times \mathcal{S}^n \rightarrow \mathbb{R}$ be*
 509 *given, where \mathcal{S}^n is the set of symmetric $n \times n$ matrices. Then we say F is degenerately*
 510 *elliptic if $F(x, r, p, X) \leq F(x, r, p, Y)$ whenever $Y \leq X$.*

511 **DEFINITION A.2** (Modulus function). *We call a function $\sigma : [0, \infty) \rightarrow [0, \infty)$ a*
 512 *modulus function if $\sigma(0) = 0$ and it is nondecreasing.*

513 In contrast to the notational convention in [Section 3](#), we follow in this appendix the
 514 convention of the viscosity literature and refer to time- and space-dependent functions
 515 as $u(t, x)$.

516 **DEFINITION A.3** (Parabolic boundary). *If $U = (0, T] \times D$ where $D \subset \mathbb{R}^n$ and*
 517 *$T \geq 0$, then $\partial_P U := \{0\} \times D \cup [0, T] \times \partial D$ is called the parabolic boundary of U .*

518 **DEFINITION A.4** (Semicontinuous envelope). *The upper (respectively, lower) semi-*
 519 *continuous envelopes of a function $u : V \rightarrow [-\infty, \infty]$ are*

$$520 \quad u^*(x) = \limsup_{r \downarrow 0} \{u(y) : y \in V, |y - x| \leq r\},$$

$$521 \quad u_*(x) = \liminf_{r \downarrow 0} \{u(y) : y \in V, |y - x| \leq r\}.$$

522 *They are, respectively, the smallest upper semicontinuous function greater than u and*
 523 *the largest lower semicontinuous function less than u .*

524 **DEFINITION A.5** (Viscosity solutions). *Let $f : U \times \mathbb{R} \times \mathbb{R}^n \times \mathcal{S}^n \rightarrow \mathbb{R}$ be given.*
 525 *An upper (resp. lower) semicontinuous function u is a viscosity subsolution (resp.*
 526 *supersolution) of*

$$528 \quad (\text{A.1a}) \quad u_t + f(t, x, u, \nabla_x u, \nabla_x^2 u) = 0, \quad (t, x) \in U,$$

$$529 \quad (\text{A.1b}) \quad h(t, x, u, \nabla_x u) = 0, \quad (t, x) \in \partial_P U,$$

531 *on $(0, T] \times D$ if at every $(t, x) \in (0, T] \times D$, when $u - \psi$ is locally maximized (resp.*
 532 *minimized) at (t, x) and ψ is $C^2((0, T] \times D)$ we have*

$$533 \quad (\text{A.2a}) \quad \psi_t(t, x) + f(t, x, u, \nabla_x \psi(t, x), \nabla_x^2 \psi(t, x)) \leq 0, \quad (t, x) \in U$$

$$534 \quad (\text{A.2b}) \quad \min\{\psi_t(t, x) + f(t, x, u, \nabla_x \psi(t, x), \nabla_x^2 \psi(t, x)),$$

$$535 \quad h(t, x, u(t, x), \nabla_x \psi(t, x))\} \leq 0, \quad (t, x) \in \partial_P U$$

537 *(respectively,*

$$538 \quad (\text{A.3a}) \quad \psi_t(t, x) + f(t, x, u, \nabla_x \psi(t, x), \nabla_x^2 \psi(t, x)) \geq 0, \quad (t, x) \in U$$

$$539 \quad (\text{A.3b}) \quad \max\{\psi_t(t, x) + f(t, x, u, \nabla_x \psi(t, x), \nabla_x^2 \psi(t, x)),$$

$$540 \quad h(t, x, u(t, x), \nabla_x \psi(t, x))\} \geq 0, \quad (t, x) \in \partial_P U.$$

541 *A function u is a viscosity solution of (A.1) if its upper semicontinuous envelope is a*
 542 *viscosity subsolution and its lower semicontinuous envelope is a viscosity supersolution.*

544 Next, we recall the following general comparison theorem, which is Theorem 4.1
545 of [17].

546 LEMMA A.6 (Generalized comparison principle). *Consider the system (A.1) on*
547 $U = (0, T) \times \Omega$ *where* $\Omega \subseteq \mathbb{R}^n$ *is a possibly unbounded domain and* $T > 0$. *Assume*
548 *that* f *satisfies the following assumptions:*

- 549 1. f *is continuous on* $U \times \mathbb{R} \times (\mathbb{R}^n \setminus \{0\}) \times \mathcal{S}^n$.
- 550 2. f *is degenerately elliptic.*
- 551 3. $-\infty < f_*(t, x, r, 0, O) = f^*(t, x, r, 0, O) < \infty$ *for all* $(t, x, r) \in U \times \mathbb{R}$, *where*
552 O *is the zero matrix.*
- 553 4. *For every* $R > 0$, *we have*

$$(A.4) \quad \sup\{|f(t, x, r, p, X)| : |p|, |X| \leq R, (t, x, r, p, S) \in U \times \mathbb{R} \times (\mathbb{R}^n \setminus \{0\}) \times \mathcal{S}^n\}$$

555 *is finite.*

- 556 5. *For every* $H > 0$, *there is a constant* c_0 *such that* $r \mapsto f(t, x, r, p, X) + c_0 r$ *is*
557 *nondecreasing for all* $(t, x, r, p, X) \in U \times \mathbb{R} \times (\mathbb{R}^n \setminus \{0\}) \times \mathcal{S}^n$ *with* $|r| \leq H$.
- 558 6. *For every* $R > \rho > 0$ *there is a modulus function* $\sigma_{R, \rho}$ *such that*

$$(A.5) \quad |f(t, x, r, p, X) - f(t, x, r, q, Y)| \leq \sigma_{R, \rho}(|p - q| + |X - Y|)$$

560 *for* $(t, x, r) \in U \times \mathbb{R}$, $\rho \leq |p|$, $|q| \leq R$, *and* $|X|, |Y| \leq R$.

- 561 7. *There is a constant* $\rho_0 > 0$ *and a modulus function* σ_1 *such that*

$$(A.6) \quad f^*(t, x, r, p, X) - f^*(t, x, r, 0, O) \leq \sigma_1(|p| + |X|)$$

$$(A.7) \quad f_*(t, x, r, p, X) - f_*(t, x, r, 0, O) \geq -\sigma_1(|p| + |X|)$$

565 *for* $(t, x, r) \in U \times \mathbb{R}$ *and* $|p|, |X| \leq \rho_0$.

- 566 8. *There is a modulus function* σ_2 *such that*

$$(A.8) \quad |f(t, x, r, p, X) - f(t, y, r, p, X)| \leq \sigma_2(|x - y|(|p| + 1))$$

568 *for any* $y \in \Omega$ *and* $(t, x, r, p, X) \in U \times \mathbb{R} \times (\mathbb{R}^n \setminus \{0\}) \times \mathcal{S}^n$.

569 *Then if* u^- *and* u^+ *are viscosity subsolutions and supersolutions of (A.1), respectively,*
570 *such that for some* $K > 0$ *independent of* $t, x, y \in (0, T] \times \Omega \times \Omega$:

- 571 • $u^-(t, x) \leq K(|x| + 1)$ *and* $u^+(t, x) \geq -K(|x| + 1)$;
- 572 • $(u^-)^*(t, x) - (u^+)_*(t, y) \leq m_T(|x - y|)$ *on* $\partial_p((0, T] \times (\Omega \times \Omega))$;
- 573 • $(u^-)^*(t, x) - (u^+)_*(t, y) \leq K(|x - y| + 1)$ *on* $\partial_p((0, T] \times (\Omega \times \Omega))$.

574 *Then there is a modulus function* σ *such that*

$$(A.9) \quad (u^-)^*(t, x) - (u^+)_*(t, y) \leq \sigma(|x - y|).$$

576 This generalized comparison principle, coupled with an argument which uses the
577 framework given in [11], known as Perron's method, gives the existence and unique-
578 ness of a viscosity solution to (3.28). Specifically, we note the parabolic version of
579 this framework uses a result like the following, which is Lemma 2.3.15 from [20]

580 LEMMA A.7 (Perron process for parabolic equations). *Consider (A.1) where* f
581 *is degenerate elliptic and continuous. Assume that* u^+ *and* u^- *are viscosity superso-*
582 *lutions and subsolutions, respectively. Then there exists a viscosity solution* u *such*
583 *that* $u^- \leq u \leq u^+$.

584 The results above are summarized in the following theorem.

585 THEOREM A.8 (Unique viscosity solution). *Suppose that f satisfies the condi-*
 586 *tions of Lemma A.6 and Lemma A.7 and that there exists a viscosity supersolution*
 587 *u^+ and a viscosity subsolution u^- to (A.1). Then there exists a unique continuous*
 588 *viscosity solution to (A.1).*

589 *Proof.* According to Lemma A.7, there exists a viscosity solution u to (A.1). To
 590 show uniqueness and continuity, let v be another viscosity solution. By definition, u^*
 591 and v^* are subsolutions and u_* and v_* are supersolutions. Then (A.9), combined with
 592 the properties of envelopes imply that

$$593 \text{ (A.10)} \quad u^* = (u^*)^* \leq (v_*)_* = v_* \leq v^* = (v^*)^* \leq (u_*)_* = u_* \leq u^*.$$

594 Thus $u = v$ and $u^* = u_*$ so that u is continuous. \square

595 **Acknowledgments.** C.D.H and R.C.B. would like to thank Michael Herty for
 596 many helpful discussions.

597

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